

Effective STL

50 Specific Ways to Improve
Your Use of the Standard
Template Library

Scott Meyers



Effective STL

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Scott Meyers



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Contents

Preface	xi
Acknowledgments	xv
Introduction	1
Chapter 1: Containers	11
Item 1: Choose your containers with care.	11
Item 2: Beware the illusion of container-independent code.	15
Item 3: Make copying cheap and correct for objects in containers.	20
Item 4: Call empty instead of checking size() against zero.	23
Item 5: Prefer range member functions to their single-element counterparts.	24
Item 6: Be alert for C++'s most vexing parse.	33
Item 7: When using containers of newed pointers, remember to delete the pointers before the container is destroyed.	36
Item 8: Never create containers of auto_ptrs.	40
Item 9: Choose carefully among erasing options.	43
Item 10: Be aware of allocator conventions and restrictions.	48
Item 11: Understand the legitimate uses of custom allocators.	54
Item 12: Have realistic expectations about the thread safety of STL containers.	58
Chapter 2: vector and string	63
Item 13: Prefer vector and string to dynamically allocated arrays.	63
Item 14: Use reserve to avoid unnecessary reallocations.	66
Item 15: Be aware of variations in string implementations.	68

Item 16:	Know how to pass vector and string data to legacy APIs.	74
Item 17:	Use “the swap trick” to trim excess capacity.	77
Item 18:	Avoid using <code>vector<bool></code> .	79

Chapter 3: Associative Containers **83**

Item 19:	Understand the difference between equality and equivalence.	83
Item 20:	Specify comparison types for associative containers of pointers.	88
Item 21:	Always have comparison functions return false for equal values.	92
Item 22:	Avoid in-place key modification in set and multiset.	95
Item 23:	Consider replacing associative containers with sorted vectors.	100
Item 24:	Choose carefully between <code>map::operator[]</code> and <code>map::insert</code> when efficiency is important.	106
Item 25:	Familiarize yourself with the nonstandard hashed containers.	111

Chapter 4: Iterators **116**

Item 26:	Prefer iterator to <code>const_iterator</code> , <code>reverse_iterator</code> , and <code>const_reverse_iterator</code> .	116
Item 27:	Use <code>distance</code> and <code>advance</code> to convert a container’s <code>const_iterators</code> to iterators.	120
Item 28:	Understand how to use a <code>reverse_iterator</code> ’s base iterator.	123
Item 29:	Consider <code>istreambuf_iterators</code> for character-by-character input.	126

Chapter 5: Algorithms **128**

Item 30:	Make sure destination ranges are big enough.	129
Item 31:	Know your sorting options.	133
Item 32:	Follow remove-like algorithms by <code>erase</code> if you really want to remove something.	139
Item 33:	Be wary of remove-like algorithms on containers of pointers.	143
Item 34:	Note which algorithms expect sorted ranges.	146
Item 35:	Implement simple case-insensitive string comparisons via <code>mismatch</code> or <code>lexicographical_compare</code> .	150
Item 36:	Understand the proper implementation of <code>copy_if</code> .	154

Item 37: Use accumulate or for_each to summarize ranges.	156
Chapter 6: Functors, Functor Classes, Functions, etc.	162
Item 38: Design functor classes for pass-by-value.	162
Item 39: Make predicates pure functions.	166
Item 40: Make functor classes adaptable.	169
Item 41: Understand the reasons for ptr_fun, mem_fun, and mem_fun_ref.	173
Item 42: Make sure less<T> means operator<.	177
Chapter 7: Programming with the STL	181
Item 43: Prefer algorithm calls to hand-written loops.	181
Item 44: Prefer member functions to algorithms with the same names.	190
Item 45: Distinguish among count, find, binary_search, lower_bound, upper_bound, and equal_range.	192
Item 46: Consider function objects instead of functions as algorithm parameters.	201
Item 47: Avoid producing write-only code.	206
Item 48: Always #include the proper headers.	209
Item 49: Learn to decipher STL-related compiler diagnostics.	210
Item 50: Familiarize yourself with STL-related web sites.	217
Bibliography	225
Appendix A: Locales and Case-Insensitive String Comparisons	229
Appendix B: Remarks on Microsoft's STL Platforms	239
Index	245

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Preface

*It came without ribbons! It came without tags!
It came without packages, boxes or bags!*

— Dr. Seuss, *How the Grinch Stole Christmas!*, Random House, 1957

I first wrote about the Standard Template Library in 1995, when I concluded the final Item of *More Effective C++* with a brief STL overview. I should have known better. Shortly thereafter, I began receiving mail asking when I'd write *Effective STL*.

I resisted the idea for several years. At first, I wasn't familiar enough with the STL to offer advice on it, but as time went on and my experience with it grew, this concern gave way to other reservations. There was never any question that the library represented a breakthrough in efficient and extensible design, but when it came to *using* the STL, there were practical problems I couldn't overlook. Porting all but the simplest STL programs was a challenge, not only because library implementations varied, but also because template support in the underlying compilers ranged from good to awful. STL tutorials were hard to come by, so learning "the STL way of programming" was difficult, and once that hurdle was overcome, finding comprehensible and accurate reference documentation was a challenge. Perhaps most daunting, even the smallest STL usage error often led to a blizzard of compiler diagnostics, each thousands of characters long, most referring to classes, functions, or templates not mentioned in the offending source code, almost all incomprehensible. Though I had great admiration for the STL and for the people behind it, I felt uncomfortable recommending it to practicing programmers. I wasn't sure it was *possible* to use the STL effectively.

Then I began to notice something that took me by surprise. Despite the portability problems, despite the dismal documentation, despite the compiler diagnostics resembling transmission line noise, many of

my consulting clients were using the STL anyway. Furthermore, they weren't just playing with it, they were using it in production code! That was a revelation. I knew that the STL featured an elegant design, but any library for which programmers are willing to endure portability headaches, poor documentation, and incomprehensible error messages has a lot more going for it than just good design. For an increasingly large number of professional programmers, I realized, even a bad implementation of the STL was preferable to no implementation at all.

Furthermore, I knew that the situation regarding the STL would only get better. Libraries and compilers would grow more conformant with the Standard (they have), better documentation would become available (it has — consult the bibliography beginning on [page 225](#)), and compiler diagnostics would improve (for the most part, we're still waiting, but [Item 49](#) offers suggestions for how to cope while we wait). I therefore decided to chip in and do my part for the STL movement. This book is the result: 50 specific ways to improve your use of C++'s Standard Template Library.

My original plan was to write the book in the second half of 1999, and with that thought in mind, I put together an outline. But then I changed course. I suspended work on the book and developed an introductory training course on the STL, which I then taught several times to groups of programmers. About a year later, I returned to the book, significantly revising the outline based on my experiences with the training course. In the same way that my *Effective C++* has been successful by being grounded in the problems faced by real programmers, it's my hope that *Effective STL* similarly addresses the practical aspects of STL programming — the aspects most important to professional developers.

I am always on the lookout for ways to improve my understanding of C++. If you have suggestions for new guidelines for STL programming or if you have comments on the guidelines in this book, please let me know. In addition, it is my continuing goal to make this book as accurate as possible, so for each error in this book that is reported to me — be it technical, grammatical, typographical, or otherwise — I will, in future printings, gladly add to the acknowledgments the name of the first person to bring that error to my attention. Send your suggested guidelines, your comments, and your criticisms to estl@aristeia.com.

I maintain a list of changes to this book since its first printing, including bug-fixes, clarifications, and technical updates. The list is available at the *Effective STL Errata* web site, <http://www.aristeia.com/BookErrata/estl1e-errata.html>.

If you'd like to be notified when I make changes to this book, I encourage you to join my mailing list. I use the list to make announcements likely to be of interest to people who follow my work on C++. For details, consult <http://www.aristeia.com/MailingList/>.

SCOTT DOUGLAS MEYERS
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Acknowledgments

I had an enormous amount of help during the roughly two years it took me to make some sense of the STL, create a training course on it, and write this book. Of all my sources of assistance, two were particularly important. The first is Mark Rodgers. Mark generously volunteered to review my training materials as I created them, and I learned more about the STL from him than from anybody else. He also acted as a technical reviewer for this book, again providing observations and insights that improved virtually every Item.

The other outstanding source of information was several C++-related Usenet newsgroups, especially `comp.lang.c++.moderated` (“clcm”), `comp.std.c++`, and `microsoft.public.vc.stl`. For well over a decade, I’ve depended on the participants in newsgroups like these to answer my questions and challenge my thinking, and it’s difficult to imagine what I’d do without them. I am deeply grateful to the Usenet community for their help with both this book and my prior publications on C++.

My understanding of the STL was shaped by a variety of publications, the most important of which are listed in the Bibliography. I leaned especially heavily on Josuttis’ *The C++ Standard Library* [3].

This book is fundamentally a summary of insights and observations made by others, though a few of the ideas are my own. I’ve tried to keep track of where I learned what, but the task is hopeless, because a typical Item contains information garnered from many sources over a long period of time. What follows is incomplete, but it’s the best I can do. Please note that my goal here is to summarize where *I* first learned of an idea or technique, not where the idea or technique was originally developed or who came up with it.

In [Item 1](#), my observation that node-based containers offer better support for transactional semantics is based on section 5.11.2 of Josuttis’ *The C++ Standard Library* [3]. [Item 2](#) includes an example from Mark Rodgers on how typedefs help when allocator types are changed.

[Item 5](#) was motivated by Reeves' C++ *Report* column, "STL Gotchas" [17]. [Item 8](#) sprang from [Item 37](#) in Sutter's *Exceptional C++* [8], and Kevlin Henney provided important details on how containers of `auto_ptr`s fail in practice. In Usenet postings, Matt Austern provided examples of when allocators are useful, and I include his examples in [Item 11](#). [Item 12](#) is based on the discussion of thread safety at the SGI STL web site [21]. The material in [Item 13](#) on the performance implications of reference counting in a multithreaded environment is drawn from Sutter's writings on this topic [20]. The idea for [Item 15](#) came from Reeves' C++ *Report* column, "Using Standard string in the Real World, Part 2," [18]. In [Item 16](#), Mark Rodgers came up with the technique I show for having a C API write data directly into a vector. [Item 17](#) includes information from Usenet postings by Siemel Naran and Carl Barron. I stole [Item 18](#) from Sutter's C++ *Report* column, "When Is a Container Not a Container?" [12]. In [Item 20](#), Mark Rodgers contributed the idea of transforming a pointer into an object via a dereferencing functor, and Scott Lewandowski came up with the version of `DereferenceLess` I present. [Item 21](#) originated in a Doug Harrison posting to `microsoft.public.vc.stl`, but the decision to restrict the focus of that Item to equality was mine. I based [Item 22](#) on Sutter's C++ *Report* column, "Standard Library News: sets and maps" [13]; Matt Austern helped me understand the status of the Standardization Committee's Library Issue #103. [Item 23](#) was inspired by Austern's C++ *Report* article, "Why You Shouldn't Use `set` — and What to Use Instead" [15]; David Smallberg provided a neat refinement for my implementation of `DataCompare`. My description of Dinkumware's hashed containers is based on Plauger's *C/C++ Users Journal* column, "Hash Tables" [16]. Mark Rodgers doesn't agree with the overall advice of [Item 26](#), but an early motivation for that Item was his observation that some container member functions accept only arguments of type iterator. My treatment of [Item 29](#) was motivated and informed by Usenet discussions involving Matt Austern and James Kanze; I was also influenced by Kreft and Langer's C++ *Report* article, "A Sophisticated Implementation of User-Defined Inserters and Extractors" [25]. [Item 30](#) is due to a discussion in section 5.4.2 of Josuttis' *The C++ Standard Library* [3]. In [Item 31](#), Marco Dalla Gasperina contributed the example use of `nth_element` to calculate medians, and use of that algorithm for finding percentiles comes straight out of section 18.7.1 of Stroustrup's *The C++ Programming Language* [7]. [Item 32](#) was influenced by the material in section 5.6.1 of Josuttis' *The C++ Standard Library* [3]. [Item 35](#) originated in Austern's C++ *Report* column "How to Do Case-Insensitive String Comparison" [11], and James Kanze's and John Potter's `clcm` postings helped me refine my understanding of the issues involved. Stroustrup's implementation for `copy_if`, which I

show in [Item 36](#), is from section 18.6.1 of his *The C++ Programming Language* [7]. [Item 39](#) was largely motivated by the publications of Josuttis, who has written about “stateful predicates” in his *The C++ Standard Library* [3], in Standard Library Issue #92, and in his C++ Report article, “Predicates vs. Function Objects” [14]. In my treatment, I use his example and recommend a solution he has proposed, though the use of the term “pure function” is my own. Matt Austern confirmed my suspicion in [Item 41](#) about the history of the terms `mem_fun` and `mem_fun_ref`. [Item 42](#) can be traced to a lecture I got from Mark Rodgers when I considered violating that guideline. Mark Rodgers is also responsible for the insight in [Item 44](#) that non-member searches over maps and multimaps examine both components of each pair, while member searches examine only the first (key) component. [Item 45](#) contains information from various clcm contributors, including John Potter, Marcin Kasperski, Pete Becker, Dennis Yelle, and David Abrahams. David Smallberg alerted me to the utility of `equal_range` in performing equivalence-based searches and counts over sorted sequence containers. Andrei Alexandrescu helped me understand the conditions under which “the reference-to-reference problem” I describe in [Item 50](#) arises, and I modeled my example of the problem on a similar example provided by Mark Rodgers at the Boost Web Site [22].

Credit for the material in [Appendix A](#) goes to Matt Austern, of course. I’m grateful that he not only gave me permission to include it in this book, he also tweaked it to make it even better than the original.

Good technical books require a thorough pre-publication vetting, and I was fortunate to benefit from the insights of an unusually talented group of technical reviewers. Brian Kernighan and Cliff Green offered early comments on a partial draft, and complete versions of the manuscript were scrutinized by Doug Harrison, Brian Kernighan, Tim Johnson, Francis Glassborow, Andrei Alexandrescu, David Smallberg, Aaron Campbell, Jared Manning, Herb Sutter, Stephen Dewhurst, Matt Austern, Gillmer Derge, Aaron Moore, Thomas Becker, Victor Von, and, of course, Mark Rodgers. Katrina Avery did the copyediting.

One of the most challenging parts of preparing a book is finding good technical reviewers. I thank John Potter for introducing me to Jared Manning and Aaron Campbell.

Herb Sutter kindly agreed to act as my surrogate in compiling, running, and reporting on the behavior of some STL test programs under a beta version of Microsoft’s Visual Studio .NET, while Leor Zolman undertook the herculean task of testing all the code in this book. Any errors that remain are my fault, of course, not Herb’s or Leor’s.

Angelika Langer opened my eyes to the indeterminate status of some aspects of STL function objects. This book has less to say about function objects than it otherwise might, but what it does say is more likely to remain true. At least I hope it is.

This printing of the book is better than earlier printings, because I was able to address problems identified by the following sharp-eyed readers: Jon Webb, Michael Hawkins, Derek Price, Jim Scheller, Carl Manaster, Herb Sutter, Albert Franklin, George King, Dave Miller, Harold Howe, John Fuller, Tim McCarthy, John Hershberger, Igor Mikolic-Torreira, Stephan Bergmann, Robert Allan Schwartz, John Potter, David Grigsby, Sanjay Pattni, Jesper Andersen, Jing Tao Wang, André Blavier, Dan Schmidt, Bradley White, Adam Petersen, Wayne Goertel, Gabriel Netterdag, Jason Kenny, Scott Blachowicz, Seyed H. Haeri, Gareth McCaughan, Giulio Agostini, Fraser Ross, Wolfram Burkhardt, Keith Stanley, Leor Zolman, Chan Ki Lok, Motti Abramsky, Kevlin Henney, Stefan Kuhlins, Phillip Ngan, Jim Phillips, Ruediger Dreier, Guru Chandar, Charles Brockman, Day Barr, Eric Niebler, Sharad Kala, Declan Moran, Nick de Smith, David Callaway, Shlomi Frank, Andrea Griffini, Hans Eckardt, David Smallberg, Matt Page, Andy Fyfe, Vincent Stojanov, Randy Parker, Thomas Schell, Cameron Mac Minn, Mark Davis, Giora Unger, Julie Nahil, Martin Rottinger, Neil Henderson, Andrew Savige, and Molly Sharp. I'm grateful for their help in improving *Effective STL*.

My collaborators at Addison-Wesley included John Wait (my editor and now a senior VP), Alicia Carey and Susannah Buzard (his assistants n and $n+1$), John Fuller (the production coordinator), Karin Hansen (the cover designer), Jason Jones (all-around technical guru, especially with respect to the demonic software spewed forth by Adobe), Marty Rabinowitz (their boss, but he works, too), and Curt Johnson, Chanda Leary-Coutu, and Robin Bruce (all marketing people, but still very nice).

Abbi Staley made Sunday lunches a routinely pleasurable experience.

As she has for the six books and one CD that came before it, my wife, Nancy, tolerated the demands of my research and writing with her usual forbearance and offered me encouragement and support when I needed it most. She never fails to remind me that there's more to life than C++ and software.

And then there's our dog, Persephone. As I write this, it is her sixth birthday. Tonight, she and Nancy and I will visit Baskin-Robbins for ice cream. Persephone will have vanilla. One scoop. In a cup. To go.

Introduction

You're already familiar with the STL. You know how to create containers, iterate over their contents, add and remove elements, and apply common algorithms, such as `find` and `sort`. But you're not satisfied. You can't shake the sensation that the STL offers more than you're taking advantage of. Tasks that should be simple aren't. Operations that should be straightforward leak resources or behave erratically. Procedures that should be efficient demand more time or memory than you're willing to give them. Yes, you know how to use the STL, but you're not sure you're using it *effectively*.

I wrote this book for you.

In *Effective STL*, I explain how to combine STL components to take full advantage of the library's design. Such information allows you to develop simple, straightforward solutions to simple, straightforward problems, and it also helps you design elegant approaches to more complicated problems. I describe common STL usage errors, and I show you how to avoid them. That helps you dodge resource leaks, code that won't port, and behavior that is undefined. I discuss ways to optimize your code, so you can make the STL perform like the fast, sleek machine it is intended to be.

The information in this book will make you a better STL programmer. It will make you a more productive programmer. And it will make you a happier programmer. Using the STL is fun, but using it effectively is outrageous fun, the kind of fun where they have to drag you away from the keyboard, because you just can't believe the good time you're having. Even a cursory glance at the STL reveals that it is a wondrously cool library, but the coolness runs broader and deeper than you probably imagine. One of my primary goals in this book is to convey to you just how amazing the library is, because in the nearly 30 years I've been programming, I've never seen anything like the STL. You probably haven't either.

Defining, Using, and Extending the STL

There is no official definition of “the STL,” and different people mean different things when they use the term. In this book, “the STL” means the parts of C++’s Standard Library that work with iterators. That includes the standard containers (including string), parts of the iostream library, function objects, and algorithms. It excludes the standard container adapters (stack, queue, and priority_queue) as well as the containers bitset and valarray, because they lack iterator support. It doesn’t include arrays, either. True, arrays support iterators in the form of pointers, but arrays are part of the C++ *language*, not the library.

Technically, my definition of the STL excludes extensions of the standard C++ library, notably hashed containers, singly linked lists, ropes, and a variety of nonstandard function objects. Even so, an effective STL programmer needs to be aware of such extensions, so I mention them where it’s appropriate. Indeed, [Item 25](#) is devoted to an overview of nonstandard hashed containers. They’re not in the STL now, but something similar to them is almost certain to make it into the next version of the standard C++ library, and there’s value in glimpsing the future.

One of the reasons for the existence of STL extensions is that the STL is a library designed to be extended. In this book, however, I focus on *using* the STL, not on adding new components to it. You’ll find, for example, that I have little to say about writing your own algorithms, and I offer no guidance at all on writing new containers and iterators. I believe that it’s important to master what the STL already provides before you embark on increasing its capabilities, so that’s what I focus on in *Effective STL*. When you decide to create your own STL-esque components, you’ll find advice on how to do it in books like Josuttis’ *The C++ Standard Library* [\[3\]](#) and Austern’s *Generic Programming and the STL* [\[4\]](#). One aspect of STL extension I *do* discuss in this book is writing your own function objects. You can’t use the STL effectively without knowing how to do that, so I’ve devoted an entire chapter to the topic ([Chapter 6](#)).

Citations

The references to the books by Josuttis and Austern in the preceding paragraph demonstrate how I handle bibliographic citations. In general, I try to mention enough of a cited work to identify it for people who are already familiar with it. If you already know about these authors’ books, for example, you don’t have to turn to the Bibliography to find out that [\[3\]](#) and [\[4\]](#) refer to books you already know. If you’re

not familiar with a publication, of course, the Bibliography (which begins on [page 225](#)) gives you a full citation.

I cite three works often enough that I generally leave off the citation number. The first of these is the International Standard for C++ [5], which I usually refer to as simply “the Standard.” The other two are my earlier books on C++, *Effective C++* [1] and *More Effective C++* [2].

The STL and Standards

I refer to the C++ Standard frequently, because *Effective STL* focuses on portable, standard-conformant C++. In theory, everything I show in this book will work with every C++ implementation. In practice, that isn’t true. Shortcomings in compiler and STL implementations conspire to prevent some valid code from compiling or from behaving the way it’s supposed to. Where that is commonly the case, I describe the problems, and I explain how you can work around them.

Sometimes, the easiest workaround is to use a different STL implementation. [Appendix B](#) gives an example of when this is the case. In fact, the more you work with the STL, the more important it becomes to distinguish between your *compilers* and your *library implementations*. When programmers run into difficulties trying to get legitimate code to compile, it’s customary for them to blame their compilers, but with the STL, compilers can be fine, while STL implementations are faulty. To emphasize the fact that you are dependent on both your compilers and your library implementations, I refer to your *STL platforms*. An STL platform is the combination of a particular compiler and a particular STL implementation. In this book, if I mention a compiler problem, you can be sure that I mean it’s the compiler that’s the culprit. However, if I refer to a problem with your STL platform, you should interpret that as “maybe a compiler bug, maybe a library bug, possibly both.”

I generally refer to your “compilers” — *plural*. That’s an outgrowth of my longstanding belief that you improve the quality (especially the portability) of your code if you ensure that it works with more than one compiler. Furthermore, using multiple compilers generally makes it easier to unravel the Gordian nature of error messages arising from improper use of the STL. ([Item 49](#) is devoted to approaches to decoding such messages.)

Another aspect of my emphasis on standard-conforming code is my concern that you avoid constructs with undefined behavior. Such constructs may do anything at runtime. Unfortunately, this means they may do precisely what you want them to, and that can lead to a false

sense of security. Too many programmers assume that undefined behavior always leads to an obvious problem, e.g., a segmentation fault or other catastrophic failure. The results of undefined behavior can actually be much more subtle, e.g., corruption of rarely-referenced data. They can also vary across program runs. I find that a good working definition of undefined behavior is “works for me, works for you, works during development and QA, but blows up in your most important customer’s face.” It’s important to avoid undefined behavior, so I point out common situations where it can arise. You should train yourself to be alert for such situations.

Reference Counting

It’s close to impossible to discuss the STL without mentioning reference counting. As you’ll see in Items 7 and 33, designs based on containers of pointers almost invariably lead to reference counting. In addition, many string implementations are internally reference counted, and, as Item 15 explains, this may be an implementation detail you can’t afford to ignore. In this book, I assume that you are familiar with the basics of reference counting. If you’re not, most intermediate and advanced C++ texts cover the topic. In *More Effective C++*, for example, the relevant material is in Items 28 and 29. If you don’t know what reference counting is and you have no inclination to learn, don’t worry. You’ll get through this book just fine, though there may be a few sentences here and there that won’t make as much sense as they otherwise would.

string and wstring

Whatever I say about string applies equally well to its wide-character counterpart, wstring. Similarly, any time I refer to the relationship between string and char or char*, the same is true of the relationship between wstring and wchar_t or wchar_t*. In other words, just because I don’t explicitly mention wide-character strings in this book, don’t assume that the STL fails to support them. It supports them as well as char-based strings. It has to. Both string and wstring are instantiations of the same template, basic_string.

Terms, Terms, Terms

This is not an introductory book on the STL, so I assume you know the fundamentals. Still, the following terms are sufficiently important that I feel compelled to review them:

- vector, string, deque, and list are known as the *standard sequence containers*. The *standard associative containers* are set, multiset, map, and multimap.
- Iterators are divided into five categories, based on the operations they support. Very briefly, *input iterators* are read-only iterators where each iterated location may be read only once. *Output iterators* are write-only iterators where each iterated location may be written only once. Input and output iterators are modeled on reading and writing input and output streams (e.g., files). It's thus unsurprising that the most common manifestations of input and output iterators are `istream_iterators` and `ostream_iterators`, respectively.

Forward iterators have the capabilities of both input and output iterators, but they can read or write a single location repeatedly. They don't support operator--, so they can move only forward with any degree of efficiency. All standard STL containers support iterators that are more powerful than forward iterators, but, as you'll see in [Item 25](#), one design for hashed containers yields forward iterators. Containers for singly linked lists (considered in [Item 50](#)) also offer forward iterators.

Bidirectional iterators are just like forward iterators, except they can go backward as easily as they go forward. The standard associative containers all offer bidirectional iterators. So does list.

Random access iterators do everything bidirectional iterators do, but they also offer "iterator arithmetic," i.e., the ability to jump forward or backward in a single step. vector, string, and deque each provide random access iterators. Pointers into arrays act as random access iterators for the arrays.

- Any class that overloads the function call operator (i.e., operator()) is a *functor class*. Objects created from such classes are known as *function objects* or *functors*. Most places in the STL that work with function objects work equally well with real functions, so I often use the term "function objects" to mean both C++ functions as well as true function objects.
- The functions `bind1st` and `bind2nd` are known as *binders*.

A revolutionary aspect of the STL is its complexity guarantees. These guarantees bound the amount of work any STL operation is allowed to perform. This is wonderful, because it can help you determine the relative efficiency of different approaches to the same problem, regardless of the STL platform you're using. Unfortunately, the terminology

behind the complexity guarantees can be confusing if you haven't been formally introduced to the jargon of computer science. Here's a quick primer on the complexity terms I use in this book. Each refers to how long it takes to do something as a function of n , the number of elements in a container or range.

- An operation that runs in *constant time* has performance that is unaffected by changes in n . For example, inserting an element into a list is a constant-time operation. Regardless of whether the list has one element or one million, the insertion takes about the same amount of time.

Don't take the term "constant time" too literally. It doesn't mean that the amount of time it takes to do something is literally constant, it just means that it's unaffected by n . For example, two STL platforms might take dramatically different amounts of time to perform the same "constant-time" operation. This could happen if one library has a much more sophisticated implementation than another or if one compiler performs substantially more aggressive optimization.

A variant of constant time complexity is *amortized constant time*. Operations that run in amortized constant time are usually constant-time operations, but occasionally they take time that depends on n . Amortized constant time operations *typically* run in constant time.

- An operation that runs in *logarithmic time* needs more time to run as n gets larger, but the time it requires grows at a rate proportional to the logarithm of n . For example, an operation on a million items would be expected to take only about three times as long as on a hundred items, because $\log n^3 = 3 \log n$. Most search operations on associative containers (e.g., `set::find`) are logarithmic-time operations.
- The time needed to perform an operation that runs in *linear time* increases at a rate proportional to increases in n . The standard algorithm `count` runs in linear time, because it has to look at every element of the range it's given. If the range triples in size, it has to do three times as much work, and we'd expect it to take about three times as long to do it.

As a general rule, a constant-time operation runs faster than one requiring logarithmic time, and a logarithmic-time operation runs faster than one whose performance is linear. This is always true when n gets big enough, but for relatively small values of n , it's sometimes possible for an operation with a worse theoretical complexity to outperform an operation with a better theoretical complexity. If you'd like to know more about STL complexity guarantees, turn to Josuttis' *The C++ Standard Library* [3].

As a final note on terminology, recall that each element in a map or multimap has two components. I generally call the first component the *key* and the second component the *value*. Given

```
map<string, double> m;
```

for example, the string is the key and the double is the value.

Code Examples

This book is filled with example code, and I explain each example when I introduce it. Still, it's worth knowing a few things in advance.

You can see from the map example above that I routinely omit `#includes` and ignore the fact that STL components are in namespace `std`. When defining the map `m`, I could have written this,

```
#include <map>
#include <string>

using std::map;
using std::string;

map<string, double> m;
```

but I prefer to save us both the noise.

When I declare a formal type parameter for a template, I use `typename` instead of `class`. That is, instead of writing this,

```
template<class T>
class Widget { ... };
```

I write this:

```
template<typename T>
class Widget { ... };
```

In this context, `class` and `typename` mean exactly the same thing, but I find that `typename` more clearly expresses what I usually want to say: that *any* type will do; `T` need not be a class. If you prefer to use `class` to declare type parameters, go right ahead. Whether to use `typename` or `class` in this context is purely a matter of style.

It is not a matter of style in a different context. To avoid potential parsing ambiguities (the details of which I'll spare you), you are required to use `typename` to precede type names that are dependent on formal type parameters. Such types are known as *dependent types*, and an example will help clarify what I'm talking about. Suppose you'd like to write a template for a function that, given an STL container, returns whether the last element in the container is greater than the first element. Here's one way to do it:

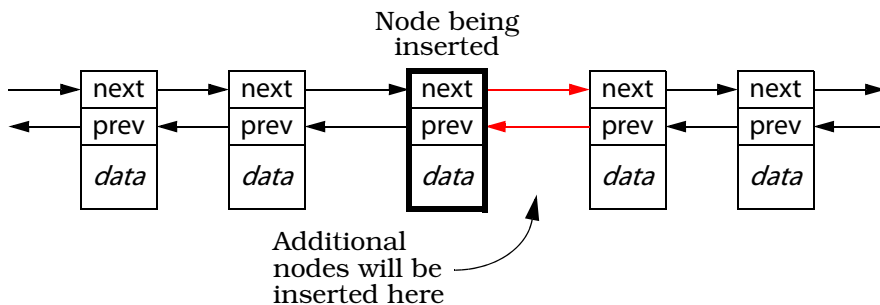
```

template<typename C>
bool lastGreaterThanFirst(const C& container)
{
    if (container.empty()) return false;
    typename C::const_iterator begin(container.begin());
    typename C::const_iterator end(container.end());
    return *--end > *begin;
}

```

In this example, the local variables `begin` and `end` are of type `C::const_iterator`. `const_iterator` is a type that is dependent on the formal type parameter `C`. Because `C::const_iterator` is a dependent type, you are required to precede it with the word `typename`. (Some compilers incorrectly accept the code without the `typename`s, but such code isn't portable.)

I hope you've noticed my use of color in the examples above. It's there to focus your attention on parts of the code that are particularly important. Often, I highlight the differences between related examples, such as when I showed the two possible ways to declare the parameter `T` in the `Widget` example. This use of color to call out especially noteworthy parts of examples carries over to diagrams, too. For instance, this diagram from [Item 5](#) uses color to identify the two pointers that are affected when a new element is inserted into a list:



I also use color for chapter numbers, but such use is purely gratuitous. This being my first two-color book, I hope you'll forgive me a little chromatic exuberance.

Two of my favorite parameter names are `lhs` and `rhs`. They stand for “left-hand side” and “right-hand side,” respectively, and I find them especially useful when declaring operators. Here's an example from [Item 19](#):

```

class Widget { ... };
bool operator==(const Widget& lhs, const Widget& rhs);

```

When this function is called in a context like this,

```
if (x == y) ... // assume x and y are Widgets
```

`x`, which is on the left-hand side of the “==”, is known as lhs inside operator==, and `y` is known as rhs.

As for the class name `Widget`, that has nothing to do with GUIs or tool-kits. It's just the name I use for “some class that does something.” Sometimes, as on [page 7](#), `Widget` is a class template instead of a class. In such cases, you may find that I still refer to `Widget` as a class, even though it's really a template. Such sloppiness about the difference between classes and class templates, structs and struct templates, and functions and function templates hurts no one as long as there is no ambiguity about what is being discussed. In cases where it could be confusing, I do distinguish between templates and the classes, structs, and functions they generate.

Efficiency Items

I considered including a chapter on efficiency in *Effective STL*, but I ultimately decided that the current organization was preferable. Still, a number of Items focus on minimizing space and runtime demands. For your performance-enhancing convenience, here is the table of contents for the virtual chapter on efficiency:

Item 4 :	Call <code>empty</code> instead of checking <code>size()</code> against zero.	23
Item 5 :	Prefer range member functions to their single-element counterparts.	24
Item 14 :	Use <code>reserve</code> to avoid unnecessary reallocations.	66
Item 15 :	Be aware of variations in string implementations.	68
Item 23 :	Consider replacing associative containers with sorted vectors.	100
Item 24 :	Choose carefully between <code>map::operator[]</code> and <code>map::insert</code> when efficiency is important.	106
Item 25 :	Familiarize yourself with the nonstandard hashed containers.	111
Item 29 :	Consider <code>istreambuf_iterator</code> s for character-by-character input.	126
Item 31 :	Know your sorting options.	133
Item 44 :	Prefer member functions to algorithms with the same names.	190
Item 46 :	Consider function objects instead of functions as algorithm parameters.	201

The Guidelines in *Effective STL*

The guidelines that make up the 50 Items in this book are based on the insights and advice of the world's most experienced STL programmers. These guidelines summarize things you should almost always do — or almost always avoid doing — to get the most out of the Standard Template Library. At the same time, they're just guidelines. Under some conditions, it makes sense to violate them. For example, the title of [Item 7](#) tells you to invoke `delete` on newed pointers in a container before the container is destroyed, but the text of that Item makes clear that this applies only when the objects pointed to by those pointers should go away when the container does. This is often the case, but it's not universally true. Similarly, the title of [Item 35](#) beseeches you to use STL algorithms to perform simple case-insensitive string comparisons, but the text of the Item points out that in some cases, you'll be better off using a function that's not only outside the STL, it's not even part of standard C++!

Only you know enough about the software you're writing, the environment in which it will run, and the context in which it's being created to determine whether it's reasonable to violate the guidelines I present. Most of the time, it won't be, and the discussions that accompany each Item explain why. In a few cases, it will. Slavish devotion to the guidelines isn't appropriate, but neither is cavalier disregard. Before venturing off on your own, you should make sure you have a good reason.

1

Containers

Sure, the STL has iterators, algorithms, and function objects, but for most C++ programmers, it's the containers that stand out. More powerful and flexible than arrays, they grow (and often shrink) dynamically, manage their own memory, keep track of how many objects they hold, bound the algorithmic complexity of the operations they support, and much, much more. Their popularity is easy to understand. They're simply better than their competition, regardless of whether that competition comes from containers in other libraries or is a container type you'd write yourself. STL containers aren't just good. They're *really* good.

This chapter is devoted to guidelines applicable to all the STL containers. Later chapters focus on specific container types. The topics addressed here include selecting the appropriate container given the constraints you face; avoiding the delusion that code written for one container type is likely to work with other container types; the significance of copying operations for objects in containers; difficulties that arise when pointers or `auto_ptr`s are stored in containers; the ins and outs of erasing; what you can and cannot accomplish with custom allocators; tips on how to maximize efficiency; and considerations for using containers in a threaded environment.

That's a lot of ground to cover, but don't worry. The Items break it down into bite-sized chunks, and along the way, you're almost sure to pick up several ideas you can apply to your code *now*.

Item 1: Choose your containers with care.

You know that C++ puts a variety of containers at your disposal, but do you realize just how varied that variety is? To make sure you haven't overlooked any of your options, here's a quick review.

- The **standard STL sequence containers**, vector, string, deque, and list.

- The **standard STL associative containers**, set, multiset, map, and multimap.
- The **nonstandard sequence containers** slist and rope. slist is a singly linked list, and rope is essentially a heavy-duty string. (A “rope” is a heavy-duty “string.” Get it?) You’ll find a brief overview of these nonstandard (but commonly available) containers in [Item 50](#).
- The **nonstandard associative containers** hash_set, hash_multiset, hash_map, and hash_multimap. I examine these widely available hash-table-based variants on the standard associative containers in [Item 25](#).
- **vector<char> as a replacement for string.** [Item 13](#) describes the conditions under which such a replacement might make sense.
- **vector as a replacement for the standard associative containers.** As [Item 23](#) makes clear, there are times when vector can outperform the standard associative containers in both time and space.
- Several **standard non-STL containers**, including arrays, bitset, valarray, stack, queue, and priority_queue. Because these are non-STL containers, I have little to say about them in this book, though [Item 16](#) mentions a case where arrays are preferable to STL containers and [Item 18](#) explains why bitset may be better than vector<bool>. It’s also worth bearing in mind that arrays can be used with STL algorithms, because pointers can be used as array iterators.

That’s a panoply of options, and it’s matched in richness by the range of considerations that should go into choosing among them. Unfortunately, most discussions of the STL take a fairly narrow view of the world of containers, ignoring many issues relevant to selecting the one that is most appropriate. Even the Standard gets into this act, offering the following guidance for choosing among vector, deque, and list:

vector, list, and deque offer the programmer different complexity trade-offs and should be used accordingly. vector is the type of sequence that should be used by default. list should be used when there are frequent insertions and deletions from the middle of the sequence. deque is the data structure of choice when most insertions and deletions take place at the beginning or at the end of the sequence.

If your primary concern is algorithmic complexity, I suppose this constitutes reasonable advice, but there is so much more to be concerned with.

In a moment, we'll examine some of the important container-related issues that complement algorithmic complexity, but first I need to introduce a way of categorizing the STL containers that isn't discussed as often as it should be. That is the distinction between contiguous-memory containers and node-based containers.

Contiguous-memory containers (also known as *array-based containers*) store their elements in one or more (dynamically allocated) chunks of memory, each chunk holding more than one container element. If a new element is inserted or an existing element is erased, other elements in the same memory chunk have to be shifted up or down to make room for the new element or to fill the space formerly occupied by the erased element. This kind of movement affects both performance (see Items 5 and 14) and exception safety (as we'll soon see). The standard contiguous-memory containers are vector, string, and deque. The nonstandard rope is also a contiguous-memory container.

Node-based containers store only a single element per chunk of (dynamically allocated) memory. Insertion or erasure of a container element affects only pointers to nodes, not the contents of the nodes themselves, so element values need not be moved when something is inserted or erased. Containers representing linked lists, such as list and slist, are node-based, as are all the standard associative containers. (They're typically implemented as balanced trees.) The nonstandard hashed containers use varying node-based implementations, as you'll see in Item 25.

With this terminology out of the way, we're ready to sketch some of the questions most relevant when choosing among containers. In this discussion, I omit consideration of non-STL-like containers (e.g., arrays, bitsets, etc.), because this is, after all, a book on the STL.

- *Do you need to be able to insert a new element at an arbitrary position in the container?* If so, you need a sequence container; associative containers won't do.
- *Do you care how elements are ordered in the container?* If not, a hashed container becomes a viable choice. Otherwise, you'll want to avoid hashed containers.
- *Must the container be part of standard C++?* If so, that eliminates hashed containers, slist, and rope.
- *What category of iterators do you require?* If they must be random access iterators, you're technically limited to vector, deque, and string, but you'd probably want to consider rope, too. (See Item 50

for information on rope.) If bidirectional iterators are required, you must avoid `slist` (see [Item 50](#)) as well as one common implementation of the hashed containers (see [Item 25](#)).

- *Is it important to avoid movement of existing container elements when insertions or erasures take place?* If so, you'll need to stay away from contiguous-memory containers (see [Item 5](#)).
- *Does the data in the container need to be layout-compatible with C?* If so, you're limited to vectors (see [Item 16](#)).
- *Is lookup speed a critical consideration?* If so, you'll want to look at hashed containers (see [Item 25](#)), sorted vectors (see [Item 23](#)), and the standard associative containers — probably in that order.
- *Do you mind if the underlying container uses reference counting?* If so, you'll want to steer clear of `string`, because many `string` implementations are reference-counted (see [Item 13](#)). You'll need to avoid `rope`, too, because the definitive `rope` implementation is based on reference counting (see [Item 50](#)). You have to represent your strings somehow, of course, so you'll want to consider `vector<char>`.
- *Do you need transactional semantics for insertions and erasures?* That is, do you require the ability to reliably roll back insertions and erasures? If so, you'll want to use a node-based container. If you need transactional semantics for multiple-element insertions (e.g., the range form — see [Item 5](#)), you'll want to choose `list`, because `list` is the only standard container that offers transactional semantics for multiple-element insertions. Transactional semantics are particularly important for programmers interested in writing exception-safe code. (Transactional semantics can be achieved with contiguous-memory containers, too, but there is a performance cost, and the code is not as straightforward. To learn more about this, consult Item 17 of Sutter's *Exceptional C++* [8].)
- *Do you need to minimize iterator, pointer, and reference invalidation?* If so, you'll want to use node-based containers, because insertions and erasures on such containers never invalidate iterators, pointers, or references (unless they point to an element you are erasing). In general, insertions or erasures on contiguous-memory containers may invalidate all iterators, pointers, and references into the container.
- *Do you care if using `swap` on containers invalidates iterators, pointers, or references?* If so, you'll need to avoid `string`, because `string` is alone in the STL in invalidating iterators, pointers, and references during swaps.

- *Would it be helpful to have a sequence container with random access iterators where pointers and references to the data are not invalidated as long as nothing is erased and insertions take place only at the ends of the container?* This is a very special case, but if it's *your* case, deque is the container of your dreams. (Interestingly, deque's iterators *may* be invalidated when insertions are made only at the ends of the container. deque is the only standard STL container whose iterators may be invalidated without also invalidating its pointers and references.)

These questions are hardly the end of the matter. For example, they don't take into account the varying memory allocation strategies employed by the different container types. (Items 10 and 14 discuss some aspects of such strategies.) Still, they should be enough to convince you that, unless you have no interest in element ordering, standards conformance, iterator capabilities, layout compatibility with C, lookup speed, behavioral anomalies due to reference counting, the ease of implementing transactional semantics, or the conditions under which iterators are invalidated, you have more to think about than simply the algorithmic complexity of container operations. Such complexity is important, of course, but it's far from the entire story.

The STL gives you lots of options when it comes to containers. If you look beyond the bounds of the STL, there are even more options. Before choosing a container, be sure to consider *all* your options. A "default container"? I don't think so.

Item 2: Beware the illusion of container-independent code.

The STL is based on generalization. Arrays are generalized into containers and parameterized on the types of objects they contain. Functions are generalized into algorithms and parameterized on the types of iterators they use. Pointers are generalized into iterators and parameterized on the type of objects they point to.

That's just the beginning. Individual container types are generalized into sequence and associative containers, and similar containers are given similar functionality. Standard contiguous-memory containers (see Item 1) offer random-access iterators, while standard node-based containers (again, see Item 1) provide bidirectional iterators. Sequence containers support `push_front` and/or `push_back`, while associative containers don't. Associative containers offer logarithmic-time `lower_bound`, `upper_bound`, and `equal_range` member functions, but sequence containers don't.

With all this generalization going on, it's natural to want to join the movement. This sentiment is laudable, and when you write your own containers, iterators, and algorithms, you'll certainly want to pursue it. Alas, many programmers try to pursue it in a different manner. Instead of committing to particular types of containers in their software, they try to generalize the notion of a container so that they can use, say, a vector, but still preserve the option of replacing it with something like a deque or a list later — all without changing the code that uses it. That is, they strive to write *container-independent code*. This kind of generalization, well-intentioned though it is, is almost always misguided.

Even the most ardent advocate of container-independent code soon realizes that it makes little sense to try to write software that will work with both sequence and associative containers. Many member functions exist for only one category of container, e.g., only sequence containers support `push_front` or `push_back`, and only associative containers support `count` and `lower_bound`, etc. Even such basics as `insert` and `erase` have signatures and semantics that vary from category to category. For example, when you insert an object into a sequence container, it stays where you put it, but if you insert an object into an associative container, the container moves the object to where it belongs in the container's sort order. For another example, the form of `erase` taking an iterator returns a new iterator when invoked on a sequence container, but it returns nothing when invoked on an associative container. (Item 9 gives an example of how this can affect the code you write.)

Suppose, then, you aspire to write code that can be used with the most common sequence containers: vector, deque, and list. Clearly, you must program to the intersection of their capabilities, and that means no uses of `reserve` or `capacity` (see Item 14), because deque and list don't offer them. The presence of list also means you give up `operator[]`, and you limit yourself to the capabilities of bidirectional iterators. That, in turn, means you must stay away from algorithms that demand random access iterators, including `sort`, `stable_sort`, `partial_sort`, and `nth_element` (see Item 31).

On the other hand, your desire to support vector rules out use of `push_front` and `pop_front`, and both vector and deque put the kibosh on `splice` and the member form of `sort`. In conjunction with the constraints above, this latter prohibition means that there is no form of `sort` you can call on your “generalized sequence container.”

That's the obvious stuff. If you violate any of those restrictions, your code will fail to compile with at least one of the containers you want to be able to use. The code that *will* compile is more insidious.

The main culprit is the different rules for invalidation of iterators, pointers, and references that apply to different sequence containers. To write code that will work correctly with `vector`, `deque`, and `list`, you must assume that any operation invalidating iterators, pointers, or references in any of those containers invalidates them in the container you're using. Thus, you must assume that every call to `insert` invalidates everything, because `deque::insert` invalidates all iterators and, lacking the ability to call `capacity`, `vector::insert` must be assumed to invalidate all pointers and references. (Item 1 explains that `deque` is unique in sometimes invalidating its iterators without invalidating its pointers and references.) Similar reasoning leads to the conclusion that, unless you're erasing the last element of a container, calls to `erase` must also be assumed to invalidate everything.

Want more? You can't pass the data in the container to a C interface, because only `vector` supports that (see Item 16). You can't instantiate your container with `bool` as the type of objects to be stored, because, as Item 18 explains, `vector<bool>` doesn't always behave like a vector, and it never actually stores bools. You can't assume `list`'s constant-time insertions and erasures, because `vector` and `deque` take linear time to perform those operations.

When all is said and done, you're left with a "generalized sequence container" where you can't call `reserve`, `capacity`, `operator[]`, `push_front`, `pop_front`, `splice`, or any algorithm requiring random access iterators; a container where every call to `insert` and `erase` takes linear time and invalidates all iterators, pointers, and references; and a container incompatible with C where bools can't be stored. Is that really the kind of container you want to use in your applications? I suspect not.

If you rein in your ambition and decide you're willing to drop support for `list`, you still give up `reserve`, `capacity`, `push_front`, and `pop_front`; you still must assume that all calls to `insert` and `erase` take linear time and invalidate everything; you still lose layout compatibility with C; and you still can't store bools.

If you abandon the sequence containers and shoot instead for code that can work with different associative containers, the situation isn't much better. Writing for both `set` and `map` is close to impossible, because sets store single objects while maps store pairs of objects. Even writing for both `set` and `multiset` (or `map` and `multimap`) is tough. The `insert` member function taking only a value has different return types for sets/maps than for their multi cousins, and you must reli-

giously avoid making any assumptions about how many copies of a value are stored in a container. With `map` and `multimap`, you must avoid using `operator[]`, because that member function exists only for `map`.

Face the truth: it's not worth it. The different containers are *different*, and they have strengths and weaknesses that vary in significant ways. They're not designed to be interchangeable, and there's little you can do to paper that over. If you try, you're merely tempting fate, and fate doesn't like to be tempted.

Still, the day will dawn when you'll realize that a container choice you made was, er, suboptimal, and you'll need to use a different container type. You now know that when you change container types, you'll not only need to fix whatever problems your compilers diagnose, you'll also need to examine all the code using the container to see what needs to be changed in light of the new container's performance characteristics and rules for invalidation of iterators, pointers, and references. If you switch from a `vector` to something else, you'll also have to make sure you're no longer relying on `vector`'s C-compatible memory layout, and if you switch to a `vector`, you'll have to ensure that you're not using it to store `bool`s.

Given the inevitability of having to change container types from time to time, you can facilitate such changes in the usual manner: by encapsulating, encapsulating, encapsulating. One of the easiest ways to do this is through the liberal use of typedefs for container types. Hence, instead of writing this,

```
class Widget { ... };
vector<Widget> vw;
Widget bestWidget;
...
vector<Widget>::iterator i = find(vw.begin(), vw.end(), bestWidget);
```

write this:

```
class Widget { ... };
typedef vector<Widget> WidgetContainer;
WidgetContainer cw;
Widget bestWidget;
...
WidgetContainer::iterator i = find(cw.begin(), cw.end(), bestWidget);
```

This makes it a lot easier to change container types, something that's especially convenient if the change in question is simply to add a custom allocator. (Such a change doesn't affect the rules for iterator/pointer/reference invalidation.)

```
class Widget { ... };
template<typename T>                               // see Item 10 for why this
SpecialAllocator { ... };                          // needs to be a template
typedef vector<Widget, SpecialAllocator<Widget> > WidgetContainer;
WidgetContainer cw;                                // still works
Widget bestWidget;
...
WidgetContainer::iterator i =
    find(cw.begin(), cw.end(), bestWidget);        // still works
```

If the encapsulating aspects of typedefs mean nothing to you, you're still likely to appreciate the work they can save, especially for iterator types. For example, if you have an object of type

```
map<string,
    vector<Widget>::iterator,
    CStringCompare>
// CStringCompare is "case-
// insensitive string compare;"
// Item 19 describes it
```

and you want to walk through the map using `const_iterator`s, do you really want to spell out

```
map<string, vector<Widget>::iterator, CStringCompare>::const_iterator
```

more than once? Once you've used the STL a little while, you'll realize that typedefs are your friends.

A typedef is just a synonym for some other type, so the encapsulation it affords is purely lexical. A typedef doesn't prevent a client from doing (or depending on) anything they couldn't already do (or depend on). You need bigger ammunition if you want to limit client exposure to the container choices you've made. You need classes.

To limit the code that may require modification if you replace one container type with another, hide the container in a class, and limit the amount of container-specific information visible through the class interface. For example, if you need to create a customer list, don't use a list directly. Instead, create a `CustomerList` class, and hide a list in its private section:

```

class CustomerList {
private:
    typedef list<Customer> CustomerContainer;
    typedef CustomerContainer::iterator CCLiterator;

    CustomerContainer customers;

public:
    ...                               // limit the amount of list-specific
};                                   // information visible through
                                   // this interface

```

At first, this may seem silly. After all a customer list is a *list*, right? Well, maybe. Later you may discover that you don't need to insert or erase customers from the middle of the list as often as you'd anticipated, but you do need to quickly identify the top 20% of your customers — a task tailor-made for the `nth_element` algorithm (see [Item 31](#)). But `nth_element` requires random access iterators. It won't work with a list. In that case, your customer “list” might be better implemented as a vector or a deque.

When you consider this kind of change, you still have to check every `CustomerList` member function and every friend to see how they'll be affected (in terms of performance and iterator/pointer/reference invalidation, etc.), but if you've done a good job of encapsulating `CustomerList`'s implementation details, the impact on `CustomerList` clients should be small. You can't write container-independent code, but *they* might be able to.

Item 3: Make copying cheap and correct for objects in containers.

Containers hold objects, but not the ones you give them. Instead, when you add an object to a container (via, e.g., `insert` or `push_back`, etc.), what goes into the container is a *copy* of the object you specify.

Once an object is in a container, it's not uncommon for it to be copied further. If you insert something into or erase something from a vector, string, or deque, existing container elements are typically moved (copied) around (see [Items 5](#) and [14](#)). If you use any of the sorting algorithms (see [Item 31](#)); `next_permutation` or `previous_permutation`; `remove`, `unique`, or their ilk (see [Item 32](#)); `rotate` or `reverse`, etc., objects will be moved (copied) around. Yes, copying objects is the STL way.

It may interest you to know how all this copying is accomplished. That's easy. An object is copied by using its copying member functions, in particular, its *copy* constructor and its *copy* assignment operator. (Clever names, no?) For a user-defined class like `Widget`, these functions are traditionally declared like this:

```
class Widget {
public:
    ...
    Widget(const Widget&);           // copy constructor
    Widget& operator=(const Widget&); // copy assignment operator
    ...
};
```

As always, if you don't declare these functions yourself, your compilers will declare them for you. Also as always, the copying of built-in types (e.g., ints, pointers, etc.) is accomplished by simply copying the underlying bits. (For details on copy constructors and assignment operators, consult any introductory book on C++. In *Effective C++*, Items 11 and 27 focus on the behavior of these functions.)

With all this copying taking place, the motivation for this Item should now be clear. If you fill a container with objects where copying is expensive, the simple act of putting the objects into the container could prove to be a performance bottleneck. The more things get moved around in the container, the more memory and cycles you'll blow on making copies. Furthermore, if you have objects where "copying" has an unconventional meaning, putting such objects into a container will invariably lead to grief. (For an example of the kind of grief it can lead to, see [Item 8](#).)

In the presence of inheritance, of course, copying leads to slicing. That is, if you create a container of base class objects and you try to insert derived class objects into it, the derivedness of the objects will be removed as the objects are copied (via the base class copy constructor) into the container:

```
vector<Widget> vw;
class SpecialWidget:           // SpecialWidget inherits from
    public Widget { ... };      // Widget above
SpecialWidget sw;
vw.push_back(sw);              // sw is copied as a base class
                                // object into vw. Its specialness
                                // is lost during the copying
```

The slicing problem suggests that inserting a derived class object into a container of base class objects is almost always an error. If you want

the resulting object to *act* like a derived class object, e.g., invoke derived class virtual functions, etc., it is always an error. (For more background on the slicing problem, consult *Effective C++*, Item 22. For another example of where it arises in the STL, see [Item 38](#).)

An easy way to make copying efficient, correct, and immune to the slicing problem is to create containers of *pointers* instead of containers of objects. That is, instead of creating a container of `Widget`, create a container of `Widget*`. Copying pointers is fast, it always does exactly what you expect (it copies the bits making up the pointer), and nothing gets sliced when a pointer is copied. Unfortunately, containers of pointers have their own STL-related headaches. You can read about them in [Items 7](#) and [33](#). As you seek to avoid those headaches while still dodging efficiency, correctness, and slicing concerns, you'll probably discover that containers of *smart pointers* are an attractive option. To learn more about this option, turn to [Item 7](#).

If all this makes it sound like the STL is copy-crazy, think again. Yes, the STL makes lots of copies, but it's generally designed to avoid copying objects *unnecessarily*. In fact, it's generally designed to avoid *creating* objects unnecessarily. Contrast this with the behavior of C's and C++'s only built-in container, the lowly array:

```
Widget w[maxNumWidgets]; // create an array of maxNumWidgets
                          // Widgets, default-constructing each one
```

This constructs `maxNumWidgets` `Widget` objects, even if we normally expect to use only a few of them or we expect to immediately overwrite each default-constructed value with values we get from someplace else (e.g., a file). Using the STL instead of an array, we can use a vector that grows when it needs to:

```
vector<Widget> vw;           // create a vector with zero Widget
                             // objects that will expand as needed
```

We can also create an empty vector that contains enough space for `maxNumWidgets` `Widgets`, but where zero `Widgets` have been constructed:

```
vector<Widget> vw;
vw.reserve(maxNumWidgets); // see Item 14 for details on reserve
```

Compared to arrays, STL containers are much more civilized. They create (by copying) only as many objects as you ask for, they do it only when you direct them to, and they use a default constructor only when you say they should. Yes, STL containers make copies, and yes, you need to understand that, but don't lose sight of the fact that they're still a big step up from arrays.

Item 4: Call `empty` instead of checking `size()` against zero.

For any container `c`, writing

```
if (c.size() == 0) ...
```

is essentially equivalent to writing

```
if (c.empty()) ...
```

That being the case, you might wonder why one construct should be preferred to the other, especially in view of the fact that `empty` is typically implemented as an inline function that simply returns whether `size` returns 0.

You should prefer the construct using `empty`, and the reason is simple: `empty` is a constant-time operation for all standard containers, but for some list implementations, `size` may take linear time.

But what makes list so troublesome? Why can't it, too, offer a constant-time `size`? The answer has much to do with the range form of list's unique splicing functions. Consider this code:

```
list<int> list1;
list<int> list2;

...
list1.splice(                                // move all nodes in list2
    list1.end(), list2,                      // from the first occurrence
    find(list2.begin(), list2.end(), 5),     // of 5 through the last
    find(list2.rbegin(), list2.rend(), 10).base() // occurrence of 10 to the
);                                           // end of list1. See Item 28
                                           // for info on the "base()" call
```

This code won't work unless `list2` contains a 10 somewhere beyond a 5, but let's assume that's not a problem. Instead, let's focus on this question: how many elements are in `list1` after the splice? Clearly, `list1` after the splice has as many elements as it did before the splice plus however many elements were spliced into it. But how many elements were spliced into it? As many as were in the range defined by `find(list2.begin(), list2.end(), 5)` and `find(list2.rbegin(), list2.rend(), 10).base()`. Okay, how many is that? Without traversing the range and counting them, there's no way to know. And therein lies the problem.

Suppose you're responsible for implementing `list`. `list` isn't just any container, it's a *standard* container, so you know your class will be widely used. You naturally want your implementation to be as efficient as possible. You figure that clients will commonly want to find out how many elements are in a list, so you'd like to make `size` a constant-

time operation. You'd thus like to design `list` so it always knows how many elements it contains.

At the same time, you know that of all the standard containers, only `list` offers the ability to splice elements from one place to another without copying any data. You reason that many list clients will choose `list` specifically because it offers high-efficiency splicing. They know that splicing a range from one list to another can be accomplished in constant time, and you know that they know it, so you certainly want to meet their expectation that `splice` is a constant-time member function.

This puts you in a quandary. If `size` is to be a constant-time operation, each list member function must update the sizes of the lists on which it operates. That includes `splice`. But the only way for the range version of `splice` to update the sizes of the lists it modifies is for it to count the number of elements being spliced, and doing that would prevent it from achieving the constant-time performance you want for *it*. If you eliminate the requirement that the range form of `splice` update the sizes of the lists it's modifying, `splice` can be made constant-time, but then `size` becomes a linear-time operation. In general, it will have to traverse its entire data structure to see how many elements it contains. No matter how you look at it, something — `size` or the range form of `splice` — has to give. One or the other can be a constant-time operation, but not both.

Different list implementations resolve this conflict in different ways, depending on whether their authors choose to maximize the efficiency of `size` or the range form of `splice`. If you happen to be using a list implementation where a constant-time range form of `splice` was given higher priority than a constant-time `size`, you'll be better off calling `empty` than `size`, because `empty` is always a constant-time operation. Even if you're not using such an implementation, you might find yourself using such an implementation in the future. For example, you might port your code to a different platform where a different implementation of the STL is available, or you might just decide to switch to a different STL implementation for your current platform.

No matter what happens, you can't go wrong if you call `empty` instead of checking to see if `size() == 0`. So call `empty` whenever you need to know whether a container has zero elements.

Item 5: Prefer range member functions to their single-element counterparts.

Quick! Given two vectors, `v1` and `v2`, what's the easiest way to make `v1`'s contents be the same as the second half of `v2`'s? Don't agonize

over the definition of “half” when `v2` has an odd number of elements, just do something reasonable.

Time’s up! If your answer was

```
v1.assign(v2.begin() + v2.size() / 2, v2.end());
```

or something quite similar, you get full credit and a gold star. If your answer involved more than one statement, but didn’t use any kind of loop, you get nearly full credit, but no gold star. If your answer involved a loop, you’ve got some room for improvement, and if your answer involved multiple loops, well, let’s just say that you really need this book.

By the way, if your response to the *answer* to the question included “Huh?”, pay close attention, because you’re going to learn something really useful.

This quiz is designed to do two things. First, it affords me an opportunity to remind you of the existence of the `assign` member function, a convenient beast that too many programmers overlook. It’s available for all the standard sequence containers (vector, string, deque, and list). Whenever you have to completely replace the contents of a container, you should think of assignment. If you’re just copying one container to another of the same type, `operator=` is the assignment function of choice, but as this example demonstrates, `assign` is available for the times when you want to give a container a completely new set of values, but `operator=` won’t do what you want.

The second reason for the quiz is to demonstrate why range member functions are superior to their single-element alternatives. A *range member function* is a member function that, like STL algorithms, uses two iterator parameters to specify a range of elements over which something should be done. Without using a range member function to solve this Item’s opening problem, you’d have to write an explicit loop, probably something like this:

```
vector<Widget> v1, v2;           // assume v1 and v2 are vectors
...                               // of Widgets
v1.clear();
for (vector<Widget>::const_iterator ci = v2.begin() + v2.size() / 2;
     ci != v2.end();
     ++ci)
    v1.push_back(*ci);
```

Item 43 examines in detail why you should try to avoid writing explicit loops, but you don’t need to read that Item to recognize that writing this code is a lot more work than is writing the call to `assign`. As we’ll

see shortly, the loop also happens to impose an efficiency penalty, but we'll deal with that in a moment.

One way to avoid the loop is to follow the advice of [Item 43](#) and employ an algorithm instead:

```
v1.clear();  
copy(v2.begin() + v2.size() / 2, v2.end(), back_inserter(v1));
```

Writing this is still more work than writing the call to assign. Furthermore, though no loop is present in this code, one certainly exists inside `copy` (see [Item 43](#)). As a result, the efficiency penalty remains. Again, I'll discuss that below. At this point, I want to digress briefly to observe that almost all uses of `copy` where the destination range is specified using an insert iterator (i.e., via `inserter`, `back_inserter`, or `front_inserter`) can be — *should* be — replaced with calls to range member functions. Here, for example, the call to `copy` can be replaced with a range version of `insert`:

```
v1.insert(v1.end(), v2.begin() + v2.size() / 2, v2.end());
```

This involves slightly less typing than the call to `copy`, but it also says more directly what is happening: data is being inserted into `v1`. The call to `copy` expresses that, too, but less directly. It puts the emphasis in the wrong place. The interesting aspect of what is happening is not that elements are being copied, it's that `v1` is having new data added to it. The `insert` member function makes that clear. The use of `copy` obscures it. There's nothing interesting about the fact that things are being copied, because the STL is *built* on the assumption that things will be copied. Copying is so fundamental to the STL, it's the topic of [Item 3](#) in this book!

Too many STL programmers overuse `copy`, so the advice I just gave bears repeating: Almost all uses of `copy` where the destination range is specified using an insert iterator should be replaced with calls to range member functions.

Returning to our assign example, we've already identified two reasons to prefer range member functions to their single-element counterparts:

- It's generally less work to write the code using the range member functions.
- Range member functions tend to lead to code that is clearer and more straightforward.

In short, range member functions yield code that is easier to write and easier to understand. What's not to like?

Alas, some will dismiss these arguments as matters of programming style, and developers enjoy arguing about style issues almost as much as they enjoy arguing about which is the One True Editor. (As if there's any doubt. It's Emacs.) It would be helpful to have a more universally agreed-upon criterion for establishing the superiority of range member functions to their single-element counterparts. For the standard sequence containers, we have one: efficiency. When dealing with the standard sequence containers, application of single-element member functions makes more demands on memory allocators, copies objects more frequently, and/or performs redundant operations compared to range member functions that achieve the same end.

For example, suppose you'd like to copy an array of ints into the front of a vector. (The data might be in an array instead of a vector in the first place, because the data came from a legacy C API. For a discussion of the issues that arise when mixing STL containers and C APIs, see [Item 16](#).) Using the vector range insert function, it's honestly trivial:

```
int data[numValues];           // assume numValues is
                                // defined elsewhere
vector<int> v;
...
v.insert(v.begin(), data, data + numValues); // insert the ints in data
                                              // into v at the front
```

Using iterative calls to insert in an explicit loop, it would probably look more or less like this:

```
vector<int>::iterator insertLoc(v.begin());
for (int i = 0; i < numValues; ++i) {
    insertLoc = v.insert(insertLoc, data[i]);
    ++insertLoc;
}
```

Notice how we have to be careful to save the return value of insert for the next loop iteration, then increment the returned iterator. If we didn't update insertLoc after each insertion, we'd have two problems. First, all loop iterations after the first would yield undefined behavior, because each insert call would invalidate insertLoc. Second, even if insertLoc remained valid, we'd always insert at the front of the vector (i.e., at v.begin()), and the result would be that the ints copied into v would end up in reverse order.

If we follow the lead of [Item 43](#) and replace the loop with a call to copy, we get something like this:

```
copy(data, data + numValues, inserter(v, v.begin()));
```

By the time the copy template has been instantiated, the code based on copy and the code using the explicit loop will be almost identical, so for purposes of an efficiency analysis, we'll focus on the explicit loop, keeping in mind that the analysis is equally valid for the code employing copy. Looking at the explicit loop just makes it easier to understand where the efficiency hits come from. Yes, that's "hits," plural, because the code using the single-element version of insert levies up to three different performance taxes on you, none of which you pay if you use the range version of insert.

The first tax consists of unnecessary function calls. Inserting `numValues` elements into `v` one at a time naturally costs you `numValues` calls to `insert`. Using the range form of `insert`, you pay for only one function call, a savings of `numValues-1` calls. Of course, it's possible that inlining will save you from this tax, but then again, it's possible that it won't. Only one thing is sure. With the range form of `insert`, you definitely won't pay it.

Inlining won't save you from the second tax, which is the cost of inefficiently moving the existing elements in `v` to their final post-insertion positions. Each time `insert` is called to add a new value to `v`, every element above the insertion point must be moved up one position to make room for the new element. So the element at position p must be moved up to position $p+1$, etc. In our example, we're inserting `numValues` elements at the front of `v`. That means that each element in `v` prior to the insertions will have to be shifted up a total of `numValues` positions. But each will be shifted up only one position each time `insert` is called, so each element will be moved a total of `numValues` times. If `v` has n elements prior to the insertions, a total of $n*\text{numValues}$ moves will take place. In this example, `v` holds ints, so each move will probably boil down to an invocation of `memmove`, but if `v` held a user-defined type like `Widget`, each move would incur a call to that type's assignment operator or copy constructor. (Most calls would be to the assignment operator, but each time the last element in the vector was moved, that move would be accomplished by calling the element's copy constructor.) In the general case, then, inserting `numValues` new objects one at a time into the front of a `vector<Widget>` holding n elements exacts a cost of $n*\text{numValues}$ function calls: $(n-1)*\text{numValues}$ calls to the `Widget` assignment operator and `numValues` calls to the `Widget` copy constructor. Even if these calls are inlined, you're still doing the work to move the elements in `v` `numValues` times.

In contrast, the Standard requires that range `insert` functions move existing container elements directly into their final positions, i.e., at a cost of one move per element. The total cost is n moves, `numValues` to

the copy constructor for the type of objects in the container, the remainder to that type's assignment operator. Compared to the single-element insert strategy, the range insert performs $n*(\text{numValues}-1)$ fewer moves. Think about that for a minute. It means that if numValues is 100, the range form of insert would do 99% fewer moves than the code making repeated calls to the single-element form of insert!

Before I move on to the third efficiency cost of single-element member functions vis-à-vis their range counterparts, I have a minor correction. What I wrote in the previous paragraph is the truth and nothing but the truth, but it's not quite the whole truth. A range insert function can move an element into its final position in a single move only if it can determine the distance between two iterators without losing its place. This is almost always possible, because all forward iterators offer this functionality, and forward iterators are nearly ubiquitous. All iterators for the standard containers offer forward iterator functionality. So do the iterators for the nonstandard hashed containers (see [Item 25](#)). Pointers acting as iterators into arrays offer such functionality, too. In fact, the only standard iterators that don't offer forward iterator capabilities are input and output iterators. Thus, everything I wrote above is true except when the iterators passed to the range form of insert are input iterators (e.g. `istream_iterators` — see [Item 6](#)). In that case only, range insert is allowed to move elements into their final positions one place at a time, and if it's implemented that way, its advantage as regards number of element movements ceases to exist. (For output iterators, this issue fails to arise, because output iterators can't be used to specify a range for insert.)

The final performance tax levied on those so foolish as to use repeated single-element insertions instead of a single range insertion has to do with memory allocation, though it has a nasty copying side to it, too. As [Item 14](#) explains, when you try to insert an element into a vector whose memory is full, the vector allocates new memory with more capacity, copies its elements from the old memory to the new memory, destroys the elements in the old memory, and deallocates the old memory. Then it adds the element that is being inserted. [Item 14](#) also explains that vector implementations increase their capacity by some multiplicative factor each time they run out of memory, so inserting numValues new elements often results in new memory being allocated a number of times proportional to the logarithm of numValues. Inserting 1000 elements one at a time into a vector with a growth rate of 1.5 could thus result in 18 new allocations[†] (including their incumbent copying of elements). In contrast (and, by now, predictably), a range insertion can figure out how much new memory it needs before it starts inserting things (assuming it is given forward iterators), so it

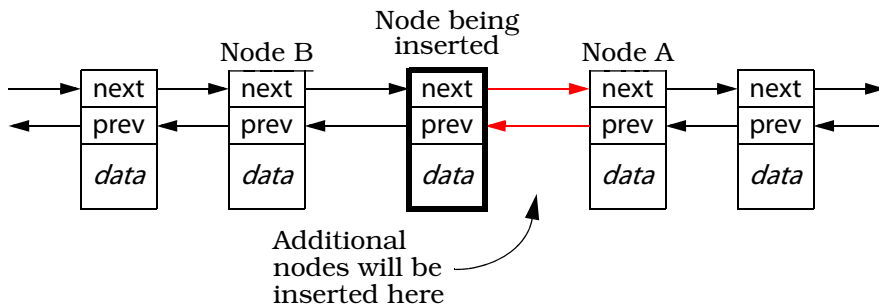
[†] Because $\log_{1.5} 1000 = 18$ (rounding up to an integer).

need not reallocate a vector's underlying memory more than once. As you can imagine, the savings can be considerable.

The analysis I've just performed is for vectors, but the same reasoning applies to strings, too. For deques, the reasoning is similar, but deques manage their memory differently from vectors and strings, so the argument about repeated memory allocations doesn't apply. The argument about moving elements an unnecessarily large number of times, however, generally does apply (though the details are different), as does the observation about the number of function calls.

Among the standard sequence containers, that leaves only list, but here, too, there is a performance advantage to using a range form of insert instead of a single-element form. The argument about repeated function calls continues to be valid, of course, but, because of the way linked lists work, the copying and memory allocation issues fail to arise. Instead, there is a new problem: repeated superfluous assignments to the next and prev pointers of some nodes in the list.

Each time an element is added to a linked list, the list node holding that element must have its next and prev pointers set, and of course the node preceding the new node (let's call it B, for "before") must set its next pointer and the node following the new node (we'll call it A, for "after") must set its prev pointer:



When a series of new nodes is added one by one by calling list's single-element insert, all but the last new node will set its next pointer *twice*, once to point to A, a second time to point to the element inserted after it. A will set its prev pointer to point to a new node each time one is inserted in front of it. If numValues nodes are inserted in front of A, numValues-1 superfluous assignments will be made to the inserted nodes' next pointers, and numValues-1 superfluous assignments will be made to A's prev pointer. All told, that's $2 * (\text{numValues} - 1)$ unnecessary pointer assignments. Pointer assignments are cheap, of course, but why pay for them if you don't have to?

By now it should be clear that you don't have to, and the key to evading the cost is to use list's range form of insert. Because that function knows how many nodes will ultimately be inserted, it can avoid the superfluous pointer assignments, using only a single assignment to each pointer to set it to its proper post-insertion value.

For the standard sequence containers, then, a lot more than programming style is on the line when choosing between single-element insertions and range insertions. For the associative containers, the efficiency case is harder to make, though the issue of extra function call overhead for repeated calls to single-element insert continues to apply. Furthermore, certain special kinds of range insertions may lead to optimization possibilities in associative containers, too, but as far as I can tell, such optimizations currently exist only in theory. By the time you read this, of course, theory may have become practice, so range insertions into associative containers may indeed be more efficient than repeated single-element insertions. Certainly they are never less efficient, so you have nothing to lose by preferring them.

Even without the efficiency argument, the fact remains that using range member functions requires less typing as you write the code, and it also yields code that is easier to understand, thus enhancing your software's long-term maintainability. Those two characteristics alone should convince you to prefer range member functions. The efficiency edge is really just a bonus.

Having droned on this long about the wonder of range member functions, it seems only appropriate that I summarize them for you. Knowing which member functions support ranges makes it a lot easier to recognize opportunities to use them. In the signatures below, the parameter type iterator literally means the iterator type for the container, i.e., `container::iterator`. The parameter type `InputIterator`, on the other hand, means that any input iterator is acceptable.

- **Range construction.** All standard containers offer a constructor of this form:

```
container::container(InputIterator begin,    // beginning of range  
                     InputIterator end);    // end of range
```

When the iterators passed to this constructor are `istream_iterators` or `istreambuf_iterators` (see [Item 29](#)), you may encounter C++'s most astonishing parse, one that causes your compilers to interpret this construct as a function declaration instead of as the definition of a new container object. [Item 6](#) tells you everything you need to know about that parse, including how to defeat it.

- **Range insertion.** All standard sequence containers offer this form of insert:

```
void container::insert(iterator position,    // where to insert the range
                      InputIterator begin,  // start of range to insert
                      InputIterator end);    // end of range to insert
```

Associative containers use their comparison function to determine where elements go, so they offer a signature that omits the position parameter:

```
void container::insert(InputIterator begin, InputIterator end);
```

When looking for ways to replace single-element inserts with range versions, don't forget that some single-element variants camouflage themselves by adopting different function names. For example, `push_front` and `push_back` both insert single elements into containers, even though they're not called `insert`. If you see a loop calling `push_front` or `push_back`, or if you see an algorithm such as `copy` being passed `front_inserter` or `back_inserter` as a parameter, you've discovered a place where a range form of `insert` is likely to be a superior strategy.

- **Range erasure.** Every standard container offers a range form of `erase`, but the return types differ for sequence and associative containers. Sequence containers provide this,

```
iterator container::erase(iterator begin, iterator end);
```

while associative containers offer this:

```
void container::erase(iterator begin, iterator end);
```

Why the difference? The claim is that having the associative container version of `erase` return an iterator (to the element following the one that was erased) would incur an unacceptable performance penalty. I'm one of many who find this claim specious, but the Standard says what the Standard says, and what the Standard says is that sequence and associative container versions of `erase` have different return types.

Most of this Item's efficiency analysis for `insert` has analogues for `erase`. The number of function calls is still greater for repeated calls to single-element `erase` than for a single call to range `erase`. Element values must still be shifted one position at a time towards their final destination when using single-element `erase`, while range `erase` can move them into their final positions in a single move.

One argument about `vector`'s and `string`'s `insert` that fails to apply to `erase` has to do with repeated allocations. (For `erase`, of course, it

would concern repeated *deallocations*.) That's because the memory for vectors and strings automatically grows to accommodate new elements, but it doesn't automatically shrink when the number of elements is reduced. (Item 17 describes how you may reduce the unnecessary memory held by a vector or string.)

One particularly important manifestation of range erase is the erase-remove idiom. You can read all about it in Item 32.

- **Range assignment.** As I noted at the beginning of this Item, all standard sequence containers offer a range form of assign:

```
void container::assign(InputIterator begin, InputIterator end);
```

So there you have it, three solid arguments for preferring range member functions to their single-element counterparts. Range member functions are easier to write, they express your intent more clearly, and they exhibit higher performance. That's a troika that's hard to beat.

Item 6: Be alert for C++'s most vexing parse.

Suppose you have a file of ints and you'd like to copy those ints into a list. This seems like a reasonable way to do it:

```
ifstream dataFile("ints.dat");
list<int> data(istream_iterator<int>(dataFile), // warning! this doesn't do
             istream_iterator<int>());         // what you think it does
```

The idea here is to pass a pair of `istream_iterator`s to `list`'s range constructor (see Item 5), thus copying the ints in the file into the list.

This code will compile, but at runtime, it won't do anything. It won't read any data out of a file. It won't even create a list. That's because the second statement doesn't declare a list and it doesn't call a constructor. What it does is ... well, what it does is so strange, I dare not tell you straight out, because you won't believe me. Instead, I have to develop the explanation, bit by bit. Are you sitting down? If not, you might want to look around for a chair...

We'll start with the basics. This line declares a function `f` taking a double and returning an int:

```
int f(double d);
```

This next line does the same thing. The parentheses around the parameter name `d` are superfluous and are ignored:

```
int f(double (d));           // same as above; parens around d are ignored
```

The line below declares the same function. It simply omits the parameter name:

```
int f(double);           // same as above; parameter name is omitted
```

Those three declaration forms should be familiar to you, though the ability to put parentheses around a parameter name may have been new. (It wasn't long ago that it was new to me.)

Let's now look at three more function declarations. The first one declares a function *g* taking a parameter that's a pointer to a function taking nothing and returning a double:

```
int g(double (*pf)());   // g takes a pointer to a function as a parameter
```

Here's another way to say the same thing. The only difference is that *pf* is declared using non-pointer syntax (a syntax that's valid in both C and C++):

```
int g(double pf());      // same as above; pf is implicitly a pointer
```

As usual, parameter names may be omitted, so here's a third declaration for *g*, one where the name *pf* has been eliminated:

```
int g(double ());        // same as above; parameter name is omitted
```

Notice the difference between parentheses *around a parameter name* (such as *d* in the second declaration for *f*) and *standing by themselves* (as in this example). Parentheses around a parameter name are ignored, but parentheses standing by themselves indicate the existence of a parameter list; they announce the presence of a parameter that is itself a pointer to a function.

Having warmed ourselves up with these declarations for *f* and *g*, we are ready to examine the code that began this Item. Here it is again:

```
list<int> data(istream_iterator<int>(dataFile),
              istream_iterator<int>());
```

Brace yourself. This declares a *function*, *data*, whose return type is *list<int>*. The function *data* takes two parameters:

- The first parameter is named *dataFile*. Its type is *istream_iterator<int>*. The parentheses around *dataFile* are superfluous and are ignored.
- The second parameter has no name. Its type is pointer to function taking nothing and returning an *istream_iterator<int>*.

Amazing, huh? But it's consistent with a universal rule in C++, which says that pretty much anything that can be parsed as a function declaration will be. If you've been programming in C++ for a while, you've

almost certainly encountered another manifestation of this rule. How many times have you seen this mistake?

```
class Widget { ... };           // assume Widget has a default constructor
Widget w();                     // uh oh...
```

This doesn't declare a `Widget` named `w`, it declares a function named `w` that takes nothing and returns a `Widget`. Learning to recognize this *faux pas* is a veritable rite of passage for C++ programmers.

All of which is interesting (in its own twisted way), but it doesn't help us say what we want to say, which is that a `list<int>` object should be initialized with the contents of a file. Now that we know what parse we have to defeat, that's easy to express. It's not legal to surround a formal parameter declaration with parentheses, but it is legal to surround an argument to a function call with parentheses, so by adding a pair of parentheses, we force compilers to see things our way:

```
list<int> data((istream_iterator<int>(dataFile)), // note new parens
              istream_iterator<int>());           // around first argument
                                                    // to list's constructor
```

This is the proper way to declare `data`, and given the utility of `istream_iterators` and range constructors (again, see [Item 5](#)), it's worth knowing how to do it.

Unfortunately, not all compilers currently know it themselves. Of the several I tested, almost half refused to accept `data`'s declaration unless it was *incorrectly* declared without the additional parentheses! To placate such compilers, you could roll your eyes and use the declaration for `data` that I've painstakingly explained is incorrect, but that would be both unportable and short-sighted. After all, compilers that currently get the parse wrong will surely correct it in the future, right? (Surely!)

A better solution is to step back from the trendy use of anonymous `istream_iterator` objects in `data`'s declaration and simply give those iterators names. The following code should work everywhere:

```
ifstream dataFile("ints.dat");
istream_iterator<int> dataBegin(dataFile);
istream_iterator<int> dataEnd;
list<int> data(dataBegin, dataEnd);
```

This use of named iterator objects runs contrary to common STL programming style, but you may decide that's a price worth paying for code that's unambiguous to both compilers and the humans who have to work with them.

Item 7: When using containers of newed pointers, remember to delete the pointers before the container is destroyed.

Containers in the STL are remarkably smart. They serve up iterators for both forward and reverse traversals (via `begin`, `end`, `rbegin`, etc.); they tell you what type of objects they contain (via their `value_type` typedef); during insertions and erasures, they take care of any necessary memory management; they report both how many objects they hold and the most they may contain (via `size` and `max_size`, respectively); and of course they automatically destroy each object they hold when they (the containers) are themselves destroyed.

Given such brainy containers, many programmers stop worrying about cleaning up after themselves. Heck, they figure, their containers will do the worrying for them. In many cases, they're right, but when the containers hold *pointers* to objects allocated with `new`, they're not right enough. Sure, a container of pointers will destroy each element it contains when it (the container) is destroyed, but the "destructor" for a pointer is a no-op! It certainly doesn't call `delete`.

As a result, the following code leads straight to a resource leak:

```
void doSomething()
{
    vector<Widget*> vwp;
    for (int i = 0; i < SOME_MAGIC_NUMBER; ++i)
        vwp.push_back(new Widget);

    ...                               // use vwp
}                                     // Widgets are leaked here!
```

Each of `vwp`'s elements is destroyed when `vwp` goes out of scope, but that doesn't change the fact that `delete` was never used for the objects conjured up with `new`. Such deletion is your responsibility, not that of your vector. This is a feature. Only you know whether the pointers *should* be deleted.

Usually, you want them to be. When that's the case, making it happen seems easy enough:

```
void doSomething()
{
    vector<Widget*> vwp;

    ...                               // as before
```

```

    for( vector<Widget*>::iterator i = vwp.begin();
        i != vwp.end();
        ++i)
        delete *i;
}

```

This works, but only if you're not terribly picky about what you mean by "works". One problem is that the new for loop does pretty much what `for_each` does, but it's not as clear as using `for_each` (see [Item 43](#)). Another is that the code isn't exception safe. If an exception is thrown between the time `vwp` is filled with pointers and the time you get around to deleting them, you've leaked resources again. Fortunately, both problems can be overcome.

To turn your `for_each`-like loop into an actual use of `for_each`, you need to turn `delete` into a function object. That's child's play, assuming you have a child who likes to play with the STL:

```

template<typename T>
struct DeleteObject:                // Item 40 describes why
    public unary_function<const T*, void> { // this inheritance is here
    void operator()(const T* ptr) const
    {
        delete ptr;
    }
};

```

Now you can do this:

```

void doSomething()
{
    ...                                // as before
    for_each(vwp.begin(), vwp.end(), DeleteObject<Widget>());
}

```

Unfortunately, this makes you specify the type of objects that `DeleteObject` will be deleting (in this case, `Widget`). That's annoying. `vwp` is a `vector<Widget*>`, so *of course* `DeleteObject` will be deleting `Widget*` pointers! Duh! This kind of redundancy is more than just annoying, because it can lead to bugs that are difficult to track down. Suppose, for example, somebody ill-advisedly decides to inherit from `string`:

```

class SpecialString: public string { ... };

```

This is risky from the get-go, because `string`, like all the standard STL containers, lacks a virtual destructor, and publicly inheriting from classes without virtual destructors is a major C++ no-no. (For details, consult any good book on C++. In *Effective C++*, the place to look is

Item 14.) Still, some people do this kind of thing, so let's consider how the following code would behave:

```
void doSomething()
{
    deque<SpecialString*> dssp;

    ...

    for_each( dssp.begin(), dssp.end(),           // undefined behavior! Deletion
              DeleteObject<string>());           // of a derived object via a base
}                                                  // class pointer where there is
                                                  // no virtual destructor
```

Note how `dssp` is declared to hold `SpecialString*` pointers, but the author of the `for_each` loop has told `DeleteObject` that it will be deleting `string*` pointers. It's easy to understand how such an error could arise. `SpecialString` undoubtedly acts a lot like a string, so one can forgive its clients if they occasionally forget that they are using `SpecialStrings` instead of strings.

We can eliminate the error (as well as reduce the number of key-strokes required of `DeleteObject`'s clients) by having compilers deduce the type of pointer being passed to `DeleteObject::operator()`. All we need to do is move the templatzation from `DeleteObject` to its `operator()`:

```
struct DeleteObject {                               // templatzation and base
                                                    // class removed here

    template<typename T>                           // templatzation added here
    void operator()(const T* ptr) const
    {
        delete ptr;
    }
};
```

Compilers know the type of pointer being passed to `DeleteObject::operator()`, so we have them automatically instantiate an `operator()` taking that type of pointer. The downside to this type deduction is that we give up the ability to make `DeleteObject` adaptable (see [Item 40](#)). Considering how `DeleteObject` is designed to be used, it's difficult to imagine how that could be a problem.

With this new version of `DeleteObject`, the code for `SpecialString` clients looks like this:

```
void doSomething()
{
    deque<SpecialString*> dssp;

    ...
```

```

    for_each( dssp.begin(), dssp.end(),
              DeleteObject());           // ah! well-defined behavior!
}

```

Straightforward and type-safe, just the way we like it.

But still not exception-safe. If an exception is thrown after the Special-Strings are newed but before invocation of the call to `for_each`, it's Leakapalooza. That problem can be addressed in a variety of ways, but the simplest is probably to replace the container of pointers with a container of *smart pointers*, typically reference-counted pointers. (If you're unfamiliar with the notion of smart pointers, you should be able to find a description in any intermediate or advanced C++ book. In *More Effective C++*, the material is in Item 28.)

The STL itself contains no reference-counting smart pointer, and writing a good one — one that works correctly all the time — is tricky enough that you don't want to do it unless you have to. I published the code for a reference-counting smart pointer in *More Effective C++* in 1996, and despite basing it on established smart pointer implementations and submitting it to extensive pre-publication reviewing by experienced developers, a small parade of valid bug reports has trickled in for years. The number of subtle ways in which reference-counting smart pointers can fail is remarkable. (For details, consult the *More Effective C++* errata list [28].)

Fortunately, there's rarely a need to write your own, because proven implementations are not difficult to find. One such smart pointer is `shared_ptr` in the Boost library (see [Item 50](#)). With Boost's `shared_ptr`, this Item's original example can be rewritten as follows:

```

void doSomething()
{
    typedef boost::shared_ptr<Widget> SPW;           // SPW = "shared_ptr
                                                    //      to Widget"
    vector<SPW> vwp;
    for (int i = 0; i < SOME_MAGIC_NUMBER; ++i)
        vwp.push_back(SPW(new Widget));             // create an SPW from a
                                                    // Widget*, then do a
                                                    // push_back on it
    ...                                              // use vwp
}                                                    // no Widgets are leaked here, not
                                                    // even if an exception is thrown
                                                    // in the code above

```

One thing you must *never* be fooled into thinking is that you can arrange for pointers to be deleted automatically by creating containers

of `auto_ptr`s. That's a horrible thought, one so perilous, I've devoted [Item 8](#) to why you should avoid it.

All you really need to remember is that STL containers are smart, but they're not smart enough to know whether to delete the pointers they contain. To avoid resource leaks when you have containers of pointers that should be deleted, you must either replace the pointers with smart reference-counting pointer objects (such as Boost's `shared_ptr`) or you must manually delete each pointer in the container before the container is destroyed.

Finally, it may have crossed your mind that if a struct like `DeleteObject` can make it easier to avoid resource leaks for containers holding pointers to objects, it should be possible to create a similar `DeleteArray` struct to make it easier to avoid resource leaks for containers holding pointers to arrays. Certainly it is *possible*, but whether it is advisable is a different matter. [Item 13](#) explains why dynamically allocated arrays are almost always inferior to `vector` and `string` objects, so before you sit down to write `DeleteArray`, please review [Item 13](#) first. With luck, you'll decide that `DeleteArray` is a struct whose time will never come.

Item 8: Never create containers of `auto_ptr`s.

Frankly, this Item shouldn't need to be in *Effective STL*. Containers of `auto_ptr` (COAPs) are prohibited. Code attempting to use them shouldn't compile. The C++ Standardization Committee expended untold effort to arrange for that to be the case.[†] I shouldn't have to say anything about COAPs, because your compilers should have plenty to say about such containers, and all of it should be uncomplimentary.

Alas, many programmers use STL platforms that fail to reject COAPs. Alas even more, many programmers see in COAPs the chimera of a simple, straightforward, efficient solution to the resource leaks that often accompany containers of pointers (see [Items 7](#) and [33](#)). As a result, many programmers are tempted to use COAPs, even though it's not supposed to be possible to create them.

I'll explain in a moment why the spectre of COAPs was so alarming that the Standardization Committee took specific steps to make them illegal. Right now, I want to focus on a disadvantage that requires no knowledge of `auto_ptr`, or even of containers: COAPs aren't portable. How could they be? The Standard for C++ forbids them, and better

[†] If you're interested in the tortured history of `auto_ptr` standardization, point your web browser to the *auto_ptr Update page* [\[29\]](#) at the *More Effective C++* web site.

STL platforms already enforce this. It's reasonable to assume that as time goes by, STL platforms that currently fail to enforce this aspect of the Standard will become more compliant, and when that happens, code that uses COAPs will be even less portable than it is now. If you value portability (and you should), you'll reject COAPs simply because they fail the portability test.

But maybe you're not of a portability mind-set. If that's the case, kindly allow me to remind you of the unique — some would say bizarre — definition of what it means to copy an `auto_ptr`.

When you copy an `auto_ptr`, ownership of the object pointed to by the `auto_ptr` is transferred to the copying `auto_ptr`, and the copied `auto_ptr` is set to `NULL`. You read that right: *to copy an `auto_ptr` is to change its value*:

```
auto_ptr<Widget> pw1(new Widget); // pw1 points to a Widget
auto_ptr<Widget> pw2(pw1);        // pw2 points to pw1's Widget;
                                   // pw1 is set to NULL. (Ownership
                                   // of the Widget is transferred
                                   // from pw1 to pw2.)

pw1 = pw2;                        // pw1 now points to the Widget
                                   // again; pw2 is set to NULL
```

This is certainly unusual, and perhaps it's interesting, but the reason you (as a user of the STL) care is that it leads to some *very* surprising behavior. For example, consider this innocent-looking code, which creates a vector of `auto_ptr<Widget>` and then sorts it using a function that compares the values of the pointed-to Widgets:

```
bool widgetAPCompare(const auto_ptr<Widget>& lhs,
                    const auto_ptr<Widget>& rhs)
{
    return *lhs < *rhs; // for this example, assume that
                        // operator< exists for Widgets
}

vector<auto_ptr<Widget> > widgets; // create a vector and then fill it
...                               // with auto_ptrs to Widgets;
                                   // remember that this should
                                   // not compile!

sort(widgets.begin(), widgets.end(), // sort the vector
     widgetAPCompare);
```

Everything here looks reasonable, and conceptually, everything is reasonable, but the results need not be reasonable at all. For example, one or more of the `auto_ptrs` in `widgets` may have been set to `NULL` during the sort. The act of sorting the vector may have changed its contents! It is worthwhile understanding how this can be.

It can be because one approach to implementing sort — a common approach, as it turns out — is to use some variation on the quicksort algorithm. The fine points of quicksort need not concern us, but the basic idea is that to sort a container, some element of the container is chosen as the “pivot element,” then a recursive sort is done on the values greater than and less than or equal to the pivot element. Within sort, such an approach could look something like this:

```
template<class RandomAccessIterator,           // this declaration for
        class Compare>                       // sort is copied straight
void sort(RandomAccessIterator first,         // out of the Standard
          RandomAccessIterator last,
          Compare comp)
{
    // this typedef is described below
    typedef typename iterator_traits<RandomAccessIterator>::value_type
        ElementType;
    RandomAccessIterator i;
    ...                                     // make i point to the pivot element
    ElementType pivotValue(*i);            // copy the pivot element into a
                                           // local temporary variable; see
                                           // discussion below
    ...                                     // do the rest of the sorting work
}
```

Unless you’re an experienced reader of STL source code, this may look intimidating, but it’s really not that bad. The only tricky part is the reference to `iterator_traits<RandomAccessIterator>::value_type`, and that’s just the fancy STL way of referring to the type of object pointed to by the iterators passed to sort. (When we refer to `iterator_traits<RandomAccessIterator>::value_type`, we must precede it by `typename`, because it’s the name of a type that’s dependent on a template parameter, in this case, `RandomAccessIterator`. For more information about this use of `typename`, turn to [page 7](#).)

The troublesome statement in the code above is this one,

```
    ElementType pivotValue(*i);
```

because it copies an element from the range being sorted into a local temporary object. In our case, the element is an `auto_ptr<Widget>`, so this act of copying silently sets the copied `auto_ptr` — the one in the vector — to `NULL`. Furthermore, when `pivotValue` goes out of scope, it will automatically delete the `Widget` it points to. By the time the call to `sort` returns, the contents of the vector will have changed, and at least one `Widget` will have been deleted. It’s possible that several vector elements will have been set to `NULL` and several `Widgets` will have been

deleted, because quicksort is a recursive algorithm, so it could well have copied a pivot element at each level of recursion.

This is a nasty trap to fall into, and that's why the Standardization Committee worked so hard to make sure you're not supposed to be able to fall into it. Honor its work on your behalf, then, by never creating containers of `auto_ptr`s, even if your STL platforms allow it.

If your goal is a container of smart pointers, this doesn't mean you're out of luck. Containers of smart pointers are fine, and [Item 50](#) describes where you can find smart pointers that mesh well with STL containers. It's just that `auto_ptr` is not such a smart pointer. Not at all.

Item 9: Choose carefully among erasing options.

Suppose you have a standard STL container, `c`, that holds ints,

```
Container<int> c;
```

and you'd like to get rid of all the objects in `c` with the value 1963. Surprisingly, the way to accomplish this task varies from container type to container type; no single approach works for all of them.

If you have a contiguous-memory container (vector, deque, or string — see [Item 1](#)), the best approach is the erase-remove idiom (see [Item 32](#)):

```
c.erase(remove(c.begin(), c.end(), 1963), // the erase-remove idiom is
          c.end());                       // the best way to get rid of
                                         // elements with a specific
                                         // value when c is a vector,
                                         // string, or deque
```

This approach works for lists, too, but, as [Item 44](#) explains, the list member function `remove` is more efficient:

```
c.remove(1963);                          // the remove member function is the
                                         // best way to get rid of elements with
                                         // a specific value when c is a list
```

When `c` is a standard associative container (i.e., a set, multiset, map, or multimap), the use of anything named `remove` is completely wrong. Such containers have no member function named `remove`, and using the `remove` algorithm might overwrite container values (see [Item 32](#)), potentially corrupting the container. (For details on such corruption, consult [Item 22](#), which also explains why trying to use `remove` on maps and multimaps won't compile, and trying to use it on sets and multisets may not compile.)

No, for associative containers, the proper way to approach the problem is to call `erase`:

```
c.erase(1963);           // the erase member function is the
                        // best way to get rid of elements with
                        // a specific value when c is a
                        // standard associative container
```

Not only does this do the right thing, it takes only logarithmic time for `set` and `map`, and, for `multiset` and `multimap`, time linear in the number of elements with the specified value. (The `remove`-based techniques for sequence containers require time linear in the number of elements in the container.) Furthermore, the associative container `erase` member function has the advantage of being based on equivalence instead of equality, a distinction whose importance is explained in [Item 19](#).

Let's now revise the problem slightly. Instead of getting rid of every object in `c` that has a particular value, let's eliminate every object for which the following predicate (see [Item 39](#)) returns `true`:

```
bool badValue(int x);           // returns whether x is "bad"
```

For the sequence containers (`vector`, `string`, `deque`, and `list`), all we need to do is replace each use of `remove` with `remove_if`, and we're done:

```
c.erase(remove_if(c.begin(), c.end(), badValue), // this is the best way to
          c.end());                               // get rid of objects
                                                // where badValue
                                                // returns true when c is
                                                // a vector, string, or
                                                // deque

c.remove_if(badValue);           // this is the best way to get rid of
                                // objects where badValue returns
                                // true when c is a list
```

For the standard associative containers, it's not quite so straightforward. There are two ways to approach the problem, one easier to code, one more efficient. The easier-but-less-efficient solution uses `remove_copy_if` to copy the values we want into a new container, then swaps the contents of the original container with those of the new one:

```
AssocContainer<int> c;           // c is now one of the
...                               // standard associative
                                // containers

AssocContainer<int> goodValues;   // temporary container
                                // to hold unremoved
                                // values

remove_copy_if(c.begin(), c.end(),           // copy desired
               inserter(goodValues,          // values from c to
                       goodValues.end()),    // goodValues
               badValue);

c.swap(goodValues);              // swap the contents of
                                // c and goodValues
```

The drawback to this approach is that it involves copying all the elements that aren't being removed, and such copying might cost us more than we're interested in paying.

We can dodge that bill by removing the elements from the original container directly. However, because associative containers offer no member function akin to `remove_if`, we must write a loop to iterate over the elements in `c`, erasing elements as we go.

Conceptually, the task is simple, and in fact, the code is simple, too. Unfortunately, the code that does the job correctly is rarely the code that springs to mind. For example, this is what many programmers come up with first:

```
AssocContainer<int> c;
...
for (AssocContainer<int>::iterator i = c.begin(); // clear, straightforward,
    i != c.end();                               // and buggy code to
    ++i) {                                     // erase every element
    if (badValue(*i)) c.erase(i);              // in c where badValue
                                              // returns true; don't
                                              // do this!
}
```

Alas, this has undefined behavior. When an element of a container is erased, all iterators that point to that element are invalidated. Once `c.erase(i)` returns, `i` has been invalidated. That's bad news for this loop, because after `erase` returns, `i` is incremented via the `++i` part of the `for` loop.

To avoid this problem, we have to make sure we have an iterator to the next element of `c` before we call `erase`. The easiest way to do that is to use postfix increment on `i` when we make the call:

```
AssocContainer<int> c;
...
for (AssocContainer<int>::iterator i = c.begin(); // the 3rd part of the for
    i != c.end();                               // loop is empty; i is now
    /* nothing */) {                             // incremented below
    if (badValue(*i)) c.erase(i++);               // for bad values, pass the
    else ++i;                                     // current i to erase and
}                                                  // increment i as a side
                                              // effect; for good values,
                                              // just increment i
```

This approach to calling `erase` works, because the value of the expression `i++` is `i`'s old value, but as a side effect, `i` is incremented. Hence, we pass `i`'s old (unincremented) value to `erase`, but we also increment `i` itself before `erase` begins executing. That's exactly what we want. As I

said, the code is simple, it's just not what most programmers come up with the first time they try.

Let's now revise the problem further. Instead of merely erasing each element for which `badValue` returns true, we also want to write a message to a log file each time an element is erased.

For the associative containers, this is as easy as easy can be, because it requires only a trivial modification to the loop we just developed:

```

ofstream logFile;                                // log file to write to
AssocContainer<int> c;
...
for (AssocContainer<int>::iterator i = c.begin(); // loop conditions are the
     i != c.end(); ) {                          // same as before
    if (badValue(*i)) {
        logFile << "Erasing " << *i << '\n';    // write log file
        c.erase(i++);                          // erase element
    }
    else ++i;
}

```

It's vector, string, and deque that now give us trouble. We can't use the erase-remove idiom any longer, because there's no way to get erase or remove to write the log file. Furthermore, we can't use the loop we just developed for associative containers, because it yields undefined behavior for vectors, strings, and deques! Recall that for such containers, invoking erase not only invalidates all iterators pointing to the erased element, it also invalidates all iterators *beyond* the erased element. In our case, that includes all iterators beyond `i`. It doesn't matter if we write `i++`, `++i`, or anything else you can think of, because none of the resulting iterators is valid.

We must take a different tack with vector, string, and deque. In particular, we must take advantage of erase's return value. That return value is exactly what we need: it's a valid iterator pointing to the element following the erased element once the erase has been accomplished. In other words, we write this:

```

for (SeqContainer<int>::iterator i = c.begin();
     i != c.end(); ) {
    if (badValue(*i)) {
        logFile << "Erasing " << *i << '\n';
        i = c.erase(i);                                // keep i valid by assigning
                                                         // erase's return value to it
    }
    else ++i;
}

```

This works wonderfully, but only for the standard sequence containers. Due to reasoning one might question (Item 5 does), `erase`'s return type for the standard associative containers is `void`.[†] For those containers, you have to use the postincrement-the-iterator-you-pass-to-erase technique. (Incidentally, this kind of difference between coding for sequence containers and coding for associative containers is an example of why it's generally ill-advised to try to write container-independent code — see Item 2.)

Lest you be left wondering what the appropriate approach for `list` is, it turns out that for purposes of iterating and erasing, you can treat `list` like a `vector`/`string`/`deque` or you can treat it like an associative container; both approaches work for `list`.

If we take stock of everything we've covered in this Item, we come to the following conclusions:

- **To eliminate all objects in a container that have a particular value:**

If the container is a `vector`, `string`, or `deque`, use the `erase-remove` idiom.

If the container is a `list`, use `list::remove`.

If the container is a standard associative container, use its `erase` member function.

- **To eliminate all objects in a container that satisfy a particular predicate:**

If the container is a `vector`, `string`, or `deque`, use the `erase-remove_if` idiom.

If the container is a `list`, use `list::remove_if`.

If the container is a standard associative container, use `remove_copy_if` and `swap`, or write a loop to walk the container elements, being sure to postincrement your iterator when you pass it to `erase`.

- **To do something inside the loop (in addition to erasing objects):**

If the container is a standard sequence container, write a loop to walk the container elements, being sure to update your iterator with `erase`'s return value each time you call it.

[†] This is true only for the forms of `erase` that take iterator arguments. Associative containers also offer a form of `erase` taking a value argument, and that form returns the number of elements erased. Here, however, we're concerned only with erasing things via iterators.

If the container is a standard associative container, write a loop to walk the container elements, being sure to postincrement your iterator when you pass it to erase.

As you can see, there's more to erasing container elements effectively than just calling `erase`. The best way to approach the matter depends on how you identify which objects to erase, the type of container they're stored in, and what (if anything) you want to do while you're erasing them. As long as you're careful and heed the advice in this Item, you'll have no trouble. If you're not careful, you run the risk of producing code that's needlessly inefficient or that yields undefined behavior.

Item 10: Be aware of allocator conventions and restrictions.

Allocators are weird. They were originally developed as an abstraction for memory models that would allow library developers to ignore the distinction between near and far pointers in certain 16-bit operating systems (i.e., DOS and its pernicious spawn), but that effort failed. Allocators were also designed to facilitate the development of memory managers that are full-fledged objects, but it turned out that that approach led to efficiency degradations in some parts of the STL. To avoid the efficiency hits, the C++ Standardization Committee added wording to the Standard that emasculated allocators as objects, yet simultaneously expressed the hope that they would suffer no loss of potency from the operation.

There's more. Like `operator new` and `operator new[]`, STL allocators are responsible for allocating (and deallocating) raw memory, but an allocator's client interface bears little resemblance to that of `operator new`, `operator new[]`, or even `malloc`. Finally (and perhaps most remarkable), most of the standard containers *never ask* their associated allocator for memory. Never. The end result is that allocators are, well, allocators are weird.

That's not their fault, of course, and at any rate, it doesn't mean they're useless. However, before I explain what allocators are good for (that's the topic of [Item 11](#)), I need to explain what they're not good for. There are a number of things that allocators *seem* to be able to do, but can't, and it's important that you know the boundaries of the field before you try to start playing. If you don't, you'll get injured for sure. Besides, the truth about allocators is so peculiar, the mere act of summarizing it is both enlightening and entertaining. At least I hope it is.

The list of restrictions on allocators begins with their vestigial typedefs for pointers and references. As I mentioned, allocators were originally conceived of as abstractions for memory models, and as such it made

sense for allocators to provide typedefs for pointers and references in the memory model they defined. In the C++ standard, the default allocator for objects of type `T` (cunningly known as `allocator<T>`) offers the typedefs `allocator<T>::pointer` and `allocator<T>::reference`, and it is expected that user-defined allocators will provide these typedefs, too.

Old C++ hands immediately recognize that this is suspect, because there's no way to fake a reference in C++. Doing so would require the ability to overload operator. ("operator dot"), and that's not permitted. In addition, creating objects that act like references is an example of the use of *proxy objects*, and proxy objects lead to a number of problems. (One such problem motivates [Item 18](#). For a comprehensive discussion of proxy objects, turn to Item 30 of *More Effective C++*, where you can read about when they work as well as when they do not.)

In the case of allocators in the STL, it's not any technical shortcomings of proxy objects that render the pointer and reference typedefs impotent, it's the fact that the Standard explicitly allows library implementers to assume that every allocator's pointer typedef is a synonym for `T*` and every allocator's reference typedef is the same as `T&`. That's right, library implementers may ignore the typedefs and use raw pointers and references directly! So even if you could somehow find a way to write an allocator that successfully provided new pointer and reference types, it wouldn't do any good, because the STL implementations you were using would be free to ignore your typedefs. Neat, huh?

While you're admiring that quirk of standardization, I'll introduce another. Allocators are objects, and that means they may have member functions, nested types and typedefs (such as pointer and reference), etc., but the Standard says that an implementation of the STL is permitted to assume that all allocator objects of the same type are equivalent and always compare equal. Offhand, that doesn't sound so awful, and there's certainly good motivation for it. Consider this code:

```
template<typename T>                // a user-defined allocator
class SpecialAllocator { ... };      // template
typedef SpecialAllocator<Widget> SAW; // SAW = "SpecialAllocator
                                     //      for Widgets"

list<Widget, SAW> L1;
list<Widget, SAW> L2;

...

L1.splice(L1.begin(), L2);           // move L2's nodes to the
                                     // front of L1
```

Recall that when list elements are spliced from one list to another, nothing is copied. Instead, a few pointers are adjusted, and the list

nodes that used to be in one list find themselves in another. This makes splicing operations both fast and exception-safe. In the example above, the nodes that were in L2 prior to the splice are in L1 after the splice.

When L1 is destroyed, of course, it must destroy all its nodes (and deallocate their memory), and because it now contains nodes that were originally part of L2, L1's allocator must deallocate the nodes that were originally allocated by L2's allocator. Now it should be clear why the Standard permits implementers of the STL to assume that allocators of the same type are equivalent. It's so memory allocated by one allocator object (such as L2's) may be safely deallocated by another allocator object (such as L1's). Without being able to make such an assumption, splicing operations would be more difficult to implement. Certainly they wouldn't be as efficient as they can be now. (The existence of splicing operations affects other parts of the STL, too. For another example, see [Item 4](#).)

That's all well and good, but the more you think about it, the more you'll realize just how draconian a restriction it is that STL implementations may assume that allocators of the same type are equivalent. It means that portable allocator objects — allocators that will function correctly under different STL implementations — may not have state. Let's be explicit about this: it means that *portable allocators may not have any nonstatic data members*, at least not any that affect their behavior. None. Nada. That means, for example, you can't have one `SpecialAllocator<int>` that allocates from one heap and a different `SpecialAllocator<int>` that allocates from a different heap. Such allocators wouldn't be equivalent, and STL implementations exist where attempts to use both allocators could lead to corrupt runtime data structures.

Notice that this is a *runtime* issue. Allocators with state will compile just fine. They just may not run the way you expect them to. The responsibility for ensuring that all allocators of a given type are equivalent is yours. Don't expect compilers to issue a warning if you violate this constraint.

In fairness to the Standardization Committee, I should point out that it included the following statement immediately after the text that permits STL implementers to assume that allocators of the same type are equivalent:

Implementors are encouraged to supply libraries that ... support non-equal instances. In such implementations, ... the

semantics of containers and algorithms when allocator instances compare non-equal are implementation-defined.

This is a lovely sentiment, but as a user of the STL who is considering the development of a custom allocator with state, it offers you next to nothing. You can take advantage of this statement only if (1) you know that the STL implementations you are using support inequivalent allocators, (2) you are willing to delve into their documentation to determine whether the implementation-defined behavior of “non-equal” allocators is acceptable to you, and (3) you’re not concerned about porting your code to STL implementations that may take advantage of the latitude expressly extended to them by the Standard. In short, this paragraph — paragraph 5 of section 20.1.5, for those who insist on knowing — is the Standard’s “I have a dream” speech for allocators. Until that dream becomes common reality, programmers concerned about portability will limit themselves to custom allocators with no state.

I remarked earlier that allocators are like `operator new` in that they allocate raw memory, but their interface is different. This becomes apparent if you look at the declaration of the most common forms of `operator new` and `allocator<T>::allocate`:

```
void* operator new(size_t bytes);  
pointer allocator<T>::allocate(size_type numObjects);  
                                // recall that “pointer” is a typedef  
                                // that’s virtually always T*
```

Both take a parameter specifying how much memory to allocate, but in the case of `operator new`, this parameter specifies a certain number of bytes, while in the case of `allocator<T>::allocate`, it specifies how many `T` objects are to fit in the memory. On a platform where `sizeof(int) == 4`, for example, you’d pass 4 to `operator new` if you wanted enough memory to hold an `int`, but you’d pass 1 to `allocator<int>::allocate`. (The type of this parameter is `size_t` in the case of `operator new`, while it’s `allocator<T>::size_type` in the case of `allocate`. In both cases, it’s an unsigned integral type, and typically `allocator<T>::size_type` is a typedef for `size_t`, anyway.) There’s nothing “wrong” about this discrepancy, but the inconsistent conventions between `operator new` and `allocator<T>::allocate` complicate the process of applying experience with custom versions of `operator new` to the development of custom allocators.

`operator new` and `allocator<T>::allocate` differ in return types, too. `operator new` returns a `void*`, which is the traditional C++ way of representing a pointer to uninitialized memory. `allocator<T>::allocate` returns a `T*` (via the pointer typedef), which is not only untraditional, it’s premeditated fraud. The pointer returned from `allocator<T>::allocate` doesn’t

point to a T object, because no T has yet been constructed! Implicit in the STL is the expectation that `allocator<T>::allocate`'s caller will eventually construct one or more T objects in the memory it returns (possibly via `allocator<T>::construct`, via `uninitialized_fill`, or via some application of `raw_storage_iterators`), though in the case of `vector::reserve` or `string::reserve`, that may never happen (see [Item 14](#)). The difference in return type between `operator new` and `allocator<T>::allocate` indicates a change in the conceptual model for uninitialized memory, and it again makes it harder to apply knowledge about implementing `operator new` to the development of custom allocators.

That brings us to the final curiosity of STL allocators, that most of the standard containers never make a single call to the allocators with which they are instantiated. Here are two examples:

```
list<int> L;                                // same as list<int, allocator<int> >;
                                           // allocator<int> is never asked to
                                           // allocate memory!

set<Widget, SAW> s;                        // recall that SAW is a typedef for
                                           // SpecialAllocator<Widget>; no
                                           // SAW will ever allocate memory!
```

This oddity is true for `list` and all the standard associative containers (`set`, `multiset`, `map`, and `multimap`). That's because these are *node-based* containers, i.e., containers based on data structures in which a new node is dynamically allocated each time a value is to be stored. In the case of `list`, the nodes are list nodes. In the case of the standard associative containers, the nodes are usually tree nodes, because the standard associative containers are typically implemented as balanced binary search trees.

Think for a moment about how a `list<T>` is likely to be implemented. The list itself will be made up of nodes, each of which holds a T object as well as pointers to the next and previous nodes in the list:

```
template<typename T,                        // possible list
        typename Allocator = allocator<T> > // implementation
class list {
private:
    Allocator alloc;                       // allocator for objects of type T

    struct ListNode {                       // nodes in the linked list
        T data;
        ListNode *prev;
        ListNode *next;
    };

    ...

};
```

When a new node is added to the list, we need to get memory for it from an allocator, but we don't need memory for a `T`, we need memory for a `ListNode` that contains a `T`. That makes our `Allocator` object all but useless, because it doesn't allocate memory for `ListNode`s, it allocates memory for `T`s. Now you understand why `list` never asks its `Allocator` to do any allocation: the allocator can't provide what `list` needs.

What `list` needs is a way to get from the allocator type it has to the corresponding allocator for `ListNode`s. This would be tough were it not that, by convention, allocators provide a typedef that does the job. The typedef is called `other`, but it's not quite that simple, because `other` is a typedef nested inside a struct called `rebind`, which itself is a template nested inside the allocator — which itself is a template!

Please don't try to think about that last sentence. Instead, look at the code below, then proceed directly to the explanation that follows.

```
template<typename T>           // the standard allocator is declared
class allocator {              // like this, but this could be a user-
public:                          // written allocator template, too

    template<typename U>
    struct rebind {
        typedef allocator<U> other;
    };

    ...
};
```

In the code implementing `list<T>`, there is a need to determine the type of the allocator for `ListNode`s that corresponds to the allocator we have for `T`s. The type of the allocator we have for `T`s is the template parameter `Allocator`. That being the case, the type of the corresponding allocator for `ListNode`s is this:

```
Allocator::rebind<ListNode>::other
```

Stay with me here. Every allocator template `A` (e.g., `std::allocator`, `SpecialAllocator`, etc.) is expected to have a nested struct template called `rebind`. `rebind` takes a single type parameter, `U`, and defines nothing but a typedef, `other`. `other` is simply a name for `A<U>`. As a result, `list<T>` can get from its allocator for `T` objects (called `Allocator`) to the corresponding allocator for `ListNode` objects by referring to `Allocator::rebind<ListNode>::other`.

Maybe this makes sense to you, maybe it doesn't. (If you stare at it long enough, it will, but you may have to stare a while. I know I had to.) As a user of the STL who may want to write a custom allocator, you don't really need to know how it works. What you do need to know is that if you choose to write allocators and use them with the stan-

dard containers, your allocators must provide the `rebind` template, because standard containers assume it will be there. (For debugging purposes, it's also helpful to know why node-based containers of `T` objects never ask for memory from the allocators for `T` objects.)

Hallelujah! We are finally done examining the idiosyncrasies of allocators. Let us therefore summarize the things you need to remember if you ever want to write a custom allocator.

- Make your allocator a template, with the template parameter `T` representing the type of objects for which you are allocating memory.
- Provide the typedefs `pointer` and `reference`, but always have `pointer` be `T*` and `reference` be `T&`.
- Never give your allocators per-object state. In general, allocators should have no nonstatic data members.
- Remember that an allocator's `allocate` member functions are passed the number of *objects* for which memory is required, not the number of bytes needed. Also remember that these functions return `T*` pointers (via the `pointer` typedef), even though no `T` objects have yet been constructed.
- Be sure to provide the nested `rebind` template on which standard containers depend.

Most of what you have to do to write your own allocator is reproduce a fair amount of boilerplate code, then tinker with a few member functions, notably `allocate` and `deallocate`. Rather than writing the boilerplate from scratch, I suggest you begin with the code at Josuttis' sample allocator web page [23] or in Austern's article, "What Are Allocators Good For?" [24].

Once you've digested the information in this Item, you'll know a lot about what allocators *cannot* do, but that's probably not what you want to know. Instead, you'd probably like to know what allocators *can* do. That's a rich topic in its own right, a topic I call "[Item 11](#)."

Item 11: Understand the legitimate uses of custom allocators.

So you've benchmarked, profiled, and experimented your way to the conclusion that the default STL memory manager (i.e., `allocator<T>`) is too slow, wastes memory, or suffers excessive fragmentation for your STL needs, and you're certain you can do a better job yourself. Or you discover that `allocator<T>` takes precautions to be thread-safe, but

you're interested only in single-threaded execution and you don't want to pay for the synchronization overhead you don't need. Or you know that objects in certain containers are typically used together, so you'd like to place them near one another in a special heap to maximize locality of reference. Or you'd like to set up a unique heap that corresponds to shared memory, then put one or more containers in that memory so they can be shared by other processes. Congratulations! Each of these scenarios corresponds to a situation where custom allocators are well suited to the problem.

For example, suppose you have special routines modeled after `malloc` and `free` for managing a heap of shared memory,

```
void* mallocShared(size_t bytesNeeded);
void freeShared(void *ptr);
```

and you'd like to make it possible to put the contents of STL containers in that shared memory. No problem:

```
template<typename T>
class SharedMemoryAllocator {
public:
    ...
    pointer allocate(size_type numObjects, const void *localityHint = 0)
    {
        return static_cast<pointer>(mallocShared(numObjects * sizeof(T)));
    }
    void deallocate(pointer ptrToMemory, size_type numObjects)
    {
        freeShared(ptrToMemory);
    }
    ...
};
```

For information on the pointer type as well as the cast and the multiplication inside `allocate`, see [Item 10](#).

You could use `SharedMemoryAllocator` like this:

```
// convenience typedef
typedef
    vector<double, SharedMemoryAllocator<double> > SharedDoubleVec;
...
{
    ...
    SharedDoubleVec v;
    ...
}
// begin some block
// create a vector whose elements
// are in shared memory
// end the block
```

The wording in the comment next to `v`'s definition is important. `v` is using a `SharedMemoryAllocator`, so the memory `v` allocates to hold its elements will come from shared memory. `v` itself, however — including all its data members — will almost certainly *not* be placed in shared memory. `v` is just a normal stack-based object, so it will be located in whatever memory the runtime system uses for all normal stack-based objects. That's almost never shared memory. To put both `v`'s contents and `v` itself into shared memory, you'd have to do something like this:

```
void *pVectorMemory =           // allocate enough shared
    mallocShared(sizeof(SharedDoubleVec)); // memory to hold a
                                         // SharedDoubleVec object

SharedDoubleVec *pv =           // use "placement new" to
    new (pVectorMemory) SharedDoubleVec; // create a SharedDoubleVec
                                         // object in the memory;
                                         // see below

...                               // use the object (via pv)

pv->~SharedDoubleVec();          // destroy the object in the
                                // shared memory

freeShared(pVectorMemory);       // deallocate the initial
                                // chunk of shared memory
```

I hope the comments make clear how this works. Fundamentally, you acquire some shared memory, then construct a vector in it that uses shared memory for its own internal allocations. When you're done with the vector, you invoke its destructor, then release the memory the vector occupied. The code isn't terribly complicated, but it's a lot more demanding than just declaring a local variable as we did above. Unless you really need a container (as opposed to its elements) to be in shared memory, I encourage you to avoid this manual four-step allocate/construct/destroy/deallocate process.

In this example, you've doubtless noticed that the code ignores the possibility that `mallocShared` might return a null pointer. Obviously, production code would have to take such a possibility into account. Also, construction of the vector in the shared memory is accomplished by "placement new." If you're unfamiliar with placement new, your favorite C++ text should be able to introduce you. If that text happens to be *More Effective C++*, you'll find that the pleasantries are exchanged in Item 8.

As a second example of the utility of allocators, suppose you have two heaps, identified by the classes `Heap1` and `Heap2`. Each heap class has static member functions for performing allocation and deallocation:

```

class Heap1 {
public:
    ...
    static void* alloc(size_t numBytes, const void *memoryBlockToBeNear);
    static void dealloc(void *ptr);
    ...
};

class Heap2 { ... };           // has the same alloc/dealloc interface

```

Further suppose you'd like to co-locate the contents of some STL containers in different heaps. Again, no problem. First you write an allocator designed to use classes like `Heap1` and `Heap2` for the actual memory management:

```

template<typename T, typename Heap>
class SpecificHeapAllocator {
public:
    ...
    pointer allocate(size_type numObjects, const void *localityHint = 0)
    {
        return static_cast<pointer> (Heap::alloc(numObjects * sizeof(T),
                                                  localityHint));
    }
    void deallocate(pointer ptrToMemory, size_type numObjects)
    {
        Heap::dealloc(ptrToMemory);
    }
    ...
};

```

Then you use `SpecificHeapAllocator` to cluster containers' elements together:

```

vector<int, SpecificHeapAllocator<int, Heap1> > v;    // put both v's and
set<int, SpecificHeapAllocator<int, Heap1> > s;        // s's elements in
                                                    // Heap1

list<Widget,
    SpecificHeapAllocator<Widget, Heap2> > L;        // put both L's and
map<int, string, less<int>,                               // m's elements in
    SpecificHeapAllocator<pair<const int, string>,         // Heap2
                          Heap2> > m;

```

In this example, it's quite important that `Heap1` and `Heap2` be *types* and not objects. The STL offers a syntax for initializing different STL containers with different allocator objects of the same type, but I'm not going to show you what it is. That's because if `Heap1` and `Heap2` were *objects* instead of types, they'd be inequivalent allocators, and

that would violate the equivalence constraint on allocators that is detailed in [Item 10](#).

As these examples demonstrate, allocators are useful in a number of contexts. As long as you obey the constraint that all allocators of the same type must be equivalent, you'll have no trouble employing custom allocators to control general memory management strategies, clustering relationships, and use of shared memory and other special heaps.

Item 12: Have realistic expectations about the thread safety of STL containers.

The world of standard C++ is rather sheltered and old-fashioned. In this rarefied world, all executables are statically linked. Neither memory-mapped files nor shared memory exist. There are no window systems, no networks, no databases, no other processes. That being the case, you should not be surprised to learn that the Standard says not a word about threading. The first expectation you should have about the thread safety of the STL, then, is that it will vary from implementation to implementation.

Of course, multithreaded programs are common, so most STL vendors strive to make their implementations work well in a threaded environment. Even when they do a good job, however, much of the burden remains on your shoulders, and it's important to understand why. There's only so much STL vendors can do to ease your multithreading pain, and you need to know what it is.

The gold standard in support for multithreading in STL containers (and the aspiration of most vendors) has been defined by SGI and is published at their STL Web Site [\[21\]](#). In essence, it says that the most you can hope for from an implementation is the following.

- **Multiple readers are safe.** Multiple threads may simultaneously read the contents of a single container, and this will work correctly. Naturally, there must not be any writers acting on the container during the reads.
- **Multiple writers to different containers are safe.** Multiple threads may simultaneously write to different containers.

That's all, and let me make clear that this is what you can *hope for*, not what you can *expect*. Some implementations offer these guarantees, but some do not.

Writing multithreaded code is hard, and many programmers wish that STL implementations were completely thread safe out of the box. Were that the case, programmers might hope to be relieved of the need to attend to concurrency control themselves. There's no doubt this would be a convenient state of affairs, but it would also be very difficult to achieve. Consider the following ways a library might try to implement such comprehensive container thread safety:

- Lock a container for the duration of each call to its member functions.
- Lock a container for the lifetime of each iterator it returns (via, e.g., calls to `begin` or `end`).
- Lock a container for the duration of each algorithm invoked on that container. (This actually makes no sense, because, as [Item 32](#) explains, algorithms have no way to identify the container on which they are operating. Nevertheless, we'll examine this option here, because it's instructive to see why it wouldn't work even if it were possible.)

Now consider the following code. It searches a `vector<int>` for the first occurrence of the value 5, and, if it finds one, changes that value to 0.

```
vector<int> v;  
...  
vector<int>::iterator first5(find(v.begin(), v.end(), 5));           // Line 1  
if (first5 != v.end()) {                                           // Line 2  
    *first5 = 0;                                                  // Line 3  
}
```

In a multithreaded environment, it's possible that a different thread will modify the data in `v` immediately after completion of Line 1. If that were to happen, the test of `first5` against `v.end` on Line 2 would be meaningless, because `v`'s values would be different from what they were at the end of Line 1. In fact, such a test could yield undefined results, because another thread could have intervened between Lines 1 and 2 and invalidated `first5`, perhaps by performing an insertion that caused the vector to reallocate its underlying memory. (That would invalidate all the vector's iterators. For details on this reallocation behavior, turn to [Item 14](#).) Similarly, the assignment to `*first5` on Line 3 is unsafe, because another thread might execute between Lines 2 and 3 in such a way as to invalidate `first5`, perhaps by erasing the element it points to (or at least *used* to point to).

None of the approaches to locking listed above would prevent these problems. The calls to `begin` and `end` in Line 1 both return too quickly

to offer any help, the iterators they generate last only until the end of that line, and `find` also returns at the end of that line.

For the code above to be thread safe, `v` must remain locked from Line 1 through Line 3, and it's difficult to imagine how an STL implementation could deduce that automatically. Bearing in mind the typically high cost of synchronization primitives (e.g., semaphores, mutexes, etc.), it's even more difficult to imagine how an implementation could do it without imposing a significant performance penalty on programs that knew *a priori* — that were *designed* in such a way — that no more than one thread had access to `v` during the course of Lines 1-3.

Such considerations explain why you can't expect any STL implementation to make your threading woes disappear. Instead, you'll have to manually take charge of synchronization control in these kinds of scenarios. In this example, you might do it like this:

```
vector<int> v;
...
getMutexFor(v);
vector<int>::iterator first5(find(v.begin(), v.end(), 5));
if (first5 != v.end()) {                                // this is now safe
    *first5 = 0;                                         // so is this
}
releaseMutexFor(v);
```

A more object-oriented solution is to create a `Lock` class that acquires a mutex in its constructor and releases it in its destructor, thus minimizing the chances that a call to `getMutexFor` will go unmatched by a call to `releaseMutexFor`. The essence of such a class (really a class template) is this:

```
template<typename Container>           // skeletal template for classes
class Lock {                           // that acquire and release mutexes
public:                                // for containers; many details
                                        // have been omitted

    Lock(const Container& container)
    : c(container)
    {
        getMutexFor(c);                // acquire mutex in the constructor
    }

    ~Lock()
    {
        releaseMutexFor(c);            // release it in the destructor
    }

private:
    const Container& c;
};
```

The idea of using a class (like `Lock`) to manage the lifetime of resources (such as mutexes) is generally known as *resource acquisition is initialization*, and you should be able to read about it in any comprehensive C++ textbook. A good place to start is Stroustrup's *The C++ Programming Language* [7], because Stroustrup popularized the idiom, but you can also turn to Item 9 of *More Effective C++*. No matter what source you consult, bear in mind that the above `Lock` is stripped to the bare essentials. An industrial-strength version would require a number of enhancements, but such a fleshing-out would have nothing to do with the STL. Furthermore, this minimalist `Lock` is enough to see how we could apply it to the example we've been considering:

```
vector<int> v;

...
{                                     // create new block
    Lock<vector<int> > lock(v);        // acquire mutex
    vector<int>::iterator first5(find(v.begin(), v.end(), 5));
    if (first5 != v.end()) {
        *first5 = 0;
    }
}                                     // close block, automatically
                                     // releasing the mutex
```

Because a `Lock` object releases the container's mutex in the `Lock`'s destructor, it's important that the `Lock` be destroyed as soon as the mutex should be released. To make that happen, we create a new block in which to define the `Lock`, and we close that block as soon as we no longer need the mutex. This sounds like we're just trading the need to call `releaseMutexFor` with the need to close a new block, but that's not an accurate assessment. If we forget to create a new block for the `Lock`, the mutex will still be released, but it may happen later than it should — when control reaches the end of the enclosing block. If we forget to call `releaseMutexFor`, we never release the mutex.

Furthermore, the `Lock`-based approach is robust in the presence of exceptions. C++ guarantees that local objects are destroyed if an exception is thrown, so `Lock` will release its mutex even if an exception is thrown while we're using the `Lock` object.[†] If we relied on manual calls to `getMutexFor` and `releaseMutexFor`, we'd never relinquish the mutex if an exception was thrown after calling `getMutexFor` but before calling `releaseMutexFor`.

Exceptions and resource management are important, but they're not the subject of this Item. This Item is about thread safety in the STL. When it comes to thread safety and STL containers, you can hope for

[†] There's a loophole in the guarantee. If an exception is not caught at all, the program will terminate. In that case, local objects (such as `lock`) may not have their destructors called. Some compilers call them, some do not. Both behaviors are valid.

a library implementation that allows multiple readers on one container and multiple writers on separate containers. You *can't* hope for the library to eliminate the need for manual concurrency control, and you can't *rely* on any thread support at all.

2

vector and string

All the STL containers are useful, but if you're like most C++ programmers, you'll find yourself reaching for vector and string more often than their compatriots. That's to be expected. vector and string are designed to replace most applications of arrays, and arrays are so useful, they've been included in every commercially successful programming language from COBOL to Java.

The Items in this chapter cover vectors and strings from a number of perspectives. We begin with a discussion of why the switch from arrays is worthwhile, then look at ways to improve vector and string performance, identify important variations in string implementations, examine how to pass vector and string data to APIs that understand only C, and learn how to eliminate excess memory allocation. We conclude with an examination of an instructive anomaly, `vector<bool>`, the little vector that couldn't.

Each of the Items in this chapter will help you take the two most useful containers in the STL and refine their application. By the time we're done, you'll know how to make them serve you even better.

Item 13: Prefer vector and string to dynamically allocated arrays.

The minute you decide to use new for a dynamic allocation, you adopt the following responsibilities:

1. You must make sure that somebody will later delete the allocation. Without a subsequent delete, your new will yield a resource leak.
2. You must ensure that the correct form of delete is used. For an allocation of a single object, "delete" must be used. For an array allocation, "delete []" is required. If the wrong form of delete is

used, results will be undefined. On some platforms, the program will crash at runtime. On others, it will silently blunder forward, sometimes leaking resources and corrupting memory as it goes.

3. You must make sure that `delete` is used exactly once. If an allocation is deleted more than once, results are again undefined.

That's quite a set of responsibilities, and I can't understand why you'd want to adopt them if it wasn't necessary. Thanks to `vector` and `string`, it isn't necessary anywhere near as often as it used to be.

Any time you find yourself getting ready to dynamically allocate an array (i.e., plotting to write `"new T[...]"`), you should consider using a `vector` or a `string` instead. (In general, use `string` when `T` is a character type and use `vector` when it's not, though later in this Item, we'll encounter a scenario where a `vector<char>` may be a reasonable design choice.) `vector` and `string` eliminate the burdens above, because they manage their own memory. Their memory grows as elements are added to these containers, and when a `vector` or `string` is destroyed, its destructor automatically destroys the elements in the container and deallocates the memory holding those elements.

In addition, `vector` and `string` are full-fledged STL sequence containers, so they put at your disposal the complete arsenal of STL algorithms that work on such containers. True, arrays can be used with STL algorithms, too, but arrays don't offer member functions like `begin`, `end`, and `size`, nor do they have nested typedefs like `iterator`, `reverse_iterator`, or `value_type`. And of course `char*` pointers can hardly compete with the scores of specialized member functions proffered by `string`. The more you work with the STL, the more jaundiced the eye with which you'll come to view built-in arrays.

If you're concerned about the legacy code you must continue to support, all of which is based on arrays, relax and use `vectors` and `strings` anyway. [Item 16](#) shows how easy it is to pass the data in `vectors` and `strings` to APIs that expect arrays, so integration with legacy code is generally not a problem.

Frankly, I can think of only one legitimate cause for concern in replacing dynamically allocated arrays with `vectors` or `strings`, and it applies only to `strings`. Many `string` implementations employ reference counting behind the scenes (see [Item 15](#)), a strategy that eliminates some unnecessary memory allocations and copying of characters and that can improve performance for many applications. In fact, the ability to optimize `strings` via reference counting was considered so important, the C++ Standardization Committee took specific steps to make sure it was a valid implementation.

Alas, one programmer's optimization is another's pessimization, and if you use reference-counted strings in a multithreaded environment, you may find that the time saved by avoiding allocations and copying is dwarfed by the time spent on behind-the-scenes concurrency control. (For details, consult Sutter's article, "Optimizations That Aren't (In a Multithreaded World)" [20].) If you're using reference-counted strings in a multithreaded environment, then, it makes sense to keep an eye out for performance problems arising from their support for thread safety.

To determine whether you're using a reference-counting implementation for string, it's often easiest to consult the documentation for your library. Because reference counting is considered an optimization, vendors generally tout it as a feature. An alternative is to look at the source code for your libraries' implementations of string. I don't generally recommend trying to figure things out from library source code, but sometimes it's the only way to find out what you need to know. If you choose this approach, remember that string is a typedef for `basic_string<char>` (and `wstring` is a typedef for `basic_string<wchar_t>`), so what you really want to look at is the template `basic_string`. The easiest thing to check is probably the class's copy constructor. Look to see if it increments a reference count somewhere. If it does, string is reference counted. If it doesn't, either string isn't reference counted or you misread the code. Ahem.

If the string implementations available to you are reference counted and you are running in a multithreaded environment where you've determined that string's reference counting support is a performance problem, you have at least three reasonable choices, none of which involves abandoning the STL. First, check to see if your library implementation is one that makes it possible to disable reference counting, often by changing the value of a preprocessor variable. This won't be portable, of course, but given the amount of work involved, it's worth investigating. Second, find or develop an alternative string implementation (or partial implementation) that doesn't use reference counting. Third, consider using a `vector<char>` instead of a string. `vector` implementations are not allowed to be reference counted, so hidden multithreading performance issues fail to arise. Of course, you forgo string's fancy member functions if you switch to `vector<char>`, but most of that functionality is available through STL algorithms anyway, so you're not so much giving up functionality as you are trading one syntax for another.

The upshot of all this is simple. If you're dynamically allocating arrays, you're probably taking on more work than you need to. To lighten your load, use vectors or strings instead.

Item 14: Use reserve to avoid unnecessary reallocations.

One of the most marvelous things about STL containers is that they automatically grow to accommodate as much data as you put into them, provided only that you don't exceed their maximum size. (To discover this maximum, just call the aptly named `max_size` member function.) For vector and string, growth is handled by doing the moral equivalent of a `realloc` whenever more space is needed. This `realloc`-like operation has four parts:

1. Allocate a new block of memory that is some multiple of the container's current capacity. In most implementations, vector and string capacities grow by a factor of between 1.5 and 2 each time.
2. Copy all the elements from the container's old memory into its new memory.
3. Destroy the objects in the old memory.
4. Deallocate the old memory.

Given all that allocation, deallocation, copying, and destruction, it should not stun you to learn that these steps can be expensive. Naturally, you don't want to perform them any more frequently than you have to. If that doesn't strike you as natural, perhaps it will when you consider that each time these steps occur, all iterators, pointers, and references into the vector or string are invalidated. That means that the simple act of inserting an element into a vector or string may also require updating other data structures that use iterators, pointers, or references into the vector or string being expanded.

The `reserve` member function allows you to minimize the number of reallocations that must be performed, thus avoiding the costs of reallocation and iterator/pointer/reference invalidation. Before I explain how `reserve` can do that, however, let me briefly recap four interrelated member functions that are sometimes confused. Among the standard containers, only vector and string offer all of these functions.

- `size()` tells you how many elements are in the container. It does *not* tell you how much memory the container has allocated for the elements it holds.

- `capacity()` tells you how many elements the container can hold in the memory it has already allocated. This is how many *total elements* the container can hold in that memory, not how many *more* elements it can hold. If you'd like to find out how much unoccupied memory a vector or string has, you must subtract `size()` from `capacity()`. If `size` and `capacity` return the same value, there is no empty space in the container, and the next insertion (via `insert` or `push_back`, etc.) will trigger the reallocation steps above.
- `resize(Container::size_type n)` forces the container to change to `n` the number of elements it holds. After the call to `resize`, `size` will return `n`. If `n` is smaller than the current size, elements at the end of the container will be destroyed. If `n` is larger than the current size, new elements will be added to the end of the container, either by copying a default-constructed element or by copying an object you provide via an additional parameter. If `n` is larger than the current capacity, a reallocation will take place before elements are added.
- `reserve(Container::size_type n)` forces the container to change its capacity to at least `n`, provided `n` is no less than the current size. This typically forces a reallocation, because the capacity needs to be increased. (If `n` is less than the current capacity, vector ignores the call and does nothing. string *may* reduce its capacity to the maximum of `size()` and `n`, but the string's size definitely remains unchanged. In my experience, using `reserve` to trim the excess capacity from a string is generally less successful than using “the swap trick,” which is the topic of [Item 17](#).) Note that calling `reserve` never changes the number of elements in a container (i.e., its size).

This recap should make it clear that reallocations (including their constituent memory allocations and deallocations, object copying and destruction, and invalidation of iterators, pointers, and references) occur whenever an element needs to be inserted and the container's capacity is insufficient. The key to avoiding reallocations, then, is to use `reserve` to set a container's capacity to a sufficiently large value as soon as possible, ideally right after the container is constructed.

For example, suppose you'd like to create a `vector<int>` holding the values 1–1000. Without using `reserve`, you might do it like this:

```
vector<int> v;  
for (int i = 1; i <= 1000; ++i) v.push_back(i);
```

This code will typically result in between two and 18 reallocations during the course of the loop.

Modifying the code to use `reserve` gives us this:

```
vector<int> v;  
v.reserve(1000);  
for (int i = 1; i <= 1000; ++i) v.push_back(i);
```

This should result in zero reallocations during the loop.

The relationship between size and capacity makes it possible to predict when an insertion will cause a vector or string to perform a reallocation, and that, in turn, makes it possible to predict when an insertion will invalidate iterators, pointers, and references into a container. For example, given this code,

```
string s;  
...  
if (s.size() < s.capacity()) {  
    s.push_back('x');  
}
```

the call to `push_back` can't invalidate iterators (other than to `s`'s end), pointers, or references into the string, because the string's capacity is guaranteed to be greater than its size. If, instead of performing a `push_back`, the code did an insert into an arbitrary location in the string, we'd still be guaranteed that no reallocation would take place during the insertion, but, in accord with the usual rule for iterator invalidation accompanying a string insertion, all iterators/pointers/references from the insertion point to the end of the string would be invalidated.

Getting back to the main point of this Item, there are two common ways to use `reserve` to avoid unneeded reallocations. The first is applicable when you know exactly or approximately how many elements will ultimately end up in your container. In that case, as in the vector code above, you simply reserve the appropriate amount of space in advance. The second way is to reserve the maximum space you could ever need, then, once you've added all your data, trim off any excess capacity. The trimming part isn't difficult, but I'm not going to show it here, because there's a trick to it. To learn the trick, turn to [Item 17](#).

Item 15: Be aware of variations in string implementations.

Bjarne Stroustrup once wrote an article with the curious title, "Sixteen Ways to Stack a Cat" [\[27\]](#). It turns out that there are almost as many ways to implement strings. As experienced and sophisticated software engineers, of course, we're supposed to pooh-pooh "imple-

mentation details,” but if Einstein is right and God is in the details, reality requires that we sometimes get religion. Even when the details don’t matter, having some idea about them puts us in a position to be *sure* they don’t matter.

For example, what is the size of a string object? In other words, what value does `sizeof(string)` return? This could be an important question if you’re keeping a close eye on memory consumption, and you’re thinking of replacing a raw `char*` pointer with a string object.

The news about `sizeof(string)` is “interesting,” which is almost certainly what you do not want to hear if you’re concerned about space. While it’s not uncommon to find string implementations in which strings are the same size as `char*` pointers, it’s also easy to find string implementations where each string is seven times that size. Why the difference? To understand that, we have to know what data a string is likely to store as well as where it might decide to store it.

Virtually every string implementation holds the following information:

- The **size** of the string, i.e., the number of characters it contains.
- The **capacity** of the memory holding the string’s characters. (For a review of the difference between a string’s size and its capacity, see [Item 14](#).)
- The **value** of the string, i.e., the characters making up the string.

In addition, a string may hold

- A copy of its **allocator**. For an explanation of why this field is optional, turn to [Item 10](#) and read about the curious rules governing allocators.

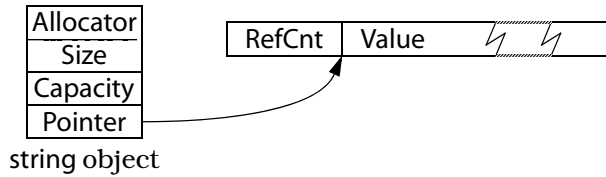
string implementations that depend on reference counting also contain

- The **reference count** for the value.

Different string implementations put these pieces of information together in different ways. To demonstrate what I mean, I’ll show you the data structures used by four different string implementations. There’s nothing special about these selections. They all come from STL implementations that are commonly used. They just happen to be the string implementations in the first four libraries I checked.

In implementation A, each string object contains a copy of its allocator, the string’s size, its capacity, and a pointer to a dynamically allocated buffer containing both the reference count (the “RefCnt”) and the string’s value. In this implementation, a string object using the default allocator is four times the size of a pointer. With a custom allocator,

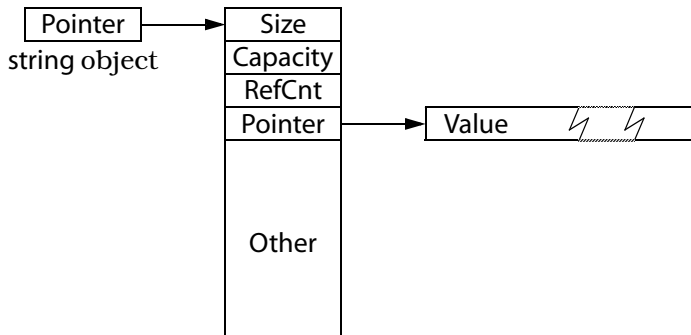
the string object would be bigger by about the size of the allocator object:



Implementation A

Implementation B's string objects are the same size as a pointer, because they contain nothing but a pointer to a struct. Again, this assumes that the default allocator is used. As in Implementation A, if a custom allocator is used, the string object's size will increase by about the size of the allocator object. In this implementation, use of the default allocator requires no space, thanks to an optimization present here but not present in Implementation A.

The object pointed to by B's string contains the string's size, capacity, and reference count, as well as a pointer to a dynamically allocated buffer holding the string's value. The object also contains some additional data related to concurrency control in multithreaded systems. Such data is outside our purview here, so I've just labeled that part of the data structure "Other:"

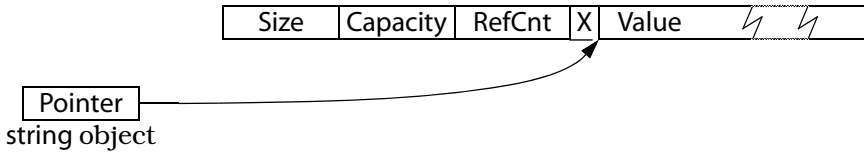


Implementation B

The box for "Other" is larger than the other boxes, because I've drawn the boxes to scale. If one box is twice the size of another, the larger box uses twice as many bytes as the smaller one. In Implementation B, the data for concurrency control is six times the size of a pointer.

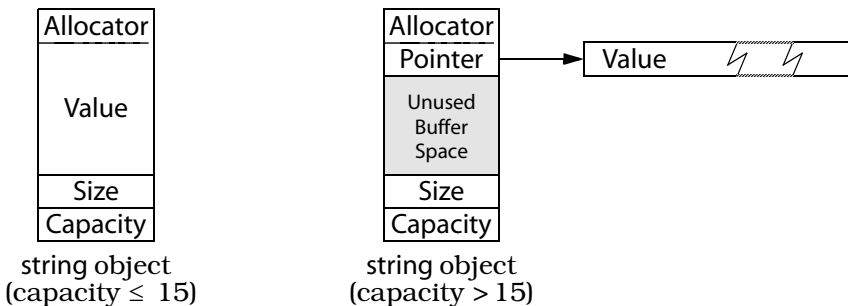
string objects under Implementation C are always the size of a pointer, but this pointer points to a dynamically allocated buffer containing everything related to the string: its size, capacity, reference count, and

value. There is no per-object allocator support. The buffer also holds some data concerning the shareability of the value, a topic we'll not consider here, so I've labeled it "X". (If you're interested in why a reference counted value might not be shareable in the first place, consult Item 29 of *More Effective C++*.)



Implementation C

Implementation D's string objects are seven times the size of a pointer (still assuming use of the default allocator). This implementation employs no reference counting, but each string contains an internal buffer large enough to represent string values of up to 15 characters. Small strings can thus be stored entirely within the string object, a feature sometimes known as the "small string optimization." When a string's capacity exceeds 15, the first part of the buffer is used as a pointer to dynamically allocated memory, and the string's value resides in that memory:



Implementation D

These diagrams do more than just prove I can read source code and draw pretty pictures. They also allow you to deduce that creation of a string in a statement such as this,

```
string s("Perse");           // Our dog is named "Persephone", but we
                             // usually just call her "Perse". Visit her web site
                             // at http://www.aristeia.com/Persephone/
```

will cost you zero dynamic allocations under Implementation D, one under Implementations A and C, and two under Implementation B (one for the object pointed to by the string object, one for the character

buffer pointed to by that object). If you're concerned about the number of times you dynamically allocate and deallocate memory, or if you're concerned about the memory overhead that often accompanies such allocations, you might want to shy away from Implementation B. On the other hand, the fact that Implementation B's data structure includes specific support for concurrency control in multithreaded systems might mean that it meets your needs better than implementations A or C, the number of dynamic allocations notwithstanding. (Implementation D requires no special support for multithreading, because it doesn't use reference counting. For more on the interaction of threading and reference counted strings, see [Item 13](#). For information on what you may reasonably hope for in the way of threading support in STL containers, consult [Item 12](#).)

In a design based on reference counting, everything outside the string object can be shared by multiple strings (if they have the same value), so something else we can observe from the diagrams is that Implementation A offers less sharing than B or C. In particular, Implementations B and C can share a string's size and capacity, thus potentially reducing the amortized per-object cost of storing that data. Interestingly, the fact that Implementation C fails to support per-object allocators means that it is the only implementation that can share allocators: all strings must use the same one! (Again, for details on the rules governing allocators, turn to [Item 10](#).) Implementation D shares no data across string objects.

An interesting aspect of string behavior you can't fully deduce from the diagrams is the policy regarding memory allocation for small strings. Some implementations refuse to allocate memory for fewer than a certain number of characters, and Implementations A and D are among them. Look again at this statement:

```
string s("Perse");           // s is a string of size 5
```

Implementation A has a minimum allocation size of 32 characters, so though *s*'s size is 5 under all implementations, its capacity under Implementation A is 31. (The 32nd character is presumably reserved for a trailing null, thus making it easy to implement the `c_str` member function.) Implementation D's minimum buffer size is 16, including room for a trailing null, so under Implementation D, *s*'s capacity is 15. Of course, what distinguishes Implementation D in this area is that the memory for strings with capacity less than 16 is contained within the string object itself. Implementation B has no minimum allocation, and under Implementation B, *s*'s capacity is 7. (Why it's not 6 or 5, I don't know. I didn't read the source code that closely, sorry.)

I hope it's obvious that the various implementations' policies on minimum allocations might be important to you if you expect to have lots of short strings and either (1) your release environment has very little memory or (2) you are concerned about locality of reference and want to cluster strings on as few pages as possible.

Clearly, string implementations have more degrees of freedom than are apparent at first glance, and equally clearly, different implementers have taken advantage of this design flexibility in different ways. Let's summarize the things that vary:

- string values may or may not be reference counted. By default, many implementations do use reference counting, but they usually offer a way to turn it off, often via a preprocessor macro. [Item 13](#) gives an example of specific conditions under which you might want to turn it off, but you might want to do so for other reasons, too. For example, reference counting helps only when strings are frequently copied, and some applications just don't copy strings often enough to justify the overhead.
- string objects may range in size from one to at least seven times the size of `char*` pointers.
- Creation of a new string value may require zero, one, or two dynamic allocations.
- string objects may or may not share information on the string's size and capacity.
- strings may or may not support per-object allocators.
- Different implementations have different policies regarding minimum allocations for character buffers.

Now, don't get me wrong. I think string is one of the most important components of the standard library, and I encourage you to use it as often as you can. [Item 13](#), for example, is devoted to why you should use string in place of dynamically allocated character arrays. At the same time, if you're to make effective use of the STL, you need to be aware of the variability in string implementations, especially if you're writing code that must run on different STL platforms and you face stringent performance requirements.

Besides, string seems so conceptually simple. Who'd have thought the implementations could be so interesting?

Item 16: Know how to pass vector and string data to legacy APIs.

Since C++ was standardized in 1998, the C++ elite haven't been terribly subtle in their attempt to nudge programmers away from arrays and towards vectors. They've been similarly overt in trying to get developers to shift from `char*` pointers to string objects. There are good reasons for making these changes, including the elimination of common programming errors (see [Item 13](#)) and the ability to take full advantage of the power of the STL algorithms (see, e.g., [Item 31](#)).

Still, obstacles remain, and one of the most common is the existence of legacy C APIs that traffic in arrays and `char*` pointers instead of vector and string objects. Such APIs will exist for a long time, so we must make peace with them if we are to use the STL effectively.

Fortunately, it's easy. If you have a vector `v` and you need to get a pointer to the data in `v` that can be viewed as an array, just use `&v[0]`. For a string `s`, the corresponding incantation is simply `s.c_str()`. But read on. As the fine print in advertising often points out, certain restrictions apply.

Given

```
vector<int> v;
```

the expression `v[0]` yields a reference to the first element in the vector, so `&v[0]` is a pointer to that first element. The elements in a vector are constrained by the C++ Standard to be stored in contiguous memory, just like an array, so if we wish to pass `v` to a C API that looks something like this,

```
void doSomething(const int* plnts, size_t numInts);
```

we can do it like this:

```
doSomething(&v[0], v.size());
```

Maybe. Probably. The only sticking point is if `v` is empty. If it is, `v.size()` is zero, and `&v[0]` attempts to produce a pointer to something that does not exist. Not good. Undefined results. A safer way to code the call is this:

```
if (!v.empty()) {  
    doSomething(&v[0], v.size());  
}
```

If you travel in the wrong circles, you may run across shady characters who will tell you that you can use `v.begin()` in place of `&v[0]`, because (these loathsome creatures will tell you) `begin` returns an iter-

ator into the vector, and for vectors, iterators are really pointers. That's often true, but as [Item 50](#) reveals, it's not always true, and you should never rely on it. The return type of `begin` is an iterator, not a pointer, and you should never use `begin` when you need to get a pointer to the data in a vector. If you're determined to type `v.begin()` for some reason, type `&*v.begin()`, because that will yield the same pointer as `&v[0]`, though it's more work for you as a typist and more obscure for people trying to make sense of your code. Frankly, if you're hanging out with people who tell you to use `v.begin()` instead of `&v[0]`, you need to rethink your social circle.

The approach to getting a pointer to container data that works for vectors isn't reliable for strings, because (1) the data for strings are not guaranteed to be stored in a single chunk of contiguous memory, and (2) the internal representation of a string is not guaranteed to end with a null character. This explains the existence of the string member function `c_str`, which returns a pointer to the value of the string in a form designed for C. We can thus pass a string `s` to this function,

```
void doSomething(const char *pString);
```

like this:

```
doSomething(s.c_str());
```

This works even if the string is of length zero. In that case, `c_str` will return a pointer to a null character. It also works if the string has embedded nulls. If it does, however, `doSomething` is likely to interpret the first embedded null as the end of the string. String objects don't care if they contain null characters, but `char*`-based C APIs do.

Look again at the `doSomething` declarations:

```
void doSomething(const int* plnts, size_t numlnts);
```

```
void doSomething(const char *pString);
```

In both cases, the pointers being passed are pointers to `const`. The vector or string data are being passed to an API that will *read* it, not modify it. This is by far the safest thing to do. For strings, it's the only thing to do, because there is no guarantee that `c_str` yields a pointer to the internal representation of the string data; it could return a pointer to an unmodifiable *copy* of the string's data, one that's correctly formatted for a C API. (If this makes the efficiency hairs on the back of your neck rise up in alarm, rest assured that the alarm is probably false. I don't know of any contemporary library implementation that takes advantage of this latitude.)

For a vector, you have a little more flexibility. If you pass `v` to a C API that modifies `v`'s elements, that's typically okay, but the called routine

must not attempt to change the number of elements in the vector. For example, it must not try to “create” new elements in a vector’s unused capacity. If it does, `v` will become internally inconsistent, because it won’t know its correct size any longer. `v.size()` will yield incorrect results. And if the called routine attempts to add data to a vector whose size and capacity (see [Item 14](#)) are the same, truly horrible things could happen. I don’t even want to contemplate them. They’re just too awful.

Did you notice my use of the word “typically” in the phrase “that’s typically okay” in the preceding paragraph? Of course you did. Some vectors have extra constraints on their data, and if you pass a vector to an API that modifies the vector’s data, you must ensure that the additional constraints continue to be satisfied. For example, [Item 23](#) explains how sorted vectors can often be a viable alternative to associative containers, but it’s important for such vectors to remain sorted. If you pass a sorted vector to an API that may modify the vector’s data, you’ll need to take into account that the vector may no longer be sorted after the call has returned.

If you have a vector that you’d like to initialize with elements from a C API, you can take advantage of the underlying layout compatibility of vectors and arrays by passing to the API the storage for the vector’s elements:

```
// C API: this function takes a pointer to an array of at most arraySize
// doubles and writes data to it. It returns the number of doubles written,
// which is never more than arraySize.
size_t fillArray(double *pArray, size_t arraySize);

vector<double> vd(maxNumDoubles);           // create a vector whose
                                           // size is maxNumDoubles

vd.resize(fillArray(&vd[0], vd.size()));    // have fillArray write data
                                           // into vd, then resize vd
                                           // to the number of
                                           // elements fillArray wrote
```

This technique works only for vectors, because only vectors are guaranteed to have the same underlying memory layout as arrays. If you want to initialize a string with data from a C API, however, you can do it easily enough. Just have the API put the data into a `vector<char>`, then copy the data from the vector to the string:

```
// C API: this function takes a pointer to an array of at most arraySize
// chars and writes data to it. It returns the number of chars written,
// which is never more than arraySize.
size_t fillString(char *pArray, size_t arraySize);
```

```

vector<char> vc(maxNumChars);           // create a vector whose
                                         // size is maxNumChars

size_t charsWritten = fillString(&vc[0], vc.size()); // have fillString write
                                                         // into vc

string s(vc.begin(), vc.begin()+charsWritten); // copy data from vc to s
                                                         // via range constructor
                                                         // (see Item 5)

```

In fact, the idea of having a C API put data into a vector and then copying the data into the STL container you really want it in always works:

```

size_t fillArray(double *pArray, size_t arraySize); // as above

vector<double> vd(maxNumDoubles); // also as above
vd.resize(fillArray(&vd[0], vd.size()));

deque<double> d(vd.begin(), vd.end()); // copy data into
                                         // deque

list<double> l(vd.begin(), vd.end()); // copy data into list

set<double> s(vd.begin(), vd.end()); // copy data into set

```

Furthermore, this hints at how STL containers other than vector or string can pass their data to C APIs. Just copy each container's data into a vector, then pass it to the API:

```

void doSomething(const int* plnts, size_t numInts); // C API (from above)

set<int> intSet; // set that will hold
... // data to pass to API

vector<int> v(intSet.begin(), intSet.end()); // copy set data into
                                              // a vector

if (!v.empty()) doSomething(&v[0], v.size()); // pass the data to
                                              // the API

```

You could copy the data into an array, too, then pass the array to the C API, but why would you want to? Unless you know the size of the container during compilation, you'd have to allocate the array dynamically, and [Item 13](#) explains why you should prefer vectors to dynamically allocated arrays anyway.

Item 17: Use “the swap trick” to trim excess capacity.

So you're writing software to support the TV game show *Give Me Lots Of Money — Now!*, and you're keeping track of the potential contestants, whom you have stored in a vector:

```

class Contestant { ... };

vector<Contestant> contestants;

```

When the show puts out a call for new contestants, it's inundated with applicants, and your vector quickly acquires a lot of elements. As the show's producers vet the prospective players, however, a relatively small number of viable candidates get moved to the front of the vector (perhaps via `partial_sort` or `partition` — see [Item 31](#)), and the candidates no longer in the running are removed from the vector (typically by calling a range form of `erase` — see [Item 5](#)). This does a fine job of reducing the size of the vector, but it does nothing to reduce its capacity. If your vector held a hundred thousand potential candidates at some point, its capacity would continue to be at least 100,000, even if later it held only, say, 10.

To avoid having your vector hold onto memory it no longer needs, you'd like to have a way to reduce its capacity from the maximum it used to the amount it currently needs. Such a reduction in capacity is commonly known as “shrink to fit.” Shrink-to-fit is easy to program, but the code is — how shall I put this? — something less than intuitive. Let me show it to you, then I'll explain how it works.

This is how you trim the excess capacity from your contestants vector:

```
vector<Contestant>(contestants.begin(), contestants.end()).swap(contestants);
```

The expression `vector<Contestant>(contestants.begin(), contestants.end())` creates a temporary vector that is a copy of `contestants`; `vector`'s range constructor (see [Item 5](#)) does the work. The resulting temporary vector has no excess capacity. We then swap the data in the temporary vector with that in `contestants`. By the time we're done, `contestants` has the trimmed capacity of the temporary, and the temporary holds the bloated capacity that used to be in `contestants`. At that point (the end of the statement), the temporary vector is destroyed, thus freeing the memory formerly used by `contestants`. *Voilà!* Shrink-to-fit.

The same trick is applicable to strings:

```
string s;
...                               // make s large, then erase most
                                   // of its characters
string(s.begin(), s.end()).swap(s); // do a “shrink-to-fit” on s
```

Now, the language police require that I inform you that there's no guarantee that this technique will truly eliminate excess capacity. Implementers are free to give vectors and strings excess capacity if they want to, and sometimes they want to. For example, they may have a minimum capacity below which they never go, or they may constrain a vector's or string's capacity to be a power of two. (In my experience,

such anomalies are more common in string implementations than in vector implementations. For examples, see [Item 15](#).) This approach to “shrink-to-fit,” then, doesn’t really mean “make the capacity as small as possible,” it means “make the capacity as small as this implementation is willing to make it given the current size of the container.” Short of switching to a different implementation of the STL, however, this is the best you can do, so when you think “shrink-to-fit” for vectors and strings, think “the swap trick.”

As an aside, a variant of the swap trick can be used both to clear a container and to reduce its capacity to the minimum your implementation offers. You simply do the swap with a temporary vector or string that is default-constructed:

```
vector<Contestant> v;
string s;

...                                // use v and s
vector<Contestant>().swap(v);      // clear v and minimize its capacity
string().swap(s);                 // clear s and minimize its capacity
```

Item 18: Avoid using vector<bool>.

As an STL container, there are really only two things wrong with vector<bool>. First, it’s not an STL container. Second, it doesn’t hold bools. Other than that, there’s not much to object to.

An object doesn’t become an STL container just because somebody says it’s one. An object becomes an STL container only if it fulfills all the container requirements laid down in section 23.1 of the Standard for C++. Among the requirements is that if c is a container of objects of type T and c supports operator[], the following must compile:

```
T *p = &c[0];                      // initialize a T* with the address
                                   // of whatever operator[] returns
```

In other words, if you use operator[] to get at one of the T objects in a Container<T>, you can get a pointer to that object by taking its address. (This assumes that T hasn’t perversely overloaded operator&.) If vector<bool> is to be a container, then, this code must compile:

```
vector<bool> v;
bool *pb = &v[0];                  // initialize a bool* with the address of
                                   // what vector<bool>::operator[] returns
```

But it won't compile. It won't, because `vector<bool>` is a pseudo-container that contains not actual bools, but a packed representation of bools that is designed to save space. In a typical implementation, each "bool" stored in the "vector" occupies a single bit, and an eight-bit byte holds eight "bools." Internally, `vector<bool>` uses the moral equivalent of bitfields to represent the bools it pretends to be holding.

Like bools, bitfields represent only two possible values, but there is an important difference between true bools and bitfields masquerading as bools. You may create a pointer to a real bool, but pointers to individual bits are forbidden.

References to individual bits are forbidden, too, and this posed a problem for the design of `vector<bool>`'s interface, because the return type of `vector<T>::operator[]` is supposed to be `T&`. That wouldn't be a problem if `vector<bool>` really held bools, but because it doesn't, `vector<bool>::operator[]` would need somehow to return a reference to a bit, and there's no such thing.

To deal with this difficulty, `vector<bool>::operator[]` returns an object that *acts like* a reference to a bit, a so-called *proxy object*. (You don't need to understand proxy objects to use the STL, but they're a C++ technique worth knowing about. For information about them, consult Item 30 of *More Effective C++* as well as the "Proxy" chapter in Gamma *et al.*'s *Design Patterns* [6].) Stripped down to the bare essentials, `vector<bool>` looks like this:

```
template <typename Allocator>
vector<bool, Allocator> {
public:
    class reference { ... };           // class to generate proxies for
                                     // references to individual bits

    reference operator[](size_type n); // operator[] returns a proxy

    ...
};
```

Now it's clear why this code won't compile:

```
vector<bool> v;
bool *pb = &v[0];           // error! the expression on the right is
                             // of type vector<bool>::reference*,
                             // not bool*
```

Because it won't compile, `vector<bool>` fails to satisfy the requirements for STL containers. Yes, `vector<bool>` is in the Standard, and yes, it *almost* satisfies the requirements for STL containers, but almost just isn't good enough. The more you write your own templates

designed to work with the STL, the more you'll appreciate that. The day will come, I assure you, when you'll write a template that will work only if taking the address of a container element yields a pointer to the contained type, and when that day comes, you'll suddenly understand the difference between being a container and *almost* being a container.

You may wonder why `vector<bool>` is in the Standard, given that it's not a container and all. The answer has to do with a noble experiment that failed, but let me defer that discussion for a moment while I address a more pressing question. To wit, if `vector<bool>` should be avoided because it's not a container, what should you use when you *need* a `vector<bool>`?

The standard library provides two alternatives that suffice almost all the time. The first is `deque<bool>`. A deque offers almost everything a vector does (the only notable omissions are `reserve` and `capacity`), and a `deque<bool>` is an STL container that really contains bools. Of course, the underlying memory for a deque isn't contiguous, so you can't pass the data behind a `deque<bool>` to a C API[†] that expects an array of bool (see [Item 16](#)), but you couldn't do that with a `vector<bool>` anyway because there's no portable way to get at the data for a `vector<bool>`. (The techniques that work for vectors in [Item 16](#) fail to compile for `vector<bool>`, because they depend on being able to get a pointer to the type of element contained in the vector. Have I mentioned that a `vector<bool>` doesn't contain bools?)

The second alternative to `vector<bool>` is `bitset`. `bitset` isn't an STL container, but it is part of the standard C++ library. Unlike STL containers, its size (number of elements) is fixed during compilation, so there is no support for inserting or erasing elements. Furthermore, because it's not an STL container, it offers no support for iterators. Like `vector<bool>`, however, it uses a compact representation that devotes only a single bit to each value it contains. It offers `vector<bool>`'s special flip member function, as well as a number of other special member functions that make sense for collections of bits. If you can live without iterators and dynamic changes in size, you'll probably find that a `bitset` will serve you well.

And now let me discuss the noble experiment that failed, the one that left behind as its residue the STL non-container `vector<bool>`. I mentioned earlier that proxy objects are often useful in C++ software development. The members of the C++ Standardization Committee

[†] This would presumably be a C99 API, because `bool` was added to C only as of that version of the language.

were well aware of this, so they decided to develop `vector<bool>` as a demonstration of how the STL could support containers whose elements are accessed via proxies. With this example in the Standard, they reasoned, practitioners would have a ready reference for implementing their own proxy-based containers.

Alas, what they discovered was that it was *not* possible to create proxy-based containers that satisfy all the requirements of STL containers. For one reason or another, however, their failed attempt at developing one remained in the Standard. One can speculate on why `vector<bool>` was retained, but practically speaking, it doesn't matter. What does matter is this: `vector<bool>` doesn't satisfy the requirements of an STL container; you're best off not using it; and `deque<bool>` and `bitset` are alternative data structures that will almost certainly satisfy your need for the capabilities promised by `vector<bool>`.

3

Associative Containers

Somewhat like the polychromatic horse in *The Wizard of Oz* movie, the associative containers are creatures of a different color. True, they share many characteristics with the sequence containers, but they differ in a number of fundamental ways. For example, they automatically keep themselves sorted; they view their contents through the lens of equivalence instead of equality; sets and maps reject duplicate entries; and maps and multimaps generally ignore half of each object they contain. Yes, the associative containers are containers, but if you'll excuse my likening vectors and strings to Kansas, we are definitely not in Kansas any more.

In the Items that follow, I explain the critical notion of equivalence, describe an important restriction on comparison functions, motivate custom comparison functions for associative containers of pointers, discuss the official and practical sides of key constness, and offer advice on improving the efficiency of associative containers.

At the end of the chapter, I examine the STL's lack of containers based on hash tables, and I survey two common (nonstandard) implementations. Though the STL proper doesn't offer hash tables, you need not write your own or do without. High-quality implementations are readily available.

Item 19: Understand the difference between equality and equivalence.

The STL is awash in comparisons of objects to see if they have the same value. For example, when you ask `find` to locate the first object in a range with a particular value, `find` has to be able to compare two objects to see if the value of one is the same as the value of the other. Similarly, when you attempt to insert a new element into a set, `set::insert` has to be able to determine whether that element's value is already in the set.

The `find` algorithm and `set`'s `insert` member function are representative of many functions that must determine whether two values are the same. Yet they do it in different ways. `find`'s definition of "the same" is *equality*, which is based on `operator==`. `set::insert`'s definition of "the same" is *equivalence*, which is usually based on `operator<`. Because these are different definitions, it's possible for one definition to dictate that two objects have the same value while the other definition decrees that they do not. As a result, you must understand the difference between equality and equivalence if you are to make effective use of the STL.

Operationally, the notion of equality is based on `operator==`. If the expression "`x == y`" returns true, `x` and `y` have equal values, otherwise they don't. That's pretty straightforward, though it's useful to bear in mind that just because `x` and `y` have equal values does not necessarily imply that all of their data members have equal values. For example, we might have a `Widget` class that internally keeps track of its last time of access,

```
class Widget {
public:
    ...
private:
    TimeStamp lastAccessed;
    ...
};
```

and we might have an `operator==` for `Widgets` that ignores this field:

```
bool operator==(const Widget& lhs, const Widget& rhs)
{
    // code that ignores the lastAccessed field
}
```

In that case, two `Widgets` would have equal values even if their `lastAccessed` fields were different.

Equivalence is based on the relative ordering of object values in a sorted range. Equivalence makes the most sense if you think about it in terms of the sort order that is part of every standard associative container (i.e., `set`, `multiset`, `map`, and `multimap`). Two objects `x` and `y` have equivalent values with respect to the sort order used by an associative container `c` if neither precedes the other in `c`'s sort order. That sounds complicated, but in practice, it's not. Consider, as an example, a `set<Widget>` `s`. Two `Widgets` `w1` and `w2` have equivalent values with respect to `s` if neither precedes the other in `s`'s sort order. The default comparison function for a `set<Widget>` is `less<Widget>`, and by default `less<Widget>` simply calls `operator<` for `Widgets`, so `w1` and `w2` have equivalent values with respect to `operator<` if the following expression is true:

```

!(w1 < w2)           // it's not true that w1 < w2
&&                 // and
!(w2 < w1)           // it's not true that w2 < w1

```

This makes sense: two values are equivalent (with respect to some ordering criterion) if neither precedes the other (according to that criterion).

In the general case, the comparison function for an associative container isn't `operator<` or even less, it's a user-defined predicate. (See [Item 39](#) for more information on predicates.) Every standard associative container makes its sorting predicate available through its `key_comp` member function, so two objects `x` and `y` have equivalent values with respect to an associative container `c`'s sorting criterion if the following evaluates to true:

```

!c.key_comp()(x, y) && !c.key_comp()(y, x) // it's not true that x precedes
                                           // y in c's sort order and it's
                                           // also not true that y precedes
                                           // x in c's sort order

```

The expression `!c.key_comp()(x, y)` looks nasty, but once you understand that `c.key_comp()` is returning a function (or a function object), the nastiness dissipates. `!c.key_comp()(x, y)` is just calling the function (or function object) returned by `key_comp`, and it's passing `x` and `y` as arguments. Then it negates the result. `c.key_comp()(x, y)` returns true only when `x` precedes `y` in `c`'s sort order, so `!c.key_comp()(x, y)` is true only when `x` *doesn't* precede `y` in `c`'s sort order.

To fully grasp the implications of equality versus equivalence, consider a case-insensitive `set<string>`, i.e., a `set<string>` where the set's comparison function ignores the case of the characters in the strings. Such a comparison function would consider “STL” and “stL” to be equivalent. [Item 35](#) shows how to implement a function, `ciStringCompare`, that performs a case-insensitive comparison, but `set` wants a comparison function *type*, not an actual function. To bridge this gap, we write a functor class whose `operator()` calls `ciStringCompare`:

```

struct CIStringCompare:           // class for case-insensitive
public                             // string comparisons;
    binary_function<string, string, bool> { // see Item 40 for info on
                                           // this base class

    bool operator()(const string& lhs,
                     const string& rhs) const
    {
        return ciStringCompare(lhs, rhs); // see Item 35 for how
    }                                     // ciStringCompare is
                                           // implemented
};

```

Given `CIStringCompare`, it's easy to create a case-insensitive `set<string>`:

```

set<string, CIStringCompare> ciss; // ciss = "case-insensitive
                                   // string set"

```

If we insert “Persephone” and “persephone” into the set, only the first string is added, because the second one is equivalent to the first:

```
ciss.insert("Persephone");           // a new element is added to the set
ciss.insert("persephone");           // no new element is added to the set
```

If we now search for the string “persephone” using set’s find member function, the search will succeed,

```
if (ciss.find("persephone") != ciss.end()) ...      // this test will succeed
```

but if we use the non-member find algorithm, the search will fail:

```
if (find(ciss.begin(), ciss.end(),
         "persephone") != ciss.end()) ...           // this test will fail
```

That’s because “persephone” is *equivalent* to “Persephone” (with respect to the comparison functor `CIStringCompare`), but it’s not *equal* to it (because `string("persephone") != string("Persephone")`). This example demonstrates one reason why you should follow the advice in [Item 44](#) and prefer member functions (like `set::find`) to their non-member counterparts (like `find`).

You might wonder why the standard associative containers are based on equivalence instead of equality. After all, most programmers have an intuition about equality that they lack for equivalence. (Were that not the case, there’d be no need for this Item.) The answer is simple at first glance, but the more closely you look, the more subtle it becomes.

The standard associative containers are kept in sorted order, so each container must have a comparison function (less, by default) that defines how to keep things sorted. Equivalence is defined in terms of this comparison function, so clients of a standard associative container need to specify only one comparison function (the one determining the sort order) for any container they use. If the associative containers used equality to determine when two objects have the same value, each associative container would need, in addition to its comparison function for sorting, a second comparison function for determining when two values are equal. (By default, this comparison function would presumably be `equal_to`, but it’s interesting to note that `equal_to` is never used as a default comparison function in the STL. When equality is needed in the STL, the convention is to simply call `operator==` directly. For example, this is what the non-member `find` algorithm does.)

Let’s suppose we had a set-like STL container called `set2CF`, “set with two comparison functions.” The first comparison function would be used to determine the sort order of the set, the second would be used

to determine whether two objects had the same value. Now consider this `set2CF`:

```
set2CF<string, CStringCompare, equal_to<string> > s;
```

Here, `s` sorts its strings internally without regard to case, and the equality criterion is the intuitive one: two strings have the same value if they are equal to one another. Let's insert our two spellings of Hades' reluctant bride (Persephone) into `s`:

```
s.insert("Persephone");  
s.insert("persephone");
```

What should this do? If we observe that `"Persephone" != "persephone"` and insert them both into `s`, what order should they go in? Remember that the sorting function can't tell them apart. Do we insert them in some arbitrary order, thus giving up the ability to traverse the contents of the set in deterministic order? (This inability to traverse associative container elements in a deterministic order already afflicts multisets and multimaps, because the Standard places no constraints on the relative ordering of equivalent values (for multisets) or keys (for multimaps).) Or do we insist on a deterministic ordering of `s`'s contents and ignore the second attempted insertion (the one for `"persephone"`)? If we do that, what happens here?

```
if (s.find("persephone") != s.end()) ...           // does this test succeed or fail?
```

Presumably `find` employs the equality check, but if we ignored the second call to insert in order to maintain a deterministic ordering of the elements in `s`, this `find` will fail, even though the insertion of `"persephone"` was ignored on the basis of its being a duplicate value!

The long and short of it is that by using only a single comparison function and by employing equivalence as the arbiter of what it means for two values to be "the same," the standard associative containers sidestep a whole host of difficulties that would arise if two comparison functions were allowed. Their behavior may seem a little odd at first (especially when one realizes that `member` and `non-member find` may return different results), but in the long run, it avoids the kind of confusion that would arise from mixing uses of equality and equivalence within standard associative containers.

Interestingly, once you leave the realm of *sorted* associative containers, the situation changes, and the issue of equality versus equivalence can be — and has been — revisited. There are two common designs for nonstandard (but widely available) associative containers based on hash tables. One design is based on equality, while the other is based on equivalence. I encourage you to turn to [Item 25](#) to learn

more about these containers and the design decisions they've adopted.

Item 20: Specify comparison types for associative containers of pointers.

Suppose you have a set of `string*` pointers and you insert the names of some animals into the set:

```
set<string*> ssp;                               // ssp = "set of string ptrs"
ssp.insert(new string("Anteater"));
ssp.insert(new string("Wombat"));
ssp.insert(new string("Lemur"));
ssp.insert(new string("Penguin"));
```

You then write the following code to print the contents of the set, expecting the strings to come out in alphabetical order. After all, sets keep their contents sorted.

```
for (set<string*>::const_iterator i = ssp.begin();    // you expect to see
     i != ssp.end();                                // this: "Anteater",
     ++i)                                           // "Lemur", "Penguin",
    cout << *i << endl;                            // "Wombat"
```

The comment describes what you expect to see, but you don't see that at all. Instead, you see four hexadecimal numbers. They are pointer values. Because the set holds pointers, `*i` isn't a string, it's a *pointer* to a string. Let this be a lesson that reminds you to adhere to the guidance of [Item 43](#) and avoid writing your own loops. If you'd used a call to the copy algorithm instead,

```
copy(ssp.begin(), ssp.end(),                      // copy the strings in
     ostream_iterator<string>(cout, "\n"));        // ssp to cout (but this
                                                    // won't compile)
```

you'd not only have typed fewer characters, you'd have found out about your error sooner, because this call to `copy` won't compile. `ostream_iterator` insists on knowing the type of object being printed, so when you tell it it's a string (by passing that as the template parameter), your compilers detect the mismatch between that and the type of objects stored in `ssp` (which is `string*`), and they refuse to compile the code. Score another one for type safety.

If you exasperatedly change the `*i` in your explicit loop to `**i`, you *might* get the output you want, but you probably won't. Yes, the animal names will be printed, but the chances of their coming out in alpha-

betical order are only 1 in 24. `ssp` keeps its contents in sorted order, but it holds pointers, so it sorts by *pointer value*, not by string value. There are 24 possible permutations for four pointer values, so there are 24 possible orders for the pointers to be stored in. Hence the odds of 1 in 24 of your seeing the strings in alphabetical order.[†]

To surmount this problem, it helps to recall that

```
set<string*> ssp;
```

is shorthand for this:

```
set<string*, less<string*> > ssp;
```

Well, to be completely accurate, it's shorthand for

```
set<string*, less<string*>, allocator<string*> > ssp;
```

but allocators don't concern us in this Item, so we'll ignore them.

If you want the `string*` pointers to be stored in the set in an order determined by the string values, you can't use the default comparison functor class `less<string*>`. You must instead write your own comparison functor class, one whose objects take `string*` pointers and order them by the values of the strings they point to. Like this:

```
struct StringPtrLess:
    public binary_function<const string*,           // see Item 40 for the
                        const string*,             // reason for this base
                        bool> {                   // class

    bool operator()(const string *ps1, const string *ps2) const
    {
        return *ps1 < *ps2;
    }
};
```

Then you can use `StringPtrLess` as `ssp`'s comparison type:

```
typedef set<string*, StringPtrLess> StringPtrSet;
StringPtrSet ssp;                                // create a set of strings and order
                                                // them as defined by StringPtrLess
...                                              // insert the same four strings as
                                                // before
```

Now your loop will finally do what you want it to do (provided you've fixed the earlier problem whereby you used `*i` instead of `**i`):

[†] Practically speaking, the 24 permutations are not equally likely, so the "1 in 24" statement is a bit misleading. Still, there are 24 different orders, and you *could* get any one of them.

```

for (StringPtrSet::const_iterator i = ssp.begin();    // prints "Anteater";
     i != ssp.end();                                // "Lemur";
     ++i)                                           // "Penguin";
    cout << *i << endl;                           // "Wombat"

```

If you want to use an algorithm instead, you could write a function that knows how to dereference `string*` pointers before printing them, then use that function in conjunction with `for_each`:

```

void print(const string *ps)                        // print to cout the
{                                                    // object pointed to
    cout << *ps << endl;                           // by ps
}

for_each(ssp.begin(), ssp.end(), print);            // invoke print on each
                                                    // element in ssp

```

Or you could get fancy and write a generic dereferencing functor class, then use that with `transform` and an `ostream_iterator`:

```

// when functors of this type are passed a T*, they return a const T&
struct Dereference {
    template <typename T>
    const T& operator()(const T *ptr) const
    {
        return *ptr;
    }
};

transform(ssp.begin(), ssp.end(),                    // "transform" each
          ostream_iterator<string>(cout, "\n"),      // element in ssp by
          Dereference());                            // dereferencing it,
                                                    // writing the results
                                                    // to cout

```

Replacement of loops with algorithms, however, is not the point, at least not for this Item. (It is the point for [Item 43](#).) The point is that anytime you create a standard associative container of pointers, you must bear in mind that the container will be sorted by the values of the pointers. Only rarely will this be what you want, so you'll almost always want to create your own functor class to serve as a comparison type.

Notice that I wrote “comparison *type*.” You may wonder why you have to go to the trouble of creating a functor class instead of simply writing a comparison function for the set. For example, you might think to try this:

```

bool stringPtrLess(const string* ps1,    // would-be comparison
                  const string* ps2)    // function for string*
{
    return *ps1 < *ps2;                // pointers to be sorted by
}                                       // string value

set<string*, stringPtrLess> ssp;        // attempt to use stringPtrLess
                                       // as ssp's comparison function;
                                       // this won't compile

```

The problem here is that each of the set template's three parameters is a *type*. Unfortunately, `stringPtrLess` isn't a type, it's a function. That's why the attempt to use `stringPtrLess` as set's comparison function won't compile. set doesn't want a function, it wants a type that it can internally instantiate to *create* a function.

Anytime you create associative containers of pointers, figure you're probably going to have to specify the container's comparison type, too. Most of the time, your comparison type will just dereference the pointers and compare the pointed-to objects (as is done in `StringPtrLess` above). That being the case, you might as well keep a template for such comparison functors close at hand. Like this:

```

struct DereferenceLess {               // see Item 7 for why we're using
                                       // a class containing a templated
                                       // operator()

    template <typename PtrType>
    bool operator()(PtrType pT1,       // parameters are passed by
                    PtrType pT2) const // value, because we expect them
    {                                  // to be (or to act like) pointers
        return *pT1 < *pT2;
    }

};

```

Such a template eliminates the need to write classes like `StringPtrLess`, because we can use `DereferenceLess` instead:

```

set<string*, DereferenceLess> ssp;    // behaves the same as
                                       // set<string*, StringPtrLess>

```

Oh, one more thing. This Item is about associative containers of pointers, but it applies equally well to containers of objects that *act* like pointers, e.g., smart pointers and iterators. If you have an associative container of smart pointers or of iterators, plan on specifying the comparison type for it, too. Fortunately, the solution for pointers tends to work for pointeresque objects, too. Just as `DereferenceLess` is likely to be suitable as the comparison type for an associative container of `T*`, it's likely to work as the comparison type for containers of iterators and smart pointers to `T` objects, too.

Item 21: Always have comparison functions return false for equal values.

Let me show you something kind of cool. Create a set where `less_equal` is the comparison type, then insert 10 into the set:

```
set<int, less_equal<int> > s;           // s is sorted by "<="
s.insert(10);                           // insert the value 10
```

Now try inserting 10 again:

```
s.insert(10);
```

For this call to insert, the set has to figure out whether 10 is already present. We know that it is, but the set is dumb as toast, so it has to check. To make it easier to understand what happens when the set does this, we'll call the 10 that was initially inserted 10_A and the 10 that we're trying to insert 10_B .

The set runs through its internal data structures looking for the place to insert 10_B . It ultimately has to check 10_B to see if it's the same as 10_A . The definition of "the same" for associative containers is equivalence (see [Item 19](#)), so the set tests to see whether 10_B is equivalent to 10_A . When performing this test, it naturally uses the set's comparison function. In this example, that's `operator<=`, because we specified `less_equal` as the set's comparison function, and `less_equal` means `operator<=`. The set thus checks to see whether this expression is true:

```
!(10A <= 10B) && !(10B <= 10A)           // test 10A and 10B for equivalence
```

Well, 10_A and 10_B are both 10, so it's clearly true that $10_A <= 10_B$. Equally clearly, $10_B <= 10_A$. The above expression thus simplifies to

```
!(true) && !(true)
```

and that simplifies to

```
false && false
```

which is simply false. That is, the set concludes that 10_A and 10_B are *not* equivalent, hence *not* the same, and it thus goes about inserting 10_B into the container alongside 10_A . Technically, this action yields undefined behavior, but the nearly universal outcome is that the set ends up with *two* copies of the value 10, and that means it's not a set any longer. By using `less_equal` as our comparison type, we've corrupted the container! Furthermore, *any* comparison function where equal values return true will do the same thing. Isn't that cool?

Okay, maybe your definition of cool isn't the same as mine. Even so, you'll still want to make sure that the comparison functions you use for associative containers always return false for equal values. You'll need to be vigilant, however. It's surprisingly easy to run afoul of this constraint.

For example, [Item 20](#) describes how to write a comparison function for containers of `string*` pointers such that the container sorts its contents by the values of the strings instead of the values of the pointers. That comparison function sorts them in ascending order, but let's suppose you're in need of a comparison function for a container of `string*` pointers that sorts in descending order. The natural thing to do is to grab the existing code and modify it. If you're not careful, you might come up with this, where I've highlighted the changes to the code in [Item 20](#):

```
struct StringPtrGreater:                // highlights show how
    public binary_function<const string*, // this code was changed
                                const string*, // from page 89. Beware,
                                bool> {      // this code is flawed!

    bool operator()(const string *ps1, const string *ps2) const
    {
        return !(*ps1 < *ps2);           // just negate the old test;
    }                                     // this is incorrect!

};
```

The idea here is to reverse the sort order by negating the test inside the comparison function. Unfortunately, negating “<” doesn't give you “>” (which is what you want), it gives you “>=”. And you now understand that “>=”, because it will return true for equal values, is an invalid comparison function for associative containers.

The comparison type you really want is this one:

```
struct StringPtrGreater:                // this is a valid
    public binary_function<const string*, // comparison type for
                                const string*, // associative containers
                                bool> {

    bool operator()(const string *ps1, const string *ps2) const
    {
        return *ps2 < *ps1;             // return whether *ps2
    }                                     // precedes *ps1 (i.e., swap
                                         // the order of the
};                                       // operands)
```

To avoid falling into this trap, all you need to remember is that the return value of a comparison function indicates whether one value precedes another in the sort order defined by that function. Equal val-

ues never precede one another, so comparison functions should always return false for equal values.

Sigh.

I know what you're thinking. You're thinking, "Sure, that makes sense for set and map, because those containers can't hold duplicates. But what about multiset and multimap? Those containers may contain duplicates, so what do I care if the container thinks that two objects of equal value aren't equivalent? It will store them both, which is what the multi containers are supposed to do. No problem, right?"

Wrong. To see why, let's go back to the original example, but this time we'll use a multiset:

```
multiset<int, less_equal<int> > s;           // s is still sorted by "<="
s.insert(10);                               // insert 10A
s.insert(10);                               // insert 10B
```

s now has two copies of 10 in it, so we'd expect that if we do an equal_range on it, we'll get back a pair of iterators that define a range containing both copies. But that's not possible. equal_range, its name notwithstanding, doesn't identify a range of equal values, it identifies a range of *equivalent* values. In this example, s's comparison function says that 10_A and 10_B are not equivalent, so there's no way that both can be in the range identified by equal_range.

You see? Unless your comparison functions always return false for equal values, you break all standard associative containers, regardless of whether they are allowed to store duplicates.

Technically speaking, comparison functions used to sort associative containers must define a "strict weak ordering" over the objects they compare. (Comparison functions passed to algorithms like sort (see [Item 31](#)) are similarly constrained.) If you're interested in the details of what it means to be a strict weak ordering, you can find them in many comprehensive STL references, including Josuttis' *The C++ Standard Library* [3], Austern's *Generic Programming and the STL* [4], and the SGI STL Web Site [21]. I've never found the details terribly illuminating, but one of the requirements of a strict weak ordering bears directly on this Item. That requirement is that any function defining a strict weak ordering must return false if it's passed two copies of the same value.

Hey! That is this Item!

Item 22: Avoid in-place key modification in set and multiset.

The motivation for this Item is easy to understand. Like all the standard associative containers, set and multiset keep their elements in sorted order, and the proper behavior of these containers is dependent on their remaining sorted. If you change the value of an element in an associative container (e.g., change a 10 to a 1000), the new value might not be in the correct location, and that would break the sortedness of the container. Simple, right?

It's especially simple for map and multimap, because programs that attempt to change the value of a key in these containers won't compile:

```
map<int, string> m;
...
m.begin()->first = 10;           // error! map keys can't be changed
multimap<int, string> mm;
...
mm.begin()->first = 20;          // error! multimap keys can't
                                // be changed, either
```

That's because the elements in an object of type `map<K, V>` or `multimap<K, V>` are of type `pair<const K, V>`. Because the type of the key is `const K`, it can't be changed. (Well, you can probably change it if you employ a `const_cast`, as we'll see below. Believe it or not, sometimes that's what you *want* to do.)

But notice that the title of this Item doesn't mention map or multimap. There's a reason for that. As the example above demonstrates, in-place key modification is impossible for map and multimap (unless you use a cast), but it may be possible for set and multiset. For objects of type `set<T>` or `multiset<T>`, the type of the elements stored in the container is simply `T`, not `const T`. Hence, the elements in a set or multiset may be changed anytime you want to. No cast is required. (Actually, things aren't quite that straightforward, but we'll come to that presently. There's no reason to get ahead of ourselves. First we crawl. Later we crawl on broken glass.)

Let us begin with an understanding of why the elements in a set or multiset aren't `const`. Suppose we have a class for employees:

```
class Employee {
public:
    ...
    const string& name() const;           // get employee name
    void setName(const string& name);     // set employee name
    const string& title() const;         // get employee title
    void setTitle(const string& title);   // set employee title
    int idNumber() const;                // get employee ID number
    ...
};
```

As you can see, there are various bits of information we can get about employees. Let us make the reasonable assumption, however, that each employee has a unique ID number, a number that is returned by the `idNumber` function. To create a set of employees, then, it could easily make sense to order the set by ID number only, like this:

```
struct IDNumberLess:
    public binary_function<Employee, Employee, bool> {    // see Item 40
    bool operator()(const Employee& lhs,
                    const Employee& rhs) const
    {
        return lhs.idNumber() < rhs.idNumber();
    }
};

typedef set<Employee, IDNumberLess> EmpIDSet;
EmpIDSet se;                                           // se is a set of employees
                                                    // sorted by ID number
```

Practically speaking, the employee ID number is the *key* for the elements in this set. The rest of the employee data is just along for the ride. That being the case, there's no reason why we shouldn't be able to change the title of a particular employee to something interesting. Like so:

```
Employee selectedID;                                // variable to hold employee
                                                    // with selected ID number

...
EmpIDSet::iterator i = se.find(selectedID);
if (i != se.end()) {
    i->setTitle("Corporate Deity");                // give employee new title;
                                                    // this shouldn't compile —
                                                    // see next page for why
```

Because all we're doing here is changing an aspect of an employee that is unrelated to the way the set is sorted (a *non-key* part of the employee), this code can't corrupt the set. That's why it's reasonable. But such reasonable code precludes having the elements of a set/multiset be `const`. And that's why they aren't.

Why, you might wonder, doesn't the same logic apply to the keys in maps and multimaps? Wouldn't it make sense to create a map from employees to, say, the country in which they live, a map where the comparison function was `IDNumberLess`, just as in the previous example? And given such a map, wouldn't it be reasonable to change an employee's title without affecting the employee's ID number, just as in the previous example?

To be brutally frank, I think it would. Being equally frank, however, it doesn't matter what I think. What matters is what the Standardization Committee thought, and what it thought is that map/multimap keys should be `const` and set/multiset values shouldn't be.

Because the values in a set or multiset are not `const`, attempts to change them may compile. The purpose of this Item is to alert you that if you do change an element in a set or multiset, you must be sure not to change a *key part* — a part of the element that affects the sort-*edness* of the container. If you do, you may corrupt the container, using the container will yield undefined results, and it will be your fault. On the other hand, this constraint applies only to the key parts of the contained objects. For all other parts of the contained elements, it's open season; change away!

Except for the broken glass. Remember the broken glass I referred to earlier? We're there now. Grab some bandages and follow me.

Even if set and multiset elements aren't `const`, there are ways for implementations to keep them from being modified. For example, an implementation could have operator* for a `set<T>::iterator` return a `const T&`. That is, it could have the result of dereferencing a set iterator be a reference-to-*const* element of the set. Under such an implementation, there'd be no way to modify set or multiset elements, because all the ways of accessing the elements would add a `const` before letting you at them.

Are such implementations legal? Arguably yes. And arguably no. The Standard is inconsistent in this area, and in accord with Murphy's Law, different implementers have interpreted it in different ways. The result is that it's not uncommon to find STL implementations where this code compiles and others where it does not:

```
EmpIDSet se;                                // as before, se is a set of employees
                                           // sorted by ID number

Employee selectedID;                        // as before, selectedID is a dummy
                                           // employee with the selected ID
                                           // number

...

EmpIDSet::iterator i = se.find(selectedID);
if (i != se.end()) {
    i->setTitle("Corporate Deity");          // some STL implementations
}                                           // accept this line, others reject it
```

Because of the equivocal state of the Standard and the differing interpretations it has engendered, code that attempts to modify elements in a set or multiset isn't portable.[†]

[†] The Standardization Committee has since clarified that elements in a set or map should not be modifiable without a cast. However, versions of the STL implemented prior to this clarification continue to be used, so, practically speaking, the material in this Item continues to apply.

So where do we stand? Encouragingly, things aren't complicated:

- If portability is not a concern, you want to change the value of an element in a set or multiset, and your STL implementation will let you get away with it, go ahead and do it. Just be sure not to change a key part of the element, i.e., a part of the element that could affect the sortedness of the container.
- If you value portability, assume that the elements in sets and multisets cannot be modified, at least not without a cast.

Ah, casts. We've seen that it can be entirely reasonable to change a non-key portion of an element in a set or a multiset, so I feel compelled to show you how to do it. How to do it correctly and portably, that is. It's not hard, but it requires a subtlety too many programmers overlook: you must cast to a *reference*. As an example, look again at the `setTitle` call we just saw that failed to compile under some implementations:

```
EmpIDSet::iterator i = se.find(selectedID);
if (i != se.end()) {
    i->setTitle("Corporate Deity");           // some STL implementations will
                                              // reject this line because *i is const
}
```

To get this to compile and behave correctly, we must cast away the constness of `*i`. Here's the correct way to do it:

```
if (i != se.end()) {
    const_cast<Employee*>(*i).setTitle("Corporate Deity"); // cast away
                                                         // constness
                                                         // of *i
}
```

This takes the object pointed to by `i`, tells your compilers to treat the result of the cast as a reference to a (non-const) `Employee`, then invoke `setTitle` on the reference. Rather than explain why this works, I'll explain why an alternative approach fails to behave the way people often expect it to.

Many people come up with this code,

```
if (i != se.end()) {
    static_cast<Employee*>(*i).setTitle("Corporate Deity"); // cast *i
                                                         // to an
                                                         // Employee
}
```

which is equivalent to the following:

```
if (i != se.end()) {
    ((Employee)(*i)).setTitle("Corporate Deity"); // same as above,
                                                         // but using C
                                                         // cast syntax
}
```

Both of these compile, and because these are equivalent, they're wrong for the same reason. At runtime, they don't modify `*i`! In both cases, the result of the cast is a temporary anonymous object that is a

copy of **i*, and `setTitle` is invoked on the anonymous object, not on **i*! **i* isn't modified, because `setTitle` is never invoked on that object, it's invoked on a *copy* of that object. Both syntactic forms are equivalent to this:

```
if (i != se.end()) {  
    Employee tempCopy(*i);           // copy *i into tempCopy  
    tempCopy.setTitle("Corporate Deity"); // modify tempCopy  
}
```

Now the importance of a cast to a reference should be clear. By casting to a reference, we avoid creating a new object. Instead, the result of the cast is a reference to an *existing* object, the object pointed to by *i*. When we invoke `setTitle` on the object indicated by the reference, we're invoking `setTitle` on **i*, and that's exactly what we want.

What I've just written is fine for sets and multisets, but when we turn to maps and multimaps, the plot thickens. Recall that a `map<K, V>` or a `multimap<K, V>` contains elements of type `pair<const K, V>`. That `const` means that the first component of the pair is *defined* to be `const`, and *that* means that attempts to modify it (even after casting away its constness) are undefined. If you're a stickler for following the rules laid down by the Standard, you'll *never* try to modify an object serving as a key to a map or multimap.

You've surely heard that casts are dangerous, and I hope this book makes clear that I believe you should avoid them whenever you can. To perform a cast is to shuck temporarily the safety of the type system, and the pitfalls we've discussed exemplify what can happen when you leave your safety net behind.

Most casts can be avoided, and that includes the ones we've just considered. If you want to change an element in a set, multiset, map, or multimap in a way that always works and is always safe, do it in five simple steps:

1. Locate the container element you want to change. If you're not sure of the best way to do that, [Item 45](#) offers guidance on how to perform an appropriate search.
2. Make a copy of the element to be modified. In the case of a map or multimap, be sure not to declare the first component of the copy `const`. After all, you want to change it!
3. Modify the copy so it has the value you want to be in the container.

4. Remove the element from the container, typically via a call to `erase` (see [Item 9](#)).
5. Insert the new value into the container. If the location of the new element in the container's sort order is likely to be the same or adjacent to that of the removed element, use the "hint" form of `insert` to improve the efficiency of the insertion from logarithmic-time to amortized constant-time. Use the iterator you got from Step 1 as the hint.

Here's the same tired employee example, this time written in a safe, portable manner:

```

EmpIDSet se;                                // as before, se is a set of employees
                                           // sorted by ID number

Employee selectedID;                        // as before, selectedID is a dummy
                                           // employee with the desired ID number

...

EmpIDSet::iterator i =
    se.find(selectedID);                    // Step 1: find element to change
if (i != se.end()) {
    Employee e(*i);                        // Step 2: copy the element
    e.setTitle("Corporate Deity");         // Step 3: modify the copy
    se.erase(i++);                        // Step 4: remove the element;
                                           // increment the iterator to maintain
                                           // its validity (see Item 9)

    se.insert(i, e);                       // Step 5: insert new value; hint that its
                                           // location is the same as that of the
                                           // original element
}

```

You'll excuse my putting it this way, but the *key* thing to remember is that with `set` and `multiset`, if you perform any in-place modifications of container elements, you are responsible for making sure that the container remains sorted.

Item 23: Consider replacing associative containers with sorted vectors.

Many STL programmers, when faced with the need for a data structure offering fast lookups, immediately think of the standard associative containers, `set`, `multiset`, `map`, and `multimap`. That's fine, as far as it goes, but it doesn't go far enough. If lookup speed is really important, it's almost certainly worthwhile to consider the nonstandard

hashed containers as well (see [Item 25](#)). With suitable hashing functions, hashed containers can be expected to offer constant-time lookups. (With poorly chosen hashing functions or with table sizes that are too small, the performance of hash table lookups may degrade significantly, but this is relatively uncommon in practice.) For many applications, the expected constant-time lookups of hashed containers are preferable to the guaranteed logarithmic-time lookups that are the hallmark of set, map and their multi companions.

Even if logarithmic-time lookup is what you want, the standard associative containers may not be your best bet. Counterintuitively, it's not uncommon for the standard associative containers to offer performance that is inferior to that of the lowly vector. If you want to make effective use of the STL, you need to understand when and how a vector can offer faster lookups than a standard associative container.

The standard associative containers are typically implemented as balanced binary search trees. A balanced binary search tree is a data structure that is optimized for a mixed combination of insertions, erasures, and lookups. That is, it's designed for applications that do some insertions, then some lookups, then maybe some more insertions, then perhaps some erasures, then a few more lookups, then more insertions or erasures, then more lookups, etc. The key characteristic of this sequence of events is that the insertions, erasures, and lookups are all mixed up. In general, there's no way to predict what the next operation on the tree will be.

Many applications use their data structures in a less chaotic manner. Their use of data structures fall into three distinct phases, which can be summarized like this:

1. **Setup.** Create a new data structure by inserting lots of elements into it. During this phase, almost all operations are insertions and erasures. Lookups are rare or nonexistent.
2. **Lookup.** Consult the data structure to find specific pieces of information. During this phase, almost all operations are lookups. Insertions and erasures are rare or nonexistent. There are so many lookups, the performance of this phase makes the performance of the other phases incidental.
3. **Reorganize.** Modify the contents of the data structure, perhaps by erasing all the current data and inserting new data in its place. Behaviorally, this phase is equivalent to phase 1. Once this phase is completed, the application returns to phase 2.

For applications that use their data structures in this way, a vector is likely to offer better performance (in both time and space) than an

associative container. But not just any vector will do. It has to be a *sorted* vector, because only sorted containers work correctly with the lookup algorithms `binary_search`, `lower_bound`, `equal_range`, etc. (see [Item 34](#)). But why should a binary search through a (sorted) vector offer better performance than a binary search through a binary search tree? Because some things are trite but true, and one of them is that size matters. Others are less trite but no less true, and one of those is that locality of reference matters, too.

Consider first the size issue. Suppose we need a container to hold `Widget` objects, and, because lookup speed is important to us, we are considering both an associative container of `Widgets` and a sorted vector<`Widget`>. If we choose an associative container, we'll almost certainly be using a balanced binary tree. Such a tree would be made up of tree nodes, each holding not only a `Widget`, but also a pointer to the node's left child, a pointer to its right child, and (typically) a pointer to its parent. That means that the space overhead for storing a `Widget` in an associative container would be at least three pointers.

In contrast, there is no overhead when we store a `Widget` in a vector; we simply store a `Widget`. The vector itself has overhead, of course, and there may be empty (reserved) space at the end of the vector (see [Item 14](#)), but the per-vector overhead is typically insignificant (usually three machine words, e.g., three pointers or two pointers and an int), and the empty space at the end can be lopped off via “the swap trick” if necessary (see [Item 17](#)). Even if the extra space is not eliminated, it's unimportant for the analysis below, because that memory won't be referenced when doing a lookup.

Assuming our data structures are big enough, they'll be split across multiple memory pages, but the vector will require fewer pages than the associative container. That's because the vector requires no per-`Widget` overhead, while the associative container exacts three pointers per `Widget`. To see why this is important, suppose you're working on a system where a `Widget` is 12 bytes in size, pointers are 4 bytes, and a memory page holds 4096 (4K) bytes. Ignoring the per-container overhead, you can fit 341 `Widgets` on a page when they are stored in a vector, but you can fit at most 170 when they are stored in an associative container. You'll thus use about twice as much memory for the associative container as you would for the vector. If you're working in an environment where virtual memory is available, it's easy to see how that can translate into a lot more page faults, therefore a system that is significantly slower for large sets of data.

I'm actually being optimistic about the associative containers here, because I'm assuming that the nodes in the binary trees are clustered

together on a relatively small set of memory pages. Most STL implementations use custom memory managers (implemented on top of the containers' allocators — see Items 10 and 11) to achieve such clustering, but if your STL implementation fails to take steps to improve locality of reference among tree nodes, the nodes could end up scattered all over your address space. That would lead to even more page faults. Even with the customary clustering memory managers, associative containers tend to have more problems with page faults, because, unlike contiguous-memory containers such as vector, node-based containers find it more difficult to guarantee that container elements that are close to one another in a container's traversal order are also close to one another in physical memory. Yet this is precisely the kind of memory organization that minimizes page faults when performing a binary search.

Bottom line: storing data in a sorted vector is likely to consume less memory than storing the same data in a standard associative container, and searching a sorted vector via binary search is likely to be faster than searching a standard associative container when page faults are taken into account.

Of course, the big drawback of a sorted vector is that it must remain sorted! When a new element is inserted, everything beyond the new element must be moved up by one. That's as expensive as it sounds, and it gets even more expensive if the vector has to reallocate its underlying memory (see [Item 14](#)), because then *all* the elements in the vector typically have to be copied. Similarly, if an element is removed from the vector, all the elements beyond it must be moved down. Insertions and erasures are expensive for vectors, but they're cheap for associative containers. That's why it makes sense to consider using a sorted vector instead of an associative container only when you know that your data structure is used in such a way that lookups are almost never mixed with insertions and erasures.

This Item has featured a lot of text, but it's been woefully short on examples, so let's take a look at a code skeleton for using a sorted vector instead of a set:

[illegible]

```

Widget w;                                // object for value to look up
...                                       // start Lookup phase
if (binary_search(vw.begin(), vw.end(), w)) ... // lookup via binary_search
vector<Widget>::iterator i =
    lower_bound(vw.begin(), vw.end(), w); // lookup via lower_bound;
if (i != vw.end() && !(w < *i)) ...       // see Item 19 for an explana-
                                           // tion of the "!(w < *i)" test

pair<vector<Widget>::iterator,
    vector<Widget>::iterator> range =
    equal_range(vw.begin(), vw.end(), w); // lookup via equal_range
if (range.first != range.second) ...

...                                     // end Lookup phase, start
                                           // Reorganize phase

sort(vw.begin(), vw.end());             // begin new Lookup phase...

```

As you can see, it's all pretty straightforward. The hardest thing about it is deciding among the search algorithms (e.g., `binary_search`, `lower_bound`, etc.), and [Item 45](#) helps you do that.

Things get a bit more interesting when you decide to replace a map or multimap with a vector, because the vector must hold pair objects. After all, that's what map and multimap hold. Recall, however, that if you declare an object of type `map<K, V>` (or its multimap equivalent), the type of elements stored in the map is `pair<const K, V>`. To emulate a map or multimap using a vector, you must omit the `const`, because when you sort the vector, the values of its elements will get moved around via assignment, and that means that both components of the pair must be assignable. When using a vector to emulate a `map<K, V>`, then, the type of the data stored in the vector will be `pair<K, V>`, not `pair<const K, V>`.

maps and multimaps keep their elements in sorted order, but they look only at the key part of the element (the first component of the pair) for sorting purposes, and you must do the same when sorting a vector. You'll need to write a custom comparison function for your pairs, because pair's `operator<` looks at *both* components of the pair.

Interestingly, you'll need a second comparison function for performing lookups. The comparison function you'll use for sorting will take two pair objects, but lookups are performed given only a key value. The comparison function for lookups, then, must take an object of the key type (the value being searched for) and a pair (one of the pairs stored in the vector) — two different types. As an additional twist, you can't know whether the key value or the pair will be passed as the first argument, so you really need two comparison functions for lookups: one

where the key value is passed first and one where the pair is passed first.

Here's an example of how to put all the pieces together:

```
typedef pair<string, int> Data;           // type held in the "map"
                                         // in this example

class DataCompare {                     // class for comparison
public:                                  // functions

    bool operator()(const Data& lhs,      // comparison func
                    const Data& rhs) const // for sorting
    {
        return keyLess(lhs.first, rhs.first); // keyLess is below
    }

    bool operator()(const Data& lhs,      // comparison func
                    const Data::first_type& k) const // for lookups
                                                // (form 1)
    {
        return keyLess(lhs.first, k);
    }

    bool operator()(const Data::first_type& k, // comparison func
                    const Data& rhs) const    // for lookups
                                                // (form 2)
    {
        return keyLess(k, rhs.first);
    }

private:
    bool keyLess(const Data::first_type& k1, // the "real"
                 const Data::first_type& k2) const // comparison
    {
        return k1 < k2; // function
    }
};
```

In this example, we assume that our sorted vector will be emulating a `map<string, int>`. The code is pretty much a literal translation of the discussion above, except for the presence of the member function `keyLess`. That function exists to ensure consistency among the various `operator()` functions. Each such function simply compares two key values, so rather than duplicate the logic, we put the test inside `keyLess` and have the `operator()` functions return whatever `keyLess` does. This admirable act of software engineering enhances the maintainability of `DataCompare`, but there is a minor drawback. Its provision for `operator()` functions with different parameter types renders the resulting function objects unadaptable (see [Item 40](#)). Oh well.

Using a sorted vector as a map is essentially the same as using it as a set. The only major difference is the need to use `DataCompare` objects as comparison functions:

```

vector<Data> vd;                                // alternative to
                                                // map<string, int>

...                                              // Setup phase: lots of
                                                // insertions, few lookups

sort(vd.begin(), vd.end(), DataCompare());      // end of Setup phase. Simu-
                                                // lated maps must avoid
                                                // dupes; simulated multimaps
                                                // might use stable_sort

string s;                                        // object for value to look up
...                                              // start Lookup phase

if (binary_search(vd.begin(), vd.end(), s,
                  DataCompare())) ...          // lookup via binary_search

vector<Data>::iterator i =
    lower_bound(vd.begin(), vd.end(), s,
                DataCompare());                // lookup via lower_bound;
if (i != vd.end() && !DataCompare()(s, *i))... // again, see Item 19 for info
                                                // on the
                                                // "!"DataCompare()(s, *i)" test

pair<vector<Data>::iterator,
    vector<Data>::iterator> range =
    equal_range(vd.begin(), vd.end(), s,
                DataCompare());                // lookup via equal_range
if (range.first != range.second) ...

...                                              // end Lookup phase, start
                                                // Reorganize phase

sort(vd.begin(), vd.end(), DataCompare());      // begin new Lookup phase...

```

As you can see, once you've written `DataCompare`, things pretty much fall into place. And once in place, they'll often run faster and use less memory than the corresponding design using a real map *as long as your program uses the data structure in the phased manner described on [page 101](#)*. If your program doesn't operate on the data structure in a phased manner, use of a sorted vector instead of a standard associative container will almost certainly end up wasting time.

Item 24: Choose carefully between `map::operator[]` and `map::insert` when efficiency is important.

Let's suppose we have a `Widget` class that supports default construction as well as construction and assignment from a `double`:

```

class Widget {
public:
    Widget();
    Widget(double weight);
    Widget& operator=(double weight);
    ...
};

```

Let's now suppose we'd like to create a map from ints to Widgets, and we'd like to initialize the map with particular values. 'Tis simplicity itself:

```
map<int, Widget> m;  
m[1] = 1.50;  
m[2] = 3.67;  
m[3] = 10.5;  
m[4] = 45.8;  
m[5] = 0.0003;
```

In fact, the only thing simpler is forgetting what's really going on. That's too bad, because what's going on could incur a considerable performance hit.

The operator[] function for maps is a novel creature, unrelated to the operator[] functions for vectors, deque, and strings and equally unrelated to the built-in operator[] that works with arrays. Instead, map::operator[] is designed to facilitate "add or update" functionality. That is, given

```
map<K, V> m;
```

the expression

```
m[k] = v;
```

checks to see if the key *k* is already in the map. If not, it's added, along with *v* as its corresponding value. If *k* is already in the map, its associated value is updated to *v*.

The way this works is that operator[] returns a reference to the value object associated with *k*. *v* is then assigned to the object to which the reference (the one returned from operator[]) refers. This is straightforward when an existing key's associated value is being updated, because there's already a value object to which operator[] can return a reference. But if *k* isn't yet in the map, there's no value object for operator[] to refer to. In that case, it creates one from scratch by using the value type's default constructor. operator[] then returns a reference to this newly-created object.

Let's look again at the first part of our original example:

```
map<int, Widget> m;  
m[1] = 1.50;
```

The expression `m[1]` is shorthand for `m.operator[](1)`, so this is a call to `map::operator[]`. That function must return a reference to a Widget, because *m*'s mapped type is Widget. In this case, *m* doesn't yet have anything in it, so there is no entry in the map for the key 1. operator[]

therefore default-constructs a `Widget` to act as the value associated with 1, then returns a reference to that `Widget`. Finally, the `Widget` becomes the target of an assignment; the assigned value is 1.50.

In other words, the statement

```
m[1] = 1.50;
```

is functionally equivalent to this:

```
typedef map<int, Widget> IntWidgetMap;           // convenience
                                                // typedef

pair<IntWidgetMap::iterator, bool> result =      // create new map
m.insert(IntWidgetMap::value_type(1, Widget())); // entry with key 1
                                                // and a default-
                                                // constructed value
                                                // object; see below
                                                // for a comment on
                                                // value_type

result.first->second = 1.50;                     // assign to the
                                                // newly-constructed
                                                // value object
```

Now it should be clear why this approach may degrade performance. We first default-construct a `Widget`, then we immediately assign it a new value. If it's measurably more efficient to construct a `Widget` with the value we want instead of default-constructing the `Widget` and then doing the assignment, we'd be better off replacing our use of `operator[]` (including its attendant construction plus assignment) with a straightforward call to `insert`:

```
m.insert(IntWidgetMap::value_type(1, 1.50));
```

This has precisely the same ultimate effect as the code above, except it replaces construction (and later destruction) of a temporary *plus an assignment* with construction of a temporary from a pair of values (followed by later destruction of that temporary); there's no assignment. The more expensive the cost of assignment of the value part of the pair, the more you save by using `map::insert` instead of `map::operator[]`.

The code above takes advantage of the `value_type` typedef that's provided by every standard container. There's nothing particularly significant about this typedef, but it's important to remember that for `map` and `multimap` (as well as the nonstandard containers `hash_map` and `hash_multimap` — see [Item 25](#)), the type of the contained elements will always be some kind of pair.

I remarked earlier that `operator[]` is designed to facilitate “add or update” functionality, and we now understand that when an “add” is performed, `insert` is more efficient than `operator[]`. The situation is reversed when we do an update, i.e., when an equivalent key (see [Item 19](#)) is already in the map. To see why that is the case, look at our update options:

```
m[k] = v;                                // use operator[]
                                           // to update k's
                                           // value to be v

m.insert(IntWidgetMap::value_type(k, v)).first->second = v; // use insert to
                                                           // update k's
                                                           // value to be v
```

The syntax alone is probably enough to convince you to favor `operator[]`, but we’re focusing on efficiency here, so we’ll overlook that.

The call to `insert` requires an argument of type `IntWidgetMap::value_type` (i.e., `pair<int, Widget>`), so when we call `insert`, we must construct and destruct an object of that type. That costs us a pair constructor and destructor. That, in turn, entails a `Widget` construction and destruction, because a `pair<int, Widget>` itself contains a `Widget` object. `operator[]` uses no pair object, so it constructs and destructs no pair and no `Widget`.

Efficiency considerations thus lead us to conclude that `insert` is preferable to `operator[]` when adding an element to a map, and both efficiency and aesthetics dictate that `operator[]` is preferable when updating the value of an element that’s already in the map.

It would be nice if the STL provided a function that offered the best of both worlds, i.e., efficient add-or-update functionality in a syntactically attractive package. For example, it’s not hard to envision a calling interface such as this:

```
iterator affectedPair =
    efficientAddOrUpdate(m, k, v);          // if key k isn't in map m, efficiently
                                           // add pair (k,v) to m; otherwise
                                           // efficiently update to v the value
                                           // associated with k. Return an
                                           // iterator to the added or
                                           // modified pair
```

There’s no function like this in the STL, but, as the code below demonstrates, it’s not terribly hard to write yourself. The comments summarize what’s going on, and the paragraphs that follow provide some additional explanation.

```

template< typename MapType,           // type of map
          typename KeyArgType,       // see below for why
          typename ValueArgType>     // KeyArgType and
                                     // ValueArgType are type
                                     // parameters
typename MapType::iterator
efficientAddOrUpdate(MapType& m,
                    const KeyArgType& k,
                    const ValueArgType& v)
{
    typename MapType::iterator lb = // find where k is or should
        m.lower_bound(k);           // be; see page 7 for why
                                     // "typename" is needed
                                     // here

    if (lb != m.end() &&             // if lb points to a pair
        !(m.key_comp()(k, lb->first))) { // whose key is equiv to k...

        lb->second = v;              // update the pair's value
        return lb;                  // and return an iterator
    }                               // to that pair

    else {
        typedef typename MapType::value_type MVT;
        return m.insert(lb, MVT(k, v)); // add pair(k, v) to m and
    }                                   // return an iterator to the
}                                     // new map element

```

To perform an efficient add or update, we need to be able to find out if *k*'s value is in the map; if so, where it is; and if not, where it should be inserted. That job is tailor-made for `lower_bound` (see [Item 45](#)), so that's the function we call. To determine whether `lower_bound` found an element with the key we were looking for, we perform the second half of an equivalence test (see [Item 19](#)), being sure to use the correct comparison function for the map; that comparison function is available via `map::key_comp`. The result of the equivalence test tells us whether we're performing an add or an update.

If it's an update, the code is straightforward. The insert branch is more interesting, because it uses the "hint" form of insert. The construct `m.insert(lb, MVT(k, v))` "hints" that `lb` identifies the correct insertion location for a new element with key equivalent to *k*, and the Standard guarantees that if the hint is correct, the insertion will occur in amortized constant, rather than logarithmic, time. In `efficientAddOrUpdate`, we know that `lb` identifies the proper insertion location, so the call to insert is guaranteed to be an amortized constant-time operation.

An interesting aspect of this implementation is that `KeyArgType` and `ValueArgType` need not be the types stored in the map. They need only be *convertible* to the types stored in the map. The alternative would be to eliminate the type parameters `KeyArgType` and `ValueArgType`, using

instead `MapType::key_type` and `MapType::mapped_type`. However, if we did that, we might force unnecessary type conversions to occur at the point of call. For instance, look again at the definition of the map we've been using in this Item's examples:

```
map<int, Widget> m; // as before
```

And recall that `Widget` accepts assignment from a double:

```
class Widget { // also as before
public:
    ...
    Widget& operator=(double weight);
    ...
};
```

Now consider this call to `efficientAddOrUpdate`:

```
efficientAddOrUpdate(m, 10, 1.5);
```

Suppose that it's an update operation, i.e., `m` already contains an element whose key is 10. In that case, the template above deduces that `ValueArgType` is `double`, and the body of the function directly assigns 1.5 as a double to the `Widget` associated with the key 10. That's accomplished by an invocation of `Widget::operator=(double)`. If we had used `MapType::mapped_type` as the type of `efficientAddOrUpdate`'s third parameter, we'd have converted 1.5 to a `Widget` at the point of call, and we'd thus have paid for a `Widget` construction (and subsequent destruction) we didn't need.

Subtleties in the implementation of `efficientAddOrUpdate` may be interesting, but they're not as important as the main point of this Item, which is that you should choose carefully between `map::operator[]` and `map::insert` when efficiency is important. If you're updating an existing map element, `operator[]` is preferable, but if you're adding a new element, `insert` has the edge.

Item 25: Familiarize yourself with the nonstandard hashed containers.

It generally doesn't take much time for STL programmers to begin to wonder, "Vectors, lists, maps, sure, but where are the hash tables?" Alas, there aren't any hash tables in the standard C++ library. Everyone agrees that this is unfortunate, but the Standardization Committee felt that the work needed to add them would have unduly delayed completion of the Standard. It's a foregone conclusion that the next version of the Standard will include hash tables, but for the time being, the STL doesn't do hashing.

If you like hash tables, however, take heart. You need not do without or roll your own. STL-compatible hashed associative containers are available from multiple sources, and they even have *de facto* standard names: `hash_set`, `hash_multiset`, `hash_map`, and `hash_multimap`.

Behind these common names, different implementations, er, differ. They differ in their interfaces, their capabilities, their underlying data structures, and the relative efficiency of the operations they support. It's still possible to write reasonably portable code using hash tables, but it's not as easy as it would be had the hashed containers been standardized. (Now you know why standards are important.)

Of the several available implementations for hashed containers, the two most common are from SGI (see [Item 50](#)) and Dinkumware (see [Appendix B](#)), so in what follows, I restrict myself to the designs of the hashed containers from these vendors. STLport (again, see [Item 50](#)) also offers hashed containers, but the STLport hashed containers are based on those from SGI. For purposes of this Item, assume that whatever I write about the SGI hashed containers applies to the STLport hashed containers, too.

Hashed containers are associative containers, so it should not surprise you that, like all associative containers, they need to know the type of objects stored in the container, the comparison function for such objects, and the allocator for these objects. In addition, hashed containers require specification of a hashing function. That suggests the following declaration for hashed containers:

```
template<typename T,
        typename HashFunction,
        typename CompareFunction,
        typename Allocator = allocator<T> >
class hash_container;
```

This is quite close to the SGI declaration for hashed containers, the primary difference being that SGI provides default types for `HashFunction` and `CompareFunction`. The SGI declaration for `hash_set` looks essentially like this (I've tweaked it a bit for presentation purposes):

```
template<typename T,
        typename HashFunction = hash<T>,
        typename CompareFunction = equal_to<T>,
        typename Allocator = allocator<T> >
class hash_set;
```

A noteworthy aspect of the SGI design is the use of `equal_to` as the default comparison function. This is a departure from the conventions of the standard associative containers, where the default comparison function is `less`. This design decision signifies more than a simple

change in default comparison functions. SGI's hashed containers determine whether two objects in a hashed container have the same value by testing for *equality*, not equivalence (see [Item 19](#)). For hashed containers, this is not an unreasonable decision, because hashed associative containers, unlike their standard (typically tree-based) counterparts, are not kept in sorted order.

The Dinkumware design for hashed containers adopts some different strategies. It still allows you to specify object types, hash function types, comparison function types, and allocator types, but it moves the default hash and comparison functions into a separate traits-like class called `hash_compare`, and it makes `hash_compare` the default argument for the `HashingInfo` parameter of the container templates. (If you're unfamiliar with the notion of a "traits" class, open a good STL reference like Josuttis' *The C++ Standard Library* [3] and study the motivation and implementation of the `char_traits` and `iterator_traits` templates.)

For example, here's the Dinkumware `hash_set` declaration (again, tweaked for presentation):

```
template<typename T, typename CompareFunction>
class hash_compare;

template<typename T,
        typename HashingInfo = hash_compare<T, less<T> >,
        typename Allocator = allocator<T> >
class hash_set;
```

The interesting part of this interface design is the use of `HashingInfo`. The container's hashing and comparison functions are stored there, but the `HashingInfo` type also holds enums controlling the minimum number of buckets in the table as well as the maximum allowable ratio of container elements to buckets. When this ratio is exceeded, the number of buckets in the table is increased, and some elements in the table are rehashed. (SGI provides member functions that afford similar control over the number of table buckets and, hence, the ratio of table elements to buckets.)

After some tweaks for presentation, `hash_compare` (the default value for `HashingInfo`) looks more or less like this:

```
template<typename T, typename CompareFunction = less<T> >
class hash_compare {
public:
    enum {
        bucket_size = 4,                // max ratio of elements to buckets
        min_buckets = 8                 // minimum number of buckets
    };
};
```

```

    size_t operator()(const T&) const;    // hash function
    bool operator()(const T&,           // comparison function
                    const T&) const;
    ...
};                                       // a few things omitted, including
                                       // the use of CompareFunction

```

The overloading of `operator()` (in this case, to implement both the hashing and comparison functions) is a strategy that crops up more frequently than you might imagine. For another application of the same idea, take a look at [Item 23](#).

The Dinkumware design allows you to write your own `hash_compare`-like class (possibly by deriving from `hash_compare` itself), and as long as your class provides definitions for `bucket_size`, `min_buckets`, two `operator()` functions (one taking one argument, one taking two), plus a few things I've left out, you can use it to control the configuration and behavior of a Dinkumware `hash_set` or `hash_multiset`. Configuration control for `hash_map` and `hash_multimap` is similar.

Notice that with both the SGI and the Dinkumware designs, you can leave all the decision-making to the implementations and simply write something like this:

```

    hash_set<int> intTable;           // create a hashed set of ints

```

For this to compile, the hash table must hold an integral type (such as `int`), because the default hashing functions are generally limited to integral types. (SGI's default hashing function is slightly more flexible. [Item 50](#) will tell you where to find all the details.)

Under the hood, the SGI and Dinkumware implementations go their very separate ways. SGI employs a conventional open hashing scheme composed of an array (the buckets) of pointers to singly linked lists of elements. Dinkumware also employs open hashing, but its design is based on a novel data structure consisting of an array of iterators (essentially the buckets) into a doubly linked list of elements, where adjacent pairs of iterators identify the extent of the elements in each bucket. (For details, consult Plauger's aptly titled column, "Hash Tables" [\[16\]](#).)

As a user of these implementations, it's likely you'll be interested in the fact that the SGI implementation stores the table elements in singly linked lists, while the Dinkumware implementation uses a doubly linked list. The difference is noteworthy, because it affects the iterator categories for the two implementations. SGI's hashed containers offer forward iterators, so you forgo the ability to perform reverse iterations; there are no `rbegin` or `rend` member functions in SGI's hashed contain-

ers. The iterators for Dinkumware's hashed containers are bidirectional, so they offer both forward and reverse traversals. In terms of memory usage, the SGI design is a bit more parsimonious than that from Dinkumware.

Which design is best for you and your applications? I can't possibly know. Only you can determine that, and this Item hasn't tried to give you enough information to come to a reasonable conclusion. Instead, the goal of this Item is to make sure you know that though the STL itself lacks hashed containers, STL-compatible hashed containers (with varying interfaces, capabilities, and behavioral trade-offs) are not difficult to come by. In the case of the SGI and STLport implementations, you can even come by them for free, because they're available for free download.

4

Iterators

At first glance, iterators appear straightforward. Look more closely, however, and you'll notice that the standard STL containers offer four different iterator types: `iterator`, `const_iterator`, `reverse_iterator`, and `const_reverse_iterator`. From there it's only a matter of time until you note that of these four types, only one is accepted by containers in certain forms of insert and erase. That's when the questions begin. Why four iterator types? What is the relationship among them? Are they interconvertible? Can the different types be mixed in calls to algorithms and STL utility functions? How do these types relate to containers and their member functions?

This chapter answers these questions, and it introduces an iterator type that deserves more notice than it usually gets: `istreambuf_iterator`. If you like the STL, but you're unhappy with the performance of `istream iterators` when reading character streams, `istreambuf_iterator` could be just the tool you're looking for.

Item 26: Prefer `iterator` to `const_iterator`, `reverse_iterator`, and `const_reverse_iterator`.

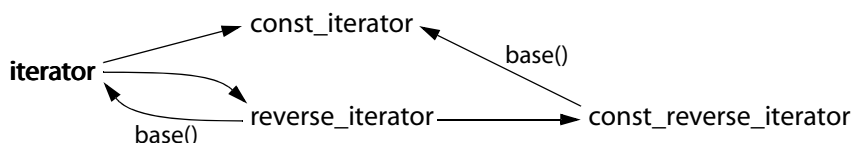
As you know, every standard container offers four types of iterator. For a container<T>, the type `iterator` acts like a `T*`, while `const_iterator` acts like a `const T*` (which you may also see written as a `T const*`; they mean the same thing). Incrementing an `iterator` or a `const_iterator` moves you to the next element in the container in a traversal starting at the front of the container and proceeding to the back. `reverse_iterator` and `const_reverse_iterator` also act like `T*` and `const T*`, respectively, except that incrementing these iterator types moves you to the next element in the container in a traversal from back to front.

Let me show you two things. First, take a look at some signatures for insert and erase in `vector<T>`:

```
iterator insert(iterator position, const T& x);  
iterator erase(iterator position);  
iterator erase(iterator rangeBegin, iterator rangeEnd);
```

Every standard container offers functions analogous to these, though the return types vary, depending on the container type. The thing to notice is that these functions demand parameters of type `iterator`. Not `const_iterator`, not `reverse_iterator`, not `const_reverse_iterator`. Always `iterator`. Though containers support four iterator types, one of those types has privileges the others do not have. That type is `iterator`. `iterator` is special.

The second thing I want to show you is this diagram, which displays the conversions that exist among iterator types.



The diagram shows that there are implicit conversions from `iterator` to `const_iterator`, from `iterator` to `reverse_iterator`, and from `reverse_iterator` to `const_reverse_iterator`. It also shows that a `reverse_iterator` may be converted into an `iterator` by using the `reverse_iterator`'s `base` member function, and a `const_reverse_iterator` may similarly be converted into a `const_iterator` via `base`. The diagram does not show that the iterators obtained from `base` may not be the ones you want. For the story on that, turn to [Item 28](#).

You'll observe that there is no way to get from a `const_iterator` to an `iterator` or from a `const_reverse_iterator` to a `reverse_iterator`. This is important, because it means that if you have a `const_iterator` or a `const_reverse_iterator`, you'll find it difficult to use those iterators with some container member functions. Such functions demand iterators, and since there's no conversion path from the `const` iterator types back to `iterator`, the `const` iterator types are largely useless if you want to use them to specify insertion positions or elements to be erased.

Don't be fooled into thinking that this means `const` iterators are useless in general. They're not. They're perfectly useful with algorithms, because algorithms don't usually care what kind of iterators they work with, as long as they are of the appropriate category. `const` iterators are also acceptable for many container member functions. It's only some forms of `insert` and `erase` that are picky.

I wrote that `const` iterators are "largely" useless if you want to specify insertion positions or elements to be erased. The implication is that

they are not completely useless. That's true. They can be useful if you can find a way to get an iterator from a `const_iterator` or from a `const_reverse_iterator`. That's often possible. It isn't *always* possible, however, and even when it is, the way to do it isn't terribly obvious. It may not be terribly efficient, either. The topic is meaty enough to justify its own Item, so turn to [Item 27](#) if you're interested in the details. For now, we have enough information to understand why it often makes sense to prefer iterators to `const` and reverse iterators:

- Some versions of `insert` and `erase` require iterators. If you want to call those functions, you're going to have to produce iterators. `const` and reverse iterators won't do.
- It's not possible to implicitly convert a `const_iterator` to an iterator, and the technique described in [Item 27](#) for generating an iterator from a `const_iterator` is neither universally applicable nor guaranteed to be efficient.
- Conversion from a `reverse_iterator` to an iterator may require iterator adjustment after the conversion. [Item 28](#) explains when and why.

All these things conspire to make working with containers easiest, most efficient, and least likely to harbor subtle bugs if you prefer iterators to their `const` and reverse colleagues.

Practically speaking, you are more likely to have a choice when it comes to iterators and `const_iterators`. The decision between iterator and `reverse_iterator` is often made for you. You either need a front-to-back traversal or a back-to-front traversal, and that's pretty much that. You choose the one you need, and if that means choosing `reverse_iterator`, you choose `reverse_iterator` and use `base` to convert it to an iterator (possibly preceded by an offset adjustment — see [Item 28](#)) when you want to make calls to container member functions that require iterators.

When choosing between iterators and `const_iterators`, there are reasons to choose iterators even when you could use a `const_iterator` and when you have no need to use the iterator as a parameter to a container member function. One of the most irksome involves comparisons between iterators and `const_iterators`. I hope we can all agree that this is reasonable code:

```
typedef deque<int> IntDeque;           // STL container and
typedef IntDeque::iterator Iter;      // iterator types are easier
typedef IntDeque::const_iterator ConstIter; // to work with if you
                                         // use some typedefs

Iter i;
ConstIter ci;
```

```

...                                     // make i and ci point into
                                     // the same container
if (i == ci) ...                       // compare an iterator
                                     // and a const_iterator

```

All we're doing here is comparing two iterators into a container, the kind of comparison that's the bread and butter of the STL. The only twist is that one object is of type `iterator` and one is of type `const_iterator`. This should be no problem. The `iterator` should be implicitly converted into a `const_iterator`, and the comparison should be performed between two `const_iterators`.

With well-designed STL implementations, this is precisely what happens, but with other implementations, the code will not compile. The reason is that such implementations declare `operator==` for `const_iterators` as a member function instead of as a non-member function, but the cause of the problem is likely to be of less interest to you than the workaround, which is to swap the order of the iterators, like this:

```

if (ci == i) ...                      // workaround for when the
                                     // above won't compile

```

This kind of problem can arise whenever you mix iterators and `const_iterators` (or `reverse_iterators` and `const_reverse_iterators`) in the same expression, not just when you are comparing them. For example, if you try to subtract one random access iterator from another,

```

if (i - ci >= 3) ...                  // if i is at least 3 beyond ci ...

```

your (valid) code may be (incorrectly) rejected if the iterators aren't of the same type. The simplest workaround in this case is to turn the iterator into a `const_iterator` through a (safe) cast:

```

if (static_cast<ConstIter>(i) - ci >= 3) ... // workaround for when the
                                     // above won't compile

```

The easiest way to guard against these kinds of problems is to minimize your mixing of iterator types, and that, in turn, leads back to preferring iterators to `const_iterators`. From the perspective of const correctness (a worthy perspective, to be sure), staying away from `const_iterators` simply to avoid potential implementation shortcomings (all of which have workarounds) seems unjustified, but in conjunction with the anointed status of iterators in some container member functions, it's hard to avoid the practical conclusion that `const_iterators` are not only less useful than iterators, sometimes they're just not worth the trouble.

Item 27: Use distance and advance to convert a container's const_iterators to iterators.

[Item 26](#) points out that some container member functions that take iterators as parameters insist on *iterators*; *const_iterators* won't do. So what do you do if you have a *const_iterator* in hand and you want to, say, insert a new value into a container at the position indicated by the iterator? Somehow you've got to turn your *const_iterator* into an iterator, and you have to take an active role in doing it, because, as [Item 26](#) explains, there is no implicit conversion from *const_iterator* to iterator.

I know what you're thinking. You're thinking, "When all else fails, get a bigger hammer." In the world of C++, that can mean only one thing: casting. Shame on you for such thoughts. Where do you get these ideas?

Let us confront this cast obsession of yours head on. Look what happens when you try to cast a *const_iterator* to an iterator:

```
typedef deque<int> IntDeque;           // convenience typedefs
typedef IntDeque::iterator Iter;
typedef IntDeque::const_iterator ConstIter;
ConstIter ci;                         // ci is a const_iterator
...
Iter i(ci);                           // error! no implicit conversion from
                                     // const_iterator to iterator
Iter i(const_cast<Iter>(ci));          // still an error! can't cast a
                                     // const_iterator to an iterator
```

This example happens to use *deque*, but the outcome would be the same for *list*, *set*, *multiset*, *map*, *multimap*, and the hashed containers described in [Item 25](#). The line using the cast *might* compile in the case of *vector* or *string*, but those are special cases we'll consider in a moment.

The reason the cast won't compile is that for these container types, *iterator* and *const_iterator* are different classes, barely more closely related to one another than *string* and *complex<double>*. Trying to cast one type to the other is nonsensical, and that's why the *const_cast* is rejected. *static_cast*, *reinterpret_cast*, and a C-style cast would lead to the same end.

Alas, the cast that won't compile might compile if the iterators' container were a *vector* or a *string*. That's because it is common for implementations of these containers to use pointers as iterators. For such

implementations, `vector<T>::iterator` is a typedef for `T*`, `vector<T>::const_iterator` is a typedef for `const T*`, `string::iterator` is a typedef for `char*`, and `string::const_iterator` is a typedef for `const char*`. With such implementations, the `const_cast` from a `const_iterator` to an `iterator` will compile and do the right thing, because the cast is converting a `const T*` into a `T*`. Even under such implementations, however, `reverse_iterators` and `const_reverse_iterators` are true classes, so you can't `const_cast` a `const_reverse_iterator` to a `reverse_iterator`. Also, as [Item 50](#) explains, even implementations where `vector` and `string` iterators are pointers might use that representation only when compiling in release mode. All these factors lead to the conclusion that casting `const` iterators to iterators is ill-advised even for `vector` and `string`, because its portability is doubtful.

If you have access to the container a `const_iterator` came from, there is a safe, portable way to get its corresponding iterator,[†] and it involves no circumvention of the type system. Here's the essence of the solution, though it must be modified slightly before it will compile:

```
typedef deque<int> IntDeque;           // as before
typedef IntDeque::iterator lter;
typedef IntDeque::const_iterator Constlter;

IntDeque d;
Constlter ci;
...                                  // make ci point into d
lter i(d.begin());                  // initialize i to d.begin()
advance(i, distance(i, ci));         // move i up to where ci is
                                     // (but see below for why this must
                                     // be tweaked before it will compile)
```

This approach is so simple and direct, it's startling. To get an iterator pointing to the same container element as a `const_iterator`, create a new iterator at the beginning of the container, then move it forward until it's as far from the beginning of the container as the `const_iterator` is! This task is facilitated by the function templates `advance` and `distance`, both of which are declared in `<iterator>`. `distance` reports how far apart two iterators into the same container are, and `advance` moves an iterator a specified distance. When `i` and `ci` point into the same container, the expression `advance(i, distance(i, ci))` makes `i` and `ci` point to the same place in the container.

Well, it would if it would compile, but it won't. To see why, look at the declaration for `distance`:

[†] I have since discovered that the approach described here may fail for string implementations that use reference counting. For details, consult jep's 8/22/01 comment regarding page 121 of this book at <http://www.aristeia.com/BookErrata/est11e-errata.html>.

```
template <typename InputIterator>
typename iterator_traits<InputIterator>::difference_type
distance(InputIterator first, InputIterator last);
```

Don't get hung up on the fact that the return type of the function is 56 characters long and mentions dependent types like `difference_type`. Instead, focus your attention on the uses of the type parameter `InputIterator`:

```
template <typename InputIterator>
typename iterator_traits<InputIterator>::difference_type
distance(InputIterator first, InputIterator last);
```

When faced with a call to `distance`, your compilers must deduce the type represented by `InputIterator` by examining the arguments used in the call. Look again at the call to `distance` in the code I said wasn't quite right:

```
advance(i, distance(i, ci));           // move i up to where ci is
```

Two parameters are passed to `distance`, `i` and `ci`. `i` is of type `Iter`, which is a typedef for `deque<int>::iterator`. To compilers, that implies that `InputIterator` in the call to `distance` is `deque<int>::iterator`. `ci`, however, is of type `ConstIter`, which is a typedef for `deque<int>::const_iterator`. *That* implies that `InputIterator` is of type `deque<int>::const_iterator`. It's not possible for `InputIterator` to be two different types at the same time, so the call to `distance` fails, typically yielding some long-winded error message that may or may not indicate that the compiler couldn't figure out what type `InputIterator` is supposed to be.

To get the call to compile, you must eliminate the ambiguity. The easiest way to do that is to explicitly specify the type parameter to be used by `distance`, thus obviating the need for your compilers to figure it out for themselves:

```
advance(i, distance<ConstIter>(i, ci)); // figure the distance between
                                         // i and ci (as const_iterators),
                                         // then move i that distance
```

We now know how to use `advance` and `distance` to get an iterator corresponding to a `const_iterator`, but we have so far sidestepped a question of considerable practical interest: How efficient is this technique? The answer is simple. It's as efficient as the iterators allow it to be. For random access iterators (such as those sported by `vector`, `string`, and `deque`), it's a constant-time operation. For bidirectional iterators (i.e., those for all other standard containers and for some implementations of the hashed containers (see [Item 25](#))), it's a linear-time operation.

Because it may take linear time to produce an iterator equivalent to a `const_iterator`, you may wish to rethink designs that require producing

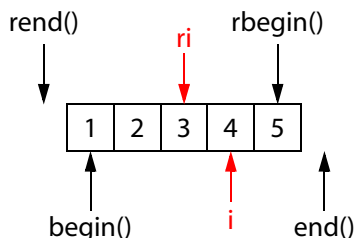
iterators from `const_iterators`. Such a consideration, in fact, helps motivate [Item 26](#), which advises you to prefer iterators over `const` and reverse iterators when dealing with containers.

Item 28: Understand how to use a reverse_iterator's base iterator.

Invoking the base member function on a `reverse_iterator` yields the “corresponding” iterator, but it’s not really clear what that means. As an example, take a look at this code, which puts the numbers 1-5 in a vector, sets a `reverse_iterator` to point to the 3, and initializes an iterator to the `reverse_iterator`’s base:

```
vector<int> v;
v.reserve(5);                                // see Item 14
for (int i = 1; i <= 5; ++i) {                // put 1-5 in the vector
    v.push_back(i);
}
vector<int>::reverse_iterator ri =             // make ri point to the 3
    find(v.rbegin(), v.rend(), 3);
vector<int>::iterator i(ri.base());            // make i the same as ri's base
```

After executing this code, things can be thought of as looking like this:



This picture is nice, displaying the characteristic offset of a `reverse_iterator` and its corresponding base iterator that mimics the offset of `rbegin()` and `rend()` with respect to `begin()` and `end()`, but it doesn’t tell you everything you need to know. In particular, it doesn’t explain how to use `i` to perform operations you’d like to perform on `ri`.

As [Item 26](#) explains, some container member functions accept only iterators as iterator parameters, so if you want to, say, insert a new element at the location identified by `ri`, you can’t do it directly; `vector`’s `insert` function won’t take `reverse_iterators`. You’d have a similar problem if you wanted to erase the element pointed to by `ri`. The erase

member functions reject `reverse_iterators`, insisting instead on iterators. To perform insertions or erasures, you must convert `reverse_iterators` into iterators via `base`, then use the iterators to get the jobs done.

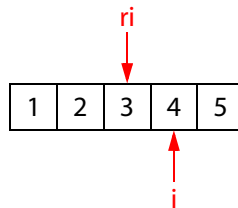
So let's suppose you *do* want to insert a new element into `v` at the position indicated by `ri`. In particular, let's assume you want to insert the value 99. Bearing in mind that `ri` is part of a traversal from right to left in the picture above and that insertion takes place in *front* of the element indicated by the iterator used to specify the insertion position, we'd expect the 99 to end up in front of the 3 with respect to a reverse traversal. After the insertion, then, `v` would look like this:

1	2	3	99	4	5
---	---	---	----	---	---

Of course, we can't use `ri` to indicate where to insert something, because it's not an iterator. We must use `i` instead. As noted above, when `ri` points at 3, `i` (which is `ri.base()`) points at 4. That's exactly where `i` needs to point for an insertion if the inserted value is to end up where it would have had we been able to use `ri` to specify the insertion location. Conclusion?

- To emulate insertion at a position specified by a `reverse_iterator` `ri`, insert at the position `ri.base()` instead. For purposes of insertion, `ri` and `ri.base()` are equivalent, and `ri.base()` is truly the iterator *corresponding* to `ri`.

Let us now consider erasing an element. Look again at the relationship between `ri` and `i` in the original vector (i.e., prior to the insertion of 99):



If we want to erase the element pointed to by `ri`, we can't just use `i`, because `i` doesn't point to the same element as `ri`. Instead, we must erase the element *preceding* `i`. Hence,

- To emulate erasure at a position specified by a `reverse_iterator` `ri`, erase at the position *preceding* `ri.base()` instead. For purposes of erasure, `ri` and `ri.base()` are *not* equivalent, and `ri.base()` is *not* the iterator corresponding to `ri`.

It's worth looking at the code to perform such an erasure, because it holds a surprise.

```
vector<int> v;
... // as above, put 1-5 in v
vector<int>::reverse_iterator ri = // as above, make ri point to the 3
    find(v.rbegin(), v.rend(), 3);
v.erase(--ri.base()); // attempt to erase at the position
                       // preceding ri.base(); for a vector,
                       // this will typically not compile
```

There's nothing wrong with this design. The expression `--ri.base()` correctly specifies the element we'd like to erase. Furthermore, this code will work with every standard container except `vector` and `string`. It might work with `vector` and `string`, too, but for many `vector` and `string` implementations, it won't compile. In such implementations, iterators (and `const_iterator`s) are implemented as built-in pointers, so the result of `ri.base()` is a pointer.

Both C and C++ dictate that pointers returned from functions shall not be modified, so for STL platforms where `string` and `vector` iterators are pointers, expressions like `--ri.base()` won't compile. To portably erase something at a position specified by a `reverse_iterator`, then, you must take pains to avoid modifying `base`'s return value. No problem. If you can't decrement the result of calling `base`, just increment the `reverse_iterator` and *then* call `base`!

```
... // as above
v.erase(++ri.base()); // erase the element pointed to by
                       // ri; this should always compile
```

Because this approach works with every standard container, it is the preferred technique for erasing an element pointed to by a `reverse_iterator`.

It should now be clear that it's not accurate to say that a `reverse_iterator`'s `base` member function returns the "corresponding" iterator. For insertion purposes, it does, but for erasure purposes, it does not. When converting `reverse_iterator`s to iterators, it's important that you know what you plan to do with the resulting iterator, because only then can you determine whether the iterator you have is the one you need.

Item 29: Consider `istreambuf_iterators` for character-by-character input.

Let's suppose you'd like to copy a text file into a string object. This seems like a pretty reasonable way to do it:

```
ifstream inputFile("interestingData.txt");
string fileData((istream_iterator<char>(inputFile)), // read inputFile into
                istream_iterator<char>());           // fileData; see below
                                                    // for why this isn't
                                                    // quite right, and see
                                                    // Item 6 for a warning
                                                    // about this syntax
```

It wouldn't take long before you'd notice that this approach fails to copy whitespace in the file into the string. That's because `istream_iterators` use `operator>>` functions to do the actual reading, and, by default, `operator>>` functions skip whitespace.

Assuming you'd like to retain the whitespace, all you need to do is override the default. Just clear the `skipws` flag for the input stream:

```
ifstream inputFile("interestingData.txt");
inputFile.unsetf(ios::skipws);                // disable the skipping of
                                              // whitespace in inputFile

string fileData((istream_iterator<char>(inputFile)),
                istream_iterator<char>());
```

Now all the characters in `inputFile` are copied into `fileData`.

Alas, you may discover that they aren't copied as quickly as you'd like. The `operator>>` functions on which `istream_iterators` depend perform formatted input, and that means they must undertake a fair amount of work on your behalf each time you call one. They have to create and destroy sentry objects (special `istream` objects that perform setup and cleanup activities for each call to `operator>>`), they have to check stream flags that might affect their behavior (e.g., `skipws`), they have to perform comprehensive checking for read errors, and, if they encounter a problem, they have to check the stream's exception mask to determine whether an exception should be thrown. Those are all important activities if you're performing formatted input, but if all you want to do is grab the next character from the input stream, it's overkill.

A more efficient approach is to use one of the STL's best kept secrets: `istreambuf_iterators`. You use `istreambuf_iterators` like `istream_iterators`, but where `istream_iterator<char>` objects use `operator>>` to read individual characters from an input stream, `istreambuf_iterator<char>` objects

go straight to the stream's buffer and read the next character directly. (More specifically, an `istreambuf_iterator<char>` object reading from an `istream` `s` will call `s.rdbuf()->sgetc()` to read `s`'s next character.)

Modifying our file-reading code to use `istreambuf_iterators` is so easy, most Visual Basic programmers need no more than two tries to get it right:

```
ifstream inputFile("interestingData.txt");
string fileData((istreambuf_iterator<char>(inputFile),
               istreambuf_iterator<char>()));
```

Notice how there's no need to "unset" the `skipws` flag here. `istreambuf_iterators` never skip any characters. Whatever's next in the stream buffer, that's what they grab.

Compared to `istream_iterators`, they grab it quickly — up to 40% faster in the simple benchmarks I performed, though don't be surprised if your mileage varies. Don't be surprised if the speed advantage increases over time, too, because `istreambuf_iterators` inhabit a seldom-visited corner of the STL where implementers haven't yet spent a lot of time on optimizations. For example, in one implementation I used, `istreambuf_iterators` were only about 5% faster than `istream_iterators` on my primitive tests. Such implementations clearly have lots of room to streamline their `istreambuf_iterator` implementations.

If you need to read the characters in a stream one by one, you don't need the power of formatted input, and you care about how long it takes to read the stream, typing three extra characters per iterator is a small price to pay for what is often a significant increase in performance. For unformatted character-by-character input, you should always consider `istreambuf_iterators`.

While you're at it, you should also consider `ostreambuf_iterators` for the corresponding unformatted character-by-character output operations. They avoid the overhead (and flexibility) of their `ostream_iterator` cousins, so they generally outperform them, too.

5

Algorithms

I noted at the beginning of [Chapter 1](#) that containers get the lion's share of the STL acclaim. In a sense, that's understandable. The containers are remarkable accomplishments, and they make life easier for legions of C++ programmers on a daily basis. Still, the STL algorithms are significant in their own right, equally capable of lightening a developer's burden. In fact, given that there are over 100 algorithms, it's easy to argue that they offer programmers a more finely honed tool set than the containers (a mere eight strong) could ever hope to match. Perhaps their number is part of the problem. Making sense of eight distinct container types is surely less work than remembering 70 algorithm names and trying to keep track of which does what.

I have two primary goals in this chapter. First, I want to introduce you to some lesser-known algorithms by showing you how they can make your life easier. Rest assured that I'm not going to punish you with lists of names to memorize. The algorithms I show you are in this chapter because they solve common problems, like performing case-insensitive string comparisons, efficiently finding the n most desirable objects in a container, summarizing the characteristics of all the objects in a range, and implementing the behavior of `copy_if` (an algorithm from the original HP STL that was dropped during standardization).

My second goal is to show you how to avoid common usage problems with the algorithms. You can't call `remove`, for example, or its cousins `remove_if` and `unique` unless you understand *exactly* what these algorithms do (and do *not* do). This is especially true when the range from which you're removing something holds pointers. Similarly, a number of algorithms work only with sorted ranges, so you need to understand which ones they are and why they impose that constraint. Finally, one of the most common algorithm-related mistakes involves asking an algorithm to write its results to a place that doesn't exist, so I explain how this absurdity can come about and how to ensure that you're not afflicted.

By the end of the chapter, you may not hold algorithms in the same high regard you probably already accord containers, but I'm hopeful you'll be willing to let them share the limelight more often than you have in the past.

Item 30: Make sure destination ranges are big enough.

STL containers automatically expand themselves to hold new objects as they are added (via `insert`, `push_front`, `push_back`, etc.). This works so well, some programmers lull themselves into the belief that they never have to worry about making room for objects in containers, because the containers themselves take care of things. If only it were so!

The problems arise when programmers *think* about inserting objects into containers, but don't tell the STL what they're thinking. Here's a common way this can manifest itself:

```
int transmogrify(int x);                // this function produces
                                        // some new value from x

vector<int> values;

...                                    // put data into values

vector<int> results;                   // apply transmogrify to
transform(values.begin(), values.end(), // each object in values,
           results.end(),               // appending the return
           transmogrify);               // values to results; this
                                        // code has a bug!
```

In this example, `transform` is told that the beginning of its destination range is `results.end()`, so that's where it starts writing the results of invoking `transmogrify` on every element of `values`. Like every algorithm that uses a destination range, `transform` writes its results by making assignments to the elements in the destination range. `transform` will thus apply `transmogrify` to `values[0]` and assign the result to `*results.end()`. It will then apply `transmogrify` to `values[1]` and assign the result to `*(results.end()+1)`. This can lead only to disaster, because *there is no object* at `*results.end()`, much less at `*(results.end()+1)`! The call to `transform` is wrong, because it's asking for assignments to be made to objects that don't exist. (Item 50 explains how a debugging implementation of the STL can detect this problem at runtime.)

Programmers who make this kind of mistake almost always intend for the results of the algorithm they're calling to be inserted into the destination container. If that's what you want to happen, you have to say so. The STL is a library, not a psychic. In this example, the way to say "please put `transform`'s results at the end of the container called `results`"

is to call `back_inserter` to generate the iterator specifying the beginning of the destination range:

```
vector<int> results;                                // apply transmogrify to
transform(values.begin(), values.end(),            // each object in values,
          back_inserter(results),                  // inserting the return
          transmogrify);                           // values at the end of
                                                    // results
```

Internally, the iterator returned by `back_inserter` causes `push_back` to be called, so you may use `back_inserter` with any container offering `push_back` (i.e., any of the standard sequence containers: `vector`, `string`, `deque`, and `list`). If you'd prefer to have an algorithm insert things at the front of a container you can use `front_inserter`. Internally, `front_inserter` makes use of `push_front`, so `front_inserter` works only for the containers offering that member function (i.e., `deque` and `list`):

```
...                                                // same as before
list<int> results;                                // results is now a list
transform(values.begin(), values.end(),            // insert transform's
          front_inserter(results),                 // results at the front of
          transmogrify);                           // results in reverse order
```

Because `front_inserter` causes each object added to `results` to be `push_fronted`, the order of the objects in `results` will be the *reverse* of the order of the corresponding objects in `values`. This is one reason why `front_inserter` isn't used as often as `back_inserter`. Another reason is that `vector` doesn't offer `push_front`, so `front_inserter` can't be used with `vectors`.

If you want `transform` to put its output at the front of `results`, but you also want the output to be in the same order as the corresponding objects in `values`, just iterate over `values` in reverse order:

```
list<int> results;                                // same as before
transform(values.rbegin(), values.rend(),          // insert transform's
          front_inserter(results),                 // results at the front of
          transmogrify);                           // results; preserve the
                                                    // relative object ordering
```

Given that `front_inserter` lets you force algorithms to insert their results at the front of a container and `back_inserter` lets you tell them to put their results at the back of a container, it's little surprise that `inserter` allows you to force algorithms to insert their results into containers at arbitrary locations:

```
vector<int> values;                                // as before
...
```

```

vector<int> results;                                // as before, except now
...                                                  // results has some data
                                                    // in it prior to the call to
                                                    // transform

transform(values.begin(), values.end(),             // insert the
           inserter(results, results.begin() + results.size() / 2), // results of
           transmogrify);                          // the trans-
                                                    // mogrifica-
                                                    // tions at
                                                    // the middle
                                                    // of results

```

Regardless of whether you use `back_inserter`, `front_inserter`, or `inserter`,[†] each insertion into the destination range is done one object at a time. [Item 5](#) explains that this can be expensive for contiguous-memory containers (vector, string, and deque), but [Item 5](#)'s suggested solution (using range member functions) can't be applied when it's an algorithm doing the inserting. In this example, `transform` will write to its destination range one value at a time, and there's nothing you can do to change that.

When the container into which you're inserting is a vector or a string, you can minimize the expense by following the advice of [Item 14](#) and calling `reserve` in advance. You'll still have to absorb the cost of shifting elements up each time an insertion takes place, but at least you'll avoid the need to reallocate the container's underlying memory:

```

vector<int> values;                                // as above
vector<int> results;
...
results.reserve(results.size() + values.size());    // ensure that results has
                                                    // the capacity for at least
                                                    // values.size() more
                                                    // elements

transform(values.begin(), values.end(),             // as above,
           inserter(results, results.begin() + results.size() / 2), // but results
           transmogrify);                          // won't do
                                                    // any reallo-
                                                    // cations

```

When using `reserve` to improve the efficiency of a series of insertions, always remember that `reserve` increases only a container's *capacity*; the container's *size* remains unchanged. Even after calling `reserve`, you must use an insert iterator (e.g., one of the iterators returned from `back_inserter`, `front_inserter`, or `inserter`) with an algorithm when you want that algorithm to add new elements to a vector or string.

[†] Unlike `back_inserter` and `front_inserter`, which just call `push_back` and `push_front`, respectively, `inserter` returns an iterator that keeps track of its most recent insertion location. Subsequent insertions take place following the most recently inserted element.

To make this absolutely clear, here is the *wrong* way to improve the efficiency of the example at the beginning of this Item (the one where we append to results the outcome of transmogrifying the data in values):

```
vector<int> values;                // as above
vector<int> results;
...
results.reserve(results.size() + values.size()); // as above
transform(values.begin(), values.end(),         // write the results of
           results.end(),                       // the transmogrifications
           transmogrify);                      // to uninitialized memory;
                                              // behavior is undefined!
```

In this code, transform blithely attempts to make assignments to the raw, uninitialized memory at the end of results. In general, this will fail at runtime, because assignment is an operation that makes sense only between two objects, not between one object and a chunk of primordial bits. Even if this code happens to do what you want it to, results won't know about the new "objects" transform "created" in its unused capacity. As far as results would be aware, its size would be the same after the call to transform as it was before. Similarly, its end iterator would point to the same place it did prior to the call to transform. Conclusion? Using reserve without also using an insert iterator leads to undefined behavior inside algorithms as well as to corrupted containers.

The correct way to code this example uses both reserve *and* an insert iterator:

```
vector<int> values;                // as above
vector<int> results;
...
results.reserve(results.size() + values.size()); // as above
transform(values.begin(), values.end(),         // write the results of
           back_inserter(results),             // the transmogrifications
           transmogrify);                      // to the end of results,
                                              // avoiding reallocations
                                              // during the process
```

So far, I've assumed that you want algorithms like transform to insert their results as new elements into a container. This is a common desire, but sometimes you want to overwrite the values of existing container elements instead of inserting new ones. When that's the case, you don't need an insert iterator, but you still need to follow the advice of this Item and make sure your destination range is big enough.

For example, suppose you want transform to overwrite results' elements. As long as results has at least as many elements as values does, that's easy. If it doesn't, you must either use `resize` to make sure it does,

```
vector<int> values;
vector<int> results;
...
if (results.size() < values.size()) {           // make sure results is at
    results.resize(values.size());             // least as big as values is
}
transform(values.begin(), values.end(),         // overwrite the first
           results.begin(),                     // values.size() elements of
           transmogrify);                       // results
```

or you can clear results and then use an insert iterator in the usual fashion:

```
...
results.clear();                               // destroy all elements in
                                              // results
results.reserve(values.size());                // reserve enough space
transform(values.begin(), values.end(),        // put transform's return
           back_inserter(results),             // values into results
           transmogrify);
```

This Item has demonstrated a number of variations on a theme, but I hope the underlying melody is what sticks in your mind. Whenever you use an algorithm requiring specification of a destination range, ensure that the destination range is big enough already or is increased in size as the algorithm runs. To increase the size as you go, use insert iterators, such as `ostream_iterators` or those returned by `back_inserter`, `front_inserter`, or `inserter`. That's all you need to remember.

Item 31: Know your sorting options.

How can I sort thee? Let me count the ways.

When many programmers think of ordering objects, only a single algorithm comes to mind: `sort`. (Some programmers think of `qsort`, but once they've read [Item 46](#), they recant and replace thoughts of `qsort` with those of `sort`.)

Now, `sort` is a wonderful algorithm, but there's no point in squandering wonder where you don't need it. Sometimes you don't need a full sort. For example, if you have a vector of `Widgets` and you'd like to

select the 20 highest-quality Widgets to send to your most loyal customers, you need to do only enough sorting to identify the 20 best Widgets; the remainder can remain unsorted. What you need is a partial sort, and there's an algorithm called `partial_sort` that does exactly what the name suggests:

```
bool qualityCompare(const Widget& lhs, const Widget& rhs)
{
    // return whether lhs's quality is greater than rhs's quality
}

...

partial_sort(widgets.begin(),                // put the best 20 elements
             widgets.begin() + 20,           // (in order) at the front of
             widgets.end(),                  // widgets
             qualityCompare);

...                                           // use widgets...
```

After the call to `partial_sort`, the first 20 elements of `widgets` are the best in the container and are in order, i.e., the highest-quality Widget is `widgets[0]`, the next highest is `widgets[1]`, etc. That makes it easy to send your best Widget to your best customer, the next best Widget to your next best customer, etc.

If all you care about is that the 20 best Widgets go to your 20 best customers, but you don't care which Widget goes to which customer, `partial_sort` gives you more than you need. In that case, all you need is the 20 best Widgets in *any* order. The STL has an algorithm that does exactly what you want, though the name isn't likely to spring to mind. It's called `nth_element`.

`nth_element` sorts a range so that the element at position *n* (which you specify) is the one that would be there if the range had been fully sorted. In addition, when `nth_element` returns, none of the elements in the positions up to *n* follow the element at position *n* in the sort order, and none of the elements in positions following *n* precede the element at position *n* in the sort order. If that sounds complicated, it's only because I have to select my words carefully. I'll explain why in a moment, but first let's look at how to use `nth_element` to make sure the best 20 Widgets are at the front of the `widgets` vector:

```
nth_element(widgets.begin(),                // put the best 20 elements
            widgets.begin() + 19,           // at the front of widgets,
            widgets.end(),                  // but don't worry about
            qualityCompare);               // their order
```

As you can see, the call to `nth_element` is essentially identical to the call to `partial_sort`. The only difference in their effect is that `partial_sort` sorts the elements in positions 1-20, while `nth_element` doesn't. Both algorithms, however, move the 20 highest-quality Widgets to the front of the vector.

That gives rise to an important question. What do these algorithms do when there are elements with the same level of quality? Suppose, for example, there are 12 elements with a quality rating of 1 (the best possible) and 15 elements with a quality rating of 2 (the next best). In that case, choosing the 20 best Widgets involves choosing the 12 with a rating of 1 and 8 of the 15 with a rating of 2. How should `partial_sort` and `nth_element` determine which of the 15 to put in the top 20? For that matter, how should `sort` figure out which order to put elements in when multiple elements have equivalent values?

`partial_sort` and `nth_element` order elements with equivalent values any way they want to, and you can't control this aspect of their behavior. (See [Item 19](#) for what it means for two values to be equivalent.) In our example, when faced with the need to choose Widgets with a quality rating of 2 to put into the last 8 spots in the vector's top 20, they'd choose whichever ones they wanted. That's not unreasonable. If you ask for the 20 best Widgets and some Widgets are equally good, you're in no position to complain as long as the 20 you get back are at least as good as the ones you didn't.

For a full sort, you have slightly more control. Some sorting algorithms are *stable*. In a stable sort, if two elements in a range have equivalent values, their relative positions are unchanged after sorting. Hence, if Widget A precedes Widget B in the (unsorted) widgets vector and both have the same quality rating, a stable sorting algorithm will guarantee that after the vector is sorted, Widget A still precedes Widget B. An algorithm that is not stable would not make this guarantee.

`partial_sort` is not stable. Neither is `nth_element`. `sort`, too, fails to offer stability, but there is an algorithm, `stable_sort`, that does what its name suggests. If you need stability when you sort, you'll probably want to use `stable_sort`. The STL does not contain stable versions of `partial_sort` or `nth_element`.

Speaking of `nth_element`, this curiously named algorithm is remarkably versatile. In addition to letting you find the top n elements of a range, it can also be used to find the median value in a range or to find the value at a particular percentile:

```

vector<Widget>::iterator begin(widgets.begin());    // convenience vars
vector<Widget>::iterator end(widgets.end());        // for widgets' begin
                                                    // and end iterators

vector<Widget>::iterator goalPosition;              // iter indicating where
                                                    // the widget of interest
                                                    // is located

// The following code finds the Widget with
// the median level of quality
goalPosition = begin + widgets.size() / 2;          // the widget of interest
                                                    // would be in the middle
                                                    // of the sorted vector

nth_element(begin, goalPosition, end,              // find the median quality
            qualityCompare);                       // value in widgets

...                                                  // goalPosition now points
                                                    // to the Widget with a
                                                    // median level of quality

// The following code finds the Widget with
// a level of quality at the 75th percentile
vector<Widget>::size_type goalOffset =              // figure out how far from
    0.25 * widgets.size();                         // the beginning the
                                                    // Widget of interest is

nth_element(begin, begin + goalOffset, end,         // find the quality value at
            qualityCompare);                       // the 75th percentile

...                                                  // begin+goalOffset now
                                                    // points to the Widget
                                                    // with the 75th percentile
                                                    // level of quality

```

`sort`, `stable_sort`, and `partial_sort` are great if you really need to put things in order, and `nth_element` fills the bill when you need to identify the top n elements or the element at a particular position, but sometimes you need something similar to `nth_element`, but not quite the same. Suppose, for example, you didn't need to identify the 20 highest-quality Widgets. Instead, you needed to identify all the Widgets with a quality rating of 1 or 2. You could, of course, sort the vector by quality and then search for the first one with a quality rating worse than 2. That would identify the beginning of a range of poor-quality Widgets.

A full sort can be a lot of work, however, and that much work is not necessary for this job. A better strategy is to use the partition algorithm, which reorders elements in a range so that all elements satisfying a particular criterion are at the beginning of the range.

For example, to move all the Widgets with a quality rating of 2 or better to the front of widgets, we define a function that identifies which Widgets make the grade,

```
bool hasAcceptableQuality(const Widget& w)
{
    // return whether w has a quality rating of 2 or better;
}
```

then pass that function to partition:

```
vector<Widget>::iterator goodEnd =      // move all widgets satisfying
    partition(widgets.begin(),          // hasAcceptableQuality to
              widgets.end(),            // the front of widgets, and
              hasAcceptableQuality);    // return an iterator to the first
                                         // widget that isn't satisfactory
```

After this call, the range from `widgets.begin()` to `goodEnd` holds all the Widgets with a quality of 1 or 2, and the range from `goodEnd` to `widgets.end()` contains all the Widgets with lower quality ratings. If it were important to maintain the relative positions of Widgets with equivalent quality levels during the partitioning, we'd naturally reach for `stable_partition` instead of `partition`.

The algorithms `sort`, `stable_sort`, `partial_sort`, and `nth_element` require random access iterators, so they may be applied only to vectors, strings, deques, and arrays. It makes no sense to sort elements in standard associative containers, because such containers use their comparison functions to remain sorted at all times. The only container where we might like to use `sort`, `stable_sort`, `partial_sort`, or `nth_element`, but can't, is `list`, and `list` compensates somewhat by offering its `sort` member function. (Interestingly, `list::sort` performs a stable sort.) If you want to sort a list, then, you can, but if you want to use `partial_sort`, or `nth_element` on the objects in a list, you have to do it indirectly. One indirect approach is to copy the elements into a container with random access iterators, then apply the desired algorithm to that. Another is to create a container of `list::iterators`, use the algorithm on that container, then access the list elements via the iterators. A third is to use the information in an ordered container of iterators to iteratively splice the list's elements into the positions you'd like them to be in. As you can see, there are lots of options.

`partition` and `stable_partition` differ from `sort`, `stable_sort`, `partial_sort`, and `nth_element` in requiring only bidirectional iterators. You can therefore use `partition` and `stable_partition` with any of the standard sequence containers.

Let's summarize your sorting options.

- If you need to perform a full sort on a vector, string, deque, or array, you can use `sort` or `stable_sort`.
- If you have a vector, string, deque, or array and you need to put only the top n elements in order, `partial_sort` is available.
- If you have a vector, string, deque, or array and you need to identify the element at position n or you need to identify the top n elements without putting them in order, `nth_element` is at your beck and call.
- If you need to separate the elements of a standard sequence container or an array into those that do and do not satisfy some criterion, you're probably looking for `partition` or `stable_partition`.
- If your data is in a list, you can use `partition` and `stable_partition` directly, and you can use `list::sort` in place of `sort` and `stable_sort`. If you need the effects offered by `partial_sort` or `nth_element`, you'll have to approach the task indirectly, but there are a number of options, as I sketched above.

In addition, you can keep things sorted at all times by storing your data in a standard associative container. You might also consider the standard non-STL container `priority_queue`, which also keeps its elements ordered all the time. (`priority_queue` is traditionally considered part of the STL, but, as I noted in the Introduction, my definition of “the STL” requires that STL containers support iterators, and `priority_queue` doesn't do iterators.)

“But what about performance?”, you wonder. Excellent question. Broadly speaking, algorithms that do more work take longer to do it, and algorithms that must sort stably take longer than algorithms that can ignore stability. We can order the algorithms we've discussed in this Item as follows, with algorithms that tend to use fewer resources (time and space) listed before those that require more:

- | | |
|----------------------------------|------------------------------|
| 1. <code>partition</code> | 4. <code>partial_sort</code> |
| 2. <code>stable_partition</code> | 5. <code>sort</code> |
| 3. <code>nth_element</code> | 6. <code>stable_sort</code> |

My advice on choosing among the sorting algorithms is to make your selection based on what you need to accomplish, not on performance considerations. If you choose an algorithm that does only what you need to do (e.g., a `partition` instead of a full sort), you're likely to end up with code that's not only the clearest expression of what you want to do, it's also the most efficient way to accomplish it using the STL.

Item 32: Follow remove-like algorithms by erase if you really want to remove something.

I begin this Item with a review of `remove`, because `remove` is the most confusing algorithm in the STL. Misunderstanding `remove` is easy, and it's important to dispel all doubt about what `remove` does, why it does it, and how it goes about doing it.

Here is the declaration for `remove`:

```
template<class ForwardIterator, class T>
ForwardIterator remove(ForwardIterator first, ForwardIterator last,
                      const T& value);
```

Like all algorithms, `remove` receives a pair of iterators identifying the range of elements over which it should operate. It does not receive a container, so `remove` doesn't know which container holds the elements it's looking at. Furthermore, it's not possible for `remove` to discover that container, because there is no way to go from an iterator to the container corresponding to that iterator.

Think for a moment about how one eliminates elements from a container. The only way to do it is to call a member function on that container, almost always some form of `erase`. (`list` has a couple of member functions that eliminate elements and are not named `erase`, but they're still member functions.) Because the only way to eliminate an element from a container is to invoke a member function on that container, and because `remove` cannot know the container holding the elements on which it is operating, *it is not possible for `remove` to eliminate elements from a container*. That explains the otherwise baffling observation that removing elements from a container never changes the number of elements in the container:

```
vector<int> v;                // create a vector<int> and fill it with
v.reserve(10);               // the values 1-10. (See Item 14 for an
for (int i = 1; i <= 10; ++i) { // explanation of the reserve call.)
    v.push_back(i);
}

cout << v.size();             // prints 10
v[3] = v[5] = v[9] = 99;      // set 3 elements to 99
remove(v.begin(), v.end(), 99); // remove all elements with value 99
cout << v.size();             // still prints 10!
```

To make sense of this example, memorize the following:

`remove` doesn't "really" remove anything, because it can't.

Repetition is good for you:

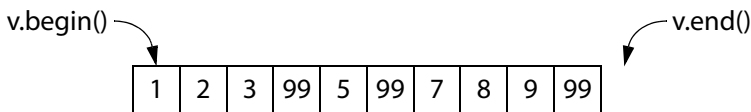
remove doesn't "really" remove anything, *because it can't*.

remove doesn't know the container it's supposed to remove things from, and without that container, there's no way for it to call the member functions that are necessary if one is to "really" remove something.

That explains what remove doesn't do, and it explains why it doesn't do it. What we need to review now is what remove *does* do.

Very briefly, remove moves elements in the range it's given until all the "unremoved" elements are at the front of the range (in the same relative order they were in originally). It returns an iterator pointing one past the last "unremoved" element. This return value is the "new logical end" of the range.

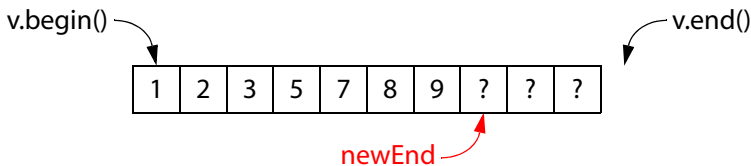
In terms of our example, this is what `v` looks like prior to calling remove,



and if we store remove's return value in a new iterator called `newEnd`,

```
vector<int>::iterator newEnd(remove(v.begin(), v.end(), 99));
```

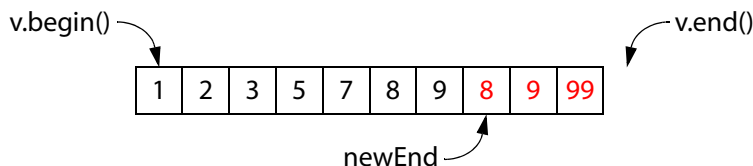
this is what `v` looks like after the call:



Here I've used question marks to indicate the values of elements that have been conceptually removed from `v` but continue to exist.

It seems logical that if the "unremoved" elements are in `v` between `v.begin()` and `newEnd`, the "removed" elements must be between `newEnd` and `v.end()`. *This is not the case!* The "removed" values aren't necessarily in `v` any longer at all. remove doesn't change the order of the elements in a range so that all the "removed" ones are at the end, it arranges for all the "unremoved" values to be at the beginning. Though the Standard doesn't require it, the elements beyond the new

logical end of the range typically *retain their old values*. After calling `remove`, `v` looks like this in every implementation I know:



As you can see, two of the “99” values that used to exist in `v` are no longer there, while one “99” remains. In general, after calling `remove`, the values removed from the range may or may not continue to exist in the range. Most people find this surprising, but why? You asked `remove` to get rid of some values, so it did. You didn’t ask it to put the removed values in a special place where you could get at them later, so it didn’t. What’s the problem? (If you don’t want to lose any values, you should probably be calling `partition` or `stable_partition` instead of `remove`. `partition` and `stable_partition` are described in [Item 31](#).)

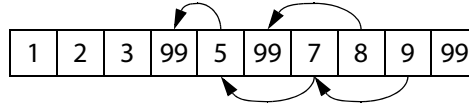
`remove`’s behavior sounds spiteful, but it’s simply a fallout of the way the algorithm operates. Internally, `remove` walks down the range, overwriting values that are to be “removed” with later values that are to be retained. The overwriting is accomplished by making assignments to the elements holding the values to be overwritten.

You can think of `remove` as performing a kind of compaction, where the values to be removed play the role of holes that are filled during compaction. For our vector `v`, it plays out as follows.

1. `remove` examines `v[0]`, sees that its value isn’t supposed to be removed, and moves on to `v[1]`. It does the same for `v[1]` and `v[2]`.
2. It sees that `v[3]` should be removed, so it notes that `v[3]`’s value may be overwritten, and it moves on to `v[4]`. This is akin to noting that `v[3]` is a “hole” that needs to be filled.
3. It sees that `v[4]`’s value should be retained, so it assigns `v[4]` to `v[3]`, notes that `v[4]` may now be overwritten, and moves on to `v[5]`. Continuing the compaction analogy, it “fills” `v[3]` with `v[4]` and notes that `v[4]` is now a hole.
4. It finds that `v[5]` should be removed, so it ignores it and moves on to `v[6]`. It continues to remember that `v[4]` is a hole waiting to be filled.
5. It sees that `v[6]` is a value that should be kept, so it assigns `v[6]` to `v[4]`, remembers that `v[5]` is now the next hole to be filled, and moves on to `v[7]`.

6. In a manner analogous to the above, it examines `v[7]`, `v[8]` and `v[9]`. It assigns `v[7]` to `v[5]` and `v[8]` to `v[6]`, ignoring `v[9]`, because the value at `v[9]` is to be removed.
7. It returns an iterator indicating the next element to be overwritten, in this case the element at `v[7]`.

You can envision the values moving around in `v` as follows:



As [Item 33](#) explains, the fact that `remove` overwrites some of the values it is removing has important repercussions when those values are pointers. For this Item, however, it's enough to understand that `remove` doesn't eliminate any elements from a container, because it can't. Only container member functions can eliminate container elements, and that's the whole point of this Item: You should follow `remove` by `erase` if you really want to remove something.

The elements you want to erase are easy to identify. They're the elements of the original range that start at the “new logical end” of the range and continue until the real end of the range. To get rid of these elements, all you need to do is call the range form of `erase` (see [Item 5](#)) with these two iterators. Because `remove` itself conveniently returns the iterator for the new logical end of the range, the call is straightforward:

```
vector<int> v;           // as before
...
v.erase(remove(v.begin(), v.end(), 99), v.end()); // really remove all
                                                    // elements with value 99
cout << v.size();       // now returns 7
```

Passing `remove`'s return value as the first argument to the range form of `erase` is so common, it's idiomatic. In fact, `remove` and `erase` are so closely allied, the two are merged in the list member function `remove`. This is the only function in the STL named `remove` that eliminates elements from a container:

```
list<int> li;           // create a list
...                     // put some values into it
li.remove(99);          // eliminate all elements with value 99;
                        // this really removes elements, so li's
                        // size may change
```

Frankly, calling this function `remove` is an inconsistency in the STL. The analogous function in the associative containers is called `erase`, and `list`'s `remove` should be called `erase`, too. It's not, however, so we all have to get used to it. The world in which we frolic may not be the best of all possible worlds, but it is the one we've got. (On the plus side, [Item 44](#) points out that, for lists, calling the `remove` member function is more efficient than applying the `erase-remove` idiom.)

Once you understand that `remove` can't "really" remove things from a container, using it in conjunction with `erase` becomes second nature. The only other thing you need to bear in mind is that `remove` isn't the only algorithm for which this is the case. There are two other "remove-like" algorithms: `remove_if` and `unique`.

The similarity between `remove` and `remove_if` is so straightforward, I won't dwell on it, but `unique` also behaves like `remove`. It is asked to remove things (adjacent repeated values) from a range without having access to the container holding the range's elements. As a result, you must also pair calls to `unique` with calls to `erase` if you really want to remove elements from a container. `unique` is also analogous to `remove` in its interaction with `list`. Just as `list::remove` really removes things (and does so more efficiently than the `erase-remove` idiom), `list::unique` really removes adjacent duplicates (also with greater efficiency than would `erase-unique`).

Item 33: Be wary of remove-like algorithms on containers of pointers.

So you conjure up a bunch of dynamically allocated `Widgets`, each of which may be certified, and you store the resulting pointers in a vector:

```
class Widget {
public:
    ...
    bool isCertified() const;           // whether the Widget is certified
    ...
};
vector<Widget*> v;                      // create a vector and fill it with
...                                   // pointers to dynamically
v.push_back(new Widget);               // allocated Widgets
...
```

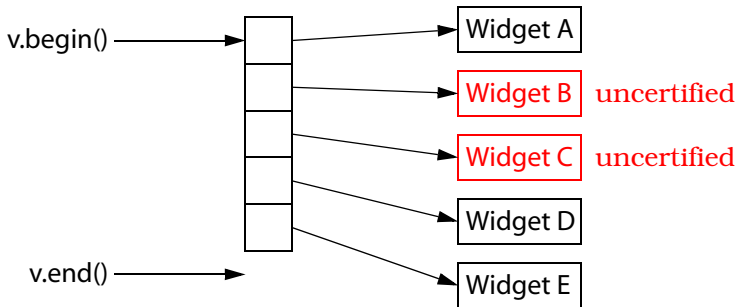
After working with `v` for a while, you decide to get rid of the uncertified `Widgets`, because you don't need them any longer. Bearing in mind [Item 43](#)'s admonition to prefer algorithm calls to explicit loops and

having read [Item 32](#)'s discourse on the relationship between `remove` and `erase`, your thoughts naturally turn to the `erase-remove` idiom, though in this case it's `remove_if` you employ:

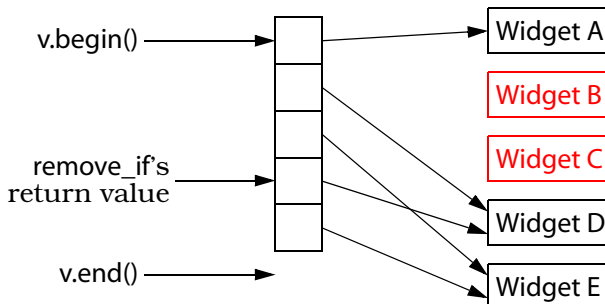
```
v.erase(remove_if(v.begin(), v.end(),           // erase ptrs to
                  not1(mem_fun(&Widget::isCertified))), // uncertified
        v.end());                               // Widgets; see
                                                // Item 41 for
                                                // info on
                                                // mem_fun
```

Suddenly you begin to worry about the call to `erase`, because you dimly recall [Item 7](#)'s discussion of how destroying a pointer in a container fails to delete what the pointer points to. This is a legitimate worry, but in this case, it comes too late. By the time `erase` is called, there's an excellent chance you have already leaked resources. Worry about `erase`, yes, but first, worry about `remove_if`.

Let's assume that prior to the `remove_if` call, `v` looks like this, where I've indicated the uncertified Widgets.



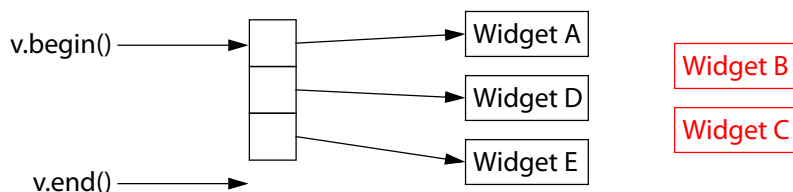
After the call to `remove_if`, `v` will typically look like this (including the iterator returned from `remove_if`):



If this transformation makes no sense, kindly turn to [Item 32](#), because it explains exactly what happens when `remove` — or, in this case, `remove_if` — is called.

The reason for the resource leak should now be apparent. The “removed” pointers to Widgets B and C have been overwritten by later “unremoved” pointers in the vector. Nothing points to the two uncertified Widgets, they can never be deleted, and their memory and other resources are leaked.

Once both `remove_if` and `erase` have returned, the situations looks as follows:



This makes the resource leak especially obvious, and it should now be clear why you should try to avoid using `remove` and similar algorithms (i.e., `remove_if` and `unique`) on containers of pointers to dynamically allocated objects. In many cases, you’ll find that the `partition` algorithm (see [Item 31](#)) is a reasonable alternative.

If you can’t avoid using `remove` on such containers, one way to eliminate this problem is to delete the pointers and set them to null prior to applying the `erase-remove` idiom, then eliminate all the null pointers in the container:

```

void delAndNullifyUncertified(Widget*& pWidget) // if *pWidget is an
{                                              // uncertified Widget,
    if (!pWidget->isCertified()) {           // delete the pointer
        delete pWidget;                     // and set it to null
        pWidget = 0;
    }
}

for_each(v.begin(), v.end(),                 // delete and set to
        delAndNullifyUncertified);           // null all ptrs to
                                              // uncertified widgets

v.erase( remove(v.begin(), v.end(),          // eliminate null ptrs
               static_cast<Widget*>(0)),     // from v; 0 must be
         v.end());                          // cast to a ptr so C++
                                              // correctly deduces
                                              // the type of
                                              // remove's 3rd param
  
```

Of course, this assumes that the vector doesn’t hold any null pointers you’d like to retain. If it does, you’ll probably have to write your own loop that erases pointers as you go. Erasing elements from a container

as you traverse that container has some subtle aspects to it, so be sure to read [Item 9](#) before considering that approach.

If you're willing to replace the container of pointers with a container of *smart* pointers that perform reference counting, the remove-related difficulties wash away, and you can use the erase-remove idiom directly:

```

template<typename T>                // RCSP = "Reference Counting
class RCSP { ... };                 //      Smart Pointer"

typedef RCSP<Widget> RCSPW;         // RCSPW = "RCSP to Widget"

vector<RCSPW> v;                    // create a vector and fill it with
...                                // smart pointers to dynamically
v.push_back(RCSPW(new Widget));    // allocated Widgets
...

v.erase(remove_if(v.begin(), v.end(), // erase the ptrs
                  not1(mem_fun(&Widget::isCertified))), // to uncertified
         v.end());                  // Widgets; no
                                   // resources are
                                   // leaked
```

For this to work, it must be possible to implicitly convert your smart pointer type (e.g., `RCSP<Widget>`) to the corresponding built-in pointer type (e.g., `Widget*`). That's because the container holds smart pointers, but the member function being called (e.g., `Widget::isCertified`) insists on built-in pointers. If no implicit conversion exists, your compilers will squawk.

If you don't happen to have a reference counting smart pointer template in your programming toolbox, you owe it to yourself to check out the `shared_ptr` template in the Boost library. For an introduction to Boost, take a look at [Item 50](#).

Regardless of how you choose to deal with containers of dynamically allocated pointers, be it by reference counting smart pointers, manual deletion and nullification of pointers prior to invoking a remove-like algorithm, or some technique of your own invention, the guidance of this Item remains the same: Be wary of remove-like algorithms on containers of pointers. Failure to heed this advice is just *asking* for resource leaks.

Item 34: Note which algorithms expect sorted ranges.

Not all algorithms are applicable to all ranges. For example, `remove` (see [Items 32](#) and [33](#)) requires forward iterators and the ability to make assignments through those iterators. As a result, it can't be applied to ranges demarcated by input iterators, nor to maps or multi-

maps, nor to some implementations of set and multiset (see [Item 22](#)). Similarly, many of the sorting algorithms (see [Item 31](#)) demand random access iterators, so it's not possible to invoke these algorithms on the elements of a list.

If you violate these kinds of rules, your code won't compile, an event likely to be heralded by lengthy and incomprehensible error messages (see [Item 49](#)). Other algorithm preconditions, however, are more subtle. Among these, perhaps the most common is that some algorithms require ranges of *sorted* values. It's important that you adhere to this requirement whenever it applies, because violating it leads not to compiler diagnostics, but to undefined runtime behavior.

A few algorithms can work with sorted or unsorted ranges, but they are most useful when they operate on sorted ranges. You should understand how these algorithms work, because that will explain why sorted ranges suit them best.

Some of you, I know, are into brute-force memorization, so here's a list of the algorithms that require the data on which they operate to be sorted:

<code>binary_search</code>	<code>lower_bound</code>
<code>upper_bound</code>	<code>equal_range</code>
<code>set_union</code>	<code>set_intersection</code>
<code>set_difference</code>	<code>set_symmetric_difference</code>
<code>merge</code>	<code>inplace_merge</code>
<code>includes</code>	

In addition, the following algorithms are typically used with sorted ranges, though they don't require them:

<code>unique</code>	<code>unique_copy</code>
---------------------	--------------------------

We'll see shortly that the definition of "sorted" has an important constraint, but first, let me try to make sense of this collection of algorithms. It's easier to remember which algorithms work with sorted ranges if you understand why such ranges are needed.

The search algorithms `binary_search`, `lower_bound`, `upper_bound`, and `equal_range` (see [Item 45](#)) require sorted ranges, because they look for values using binary search. Like the C library's `bsearch`, these algorithms promise logarithmic-time lookups, but in exchange, you must give them values that have already been put into order.

Actually, it's not quite true that these algorithms promise logarithmic-time lookup. They guarantee such performance only when they are passed random access iterators. If they're given less powerful iterators

(such as bidirectional iterators), they still perform only a logarithmic number of comparisons, but they run in linear time. That's because, lacking the ability to perform "iterator arithmetic," they need linear time to move from place to place in the range being searched.

The quartet of algorithms `set_union`, `set_intersection`, `set_difference`, and `set_symmetric_difference` offer linear-time performance of the set-theoretical operations their names suggest. Why do they demand sorted ranges? Because without them, they couldn't do their work in linear time. If you're beginning to detect a trend suggesting that algorithms requiring sorted ranges do so in order to offer better performance than they'd be able to guarantee for ranges that might not be sorted, you're right. Stay tuned. The trend will continue.

`merge` and `inplace_merge` perform what is in effect a single pass of the mergesort algorithm: they read two sorted ranges and produce a new sorted range containing all the elements from both source ranges. They run in linear time, something they couldn't do if they didn't know that the source ranges were already sorted.

The final algorithm that requires sorted source ranges is `includes`. It's used to determine whether all the objects in one range are also in another range. Because `includes` may assume that both its ranges are sorted, it promises linear-time performance. Without that guarantee, it would generally run slower.

Unlike the algorithms we've just discussed, `unique` and `unique_copy` offer well-defined behavior even on unsorted ranges. But look at how the Standard describes `unique`'s behavior (the italics are mine):

Eliminates all but the first element from every *consecutive* group of equal elements.

In other words, if you want `unique` to eliminate all duplicates from a range (i.e., to make all values in the range "unique"), you must first make sure that all duplicate values are next to one another. And guess what? That's one of the things sorting does. In practice, `unique` is usually employed to eliminate all duplicate values from a range, so you'll almost always want to make sure that the range you pass `unique` (or `unique_copy`) is sorted. (Unix developers will recognize a striking similarity between STL's `unique` and Unix's `uniq`, a similarity I suspect is anything but coincidental.)

Incidentally, `unique` eliminates elements from a range the same way `remove` does, which is to say that it only "sort of" eliminates them. If you aren't sure what this means, please turn to Items 32 and 33 immediately. It is not possible to overemphasize the importance of

understanding what `remove` and `remove-like` algorithms (including `unique`) do. Having a basic comprehension is not sufficient. If you don't *know* what they do, you will get into trouble.

Which brings me to the fine print regarding what it means for a range to be sorted. Because the STL allows you to specify comparison functions to be used during sorting, different ranges may be sorted in different ways. Given two ranges of ints, for example, one might be sorted the default way (i.e., in ascending order) while the other is sorted using `greater<int>`, hence in descending order. Given two ranges of `Widgets`, one might be sorted by price and another might be sorted by age. With so many different ways to sort things, it's critical that you give the STL consistent sorting-related information to work with. If you pass a sorted range to an algorithm that also takes a comparison function, be sure that the comparison function you pass behaves the same as the one you used to sort the range.

Here's an example of what you do *not* want to do:

```
vector<int> v;           // create a vector, put some
...                     // data into it, sort it into
sort(v.begin(), v.end(), greater<int>()); // descending order
...
// work with the vector
// (without changing it)

bool a5Exists =         // search for a 5 in the vector,
    binary_search(v.begin(), v.end(), 5); // assuming it's sorted in
                                           // ascending order!
```

By default, `binary_search` assumes that the range it's searching is sorted by “<” (i.e., the values are in ascending order), but in this example, the vector is sorted in *descending* order. You should not be surprised to learn that you get undefined results when you invoke `binary_search` (or `lower_bound`, etc.) on a range of values that is sorted in a different order from what the algorithm expects.

To get the code to behave correctly, you must tell `binary_search` to use the same comparison function that `sort` did:

```
bool a5Exists =         // search for a 5
    binary_search(v.begin(), v.end(), 5, greater<int>()); // using greater as
                                                         // the comparison
                                                         // function
```

All the algorithms that require sorted ranges (i.e., all the algorithms in this Item except `unique` and `unique_copy`) determine whether two values are “the same” by using equivalence, just like the standard associative containers (which are themselves sorted). In contrast, the default way in which `unique` and `unique_copy` determine whether two

objects are “the same” is by using equality, though you can override this default by passing these algorithms a predicate defining an alternative definition of “the same.” For a detailed discussion of the difference between equivalence and equality, consult [Item 19](#).

The eleven algorithms that require sorted ranges do so in order to offer greater efficiency than would otherwise be possible. As long as you remember to pass them only sorted ranges, and as long as you make sure that the comparison function used by the algorithms is consistent with the one used to do the sorting, you’ll revel in trouble-free search, set, and merge operations, plus you’ll find that `unique` and `unique_copy` eliminate *all* duplicate values, as you almost certainly want them to.

Item 35: Implement simple case-insensitive string comparisons via `mismatch` or `lexicographical_compare`.

One of the most frequently asked questions by STL newbies is “How do I use the STL to perform case-insensitive string comparisons?” This is a deceptively simple question. Case-insensitive string comparisons are either really easy or really hard, depending on how general you want to be. If you’re willing to ignore internationalization issues and restrict your concern to the kinds of strings `strcmp` is designed for, the task is easy. If you want to be able to handle strings of characters in languages where `strcmp` wouldn’t apply (i.e., strings holding text in just about any language except English) or where programs use a locale other than the default, the task is very hard.

In this Item, I’ll tackle the easy version of the problem, because that suffices to demonstrate how the STL can be brought to bear. (A harder version of the problem involves no more of the STL. Rather, it involves locale-dependent issues you can read about in [Appendix A](#).) To make the easy problem somewhat more challenging, I’ll tackle it twice. Programmers desiring case-insensitive string comparisons often need two different calling interfaces, one similar to `strcmp` (which returns a negative number, zero, or a positive number), the other akin to `operator<` (which returns true or false). I’ll therefore show how to implement both calling interfaces using STL algorithms.

First, however, we need a way to determine whether two characters are the same, except for their case. When internationalization issues are taken into account, this is a complicated problem. The following character-comparing function is a simplistic solution, but it’s akin to

the `strcmp` approach to string comparison, and since in this Item I consider only strings where a `strcmp`-like approach is appropriate, internationalization issues don't count, and this function will do:

```
int ciCharCompare(char c1, char c2)    // case-insensitively compare chars
{                                     // c1 and c2, returning -1 if c1 < c2,
                                     // 0 if c1==c2, and 1 if c1 > c2
    int lc1 = tolower(static_cast<unsigned char>(c1));    // see below for
    int lc2 = tolower(static_cast<unsigned char>(c2));    // info on these
                                                         // statements

    if (lc1 < lc2) return -1;
    if (lc1 > lc2) return 1;
    return 0;
}
```

This function follows the lead of `strcmp` in returning a negative number, zero, or a positive number, depending on the relationship between `c1` and `c2`. Unlike `strcmp`, `ciCharCompare` converts both parameters to lower case before performing the comparison. That's what makes it a case-insensitive character comparison.

Like many functions in `<cctype>` (and hence `<ctype.h>`), `tolower`'s parameter and return value is of type `int`, but unless that `int` is EOF, its value must be representable as an unsigned char. In both C and C++, `char` may or may not be signed (it's up to the implementation), and when `char` is signed, the only way to ensure that its value is representable as an unsigned char is to cast it to one before calling `tolower`. That explains the casts in the code above. (On implementations where `char` is already unsigned, the casts are no-ops.) It also explains the use of `int` instead of `char` to store `tolower`'s return value.

Given `ciCharCompare`, it's easy to write the first of our two case-insensitive string comparison functions, the one offering a `strcmp`-like interface. This function, `ciStringCompare`, returns a negative number, zero, or a positive number, depending on the relationship between the strings being compared. It's built around the mismatch algorithm, because mismatch identifies the first position in two ranges where the corresponding values are not the same.

Before we can call `mismatch`, we have to satisfy its preconditions. In particular, we have to make sure that if one string is shorter than the other, the shorter string is the first range passed. We'll therefore farm the real work out to a function called `ciStringCompareImpl` and have `ciStringCompare` simply make sure the arguments are passed in the cor-

correct order, adjusting the return value if the arguments have to be swapped:

```
int ciStringCompareImpl(const string& s1,           // see below for
                      const string& s2);          // implementation

int ciStringCompare(const string& s1, const string& s2)
{
    if (s1.size() <= s2.size()) return ciStringCompareImpl(s1, s2);
    else return -ciStringCompareImpl(s2, s1);
}
```

In `ciStringCompareImpl`, the heavy lifting is performed by `mismatch`. It returns a pair of iterators indicating the locations in the ranges where corresponding characters first fail to match:

```
int ciStringCompareImpl(const string& s1, const string& s2)
{
    typedef pair<string::const_iterator,           // PSCI = "pair of
               string::const_iterator> PSCI;      // string::const_iterator"

    PSCI p = mismatch(                            // see below for an
        s1.begin(), s1.end(),                     // explanation of why
        s2.begin(),                               // we need not2; see
        not2(ptr_fun(ciCharCompare)));            // Item 41 for why we
                                                // need ptr_fun

    if (p.first == s1.end()) {                    // if true, either s1 and
        if (p.second == s2.end()) return 0;       // s2 are equal or
        else return -1;                          // s1 is shorter than s2
    }

    return ciCharCompare(*p.first, *p.second);    // the relationship of the
                                                // strings is the same as
                                                // that of the
                                                // mismatched chars
}
```

With any luck, the comments make pretty clear what is going on. Fundamentally, once you know the first place where the strings differ, it's easy to determine which string, if either, precedes the other. The only thing that may seem odd is the predicate passed to `mismatch`, which is `not2(ptr_fun(ciCharCompare))`. This predicate is responsible for returning true when the characters match, because `mismatch` will stop when the predicate returns false. We can't use `ciCharCompare` for this purpose, because it returns -1, 1, or 0, and *it returns 0 when the characters match*, just like `strcmp`. If we passed `ciCharCompare` as the predicate to `mismatch`, C++ would convert `ciCharCompare`'s return type to `bool`, and of course the `bool` equivalent of zero is false, precisely the opposite of what we want! Similarly, when `ciCharCompare` returned 1 or -1, that would be interpreted as true, because, as in C, all nonzero integral values are considered true. Again, this would be the opposite of what we

want. To fix this semantic inversion, we throw a `not2` and a `ptr_fun` in front of `ciCharCompare`, and we all live happily ever after.

Our second approach to `ciStringCompare` yields a conventional STL predicate; such a function could be used as a comparison function in associative containers. The implementation is short and sweet, because all we have to do is modify `ciCharCompare` to give us a character-comparison function with a predicate interface, then turn the job of performing a string comparison over to the algorithm with the second-longest name in the STL, `lexicographical_compare`:

```
bool ciCharLess(char c1, char c2)           // return whether c1
{                                           // precedes c2 in a case-
    return                                 // insensitive comparison;
    tolower(static_cast<unsigned char>(c1)) < // Item 46 explains why a
    tolower(static_cast<unsigned char>(c2)); // function object might
}                                           // be preferable to this
                                           // function

bool ciStringCompare(const string& s1, const string& s2)
{
    return lexicographical_compare(s1.begin(), s1.end(), // see below for
                                   s2.begin(), s2.end(), // a discussion of
                                   ciCharLess);          // this algorithm
}                                                         // call
```

No, I won't keep you in suspense any longer. The longest algorithm name is `set_symmetric_difference`.

If you're familiar with the behavior of `lexicographical_compare`, the code above is as clear as clear can be. If you're not, it's probably about as clear as concrete. Fortunately, it's not hard to replace the concrete with glass.

`lexicographical_compare` is a generalized version of `strcmp`. Where `strcmp` works only with character arrays, however, `lexicographical_compare` works with ranges of values of any type. Also, while `strcmp` always compares two characters to see if their relationship is equal, less than, or greater than one another, `lexicographical_compare` may be passed an arbitrary predicate that determines whether two values satisfy a user-defined criterion.

In the call above, `lexicographical_compare` is asked to find the first position where `s1` and `s2` differ, based on the results of calls to `ciCharLess`. If, using the characters at that position, `ciCharLess` returns true, so does `lexicographical_compare`: if, at the first position where the characters differ, the character from the first string precedes the corresponding character from the second string, the first string precedes the second one. Like `strcmp`, `lexicographical_compare` considers two ranges

of equal values to be equal, hence it returns false for such ranges: the first range does *not* precede the second. Also like `strcmp`, if the first range ends before a difference in corresponding values is found, `lexicographical_compare` returns true: a prefix precedes any range for which it is a prefix.

Enough about `mismatch` and `lexicographical_compare`. Though I focus on portability in this book, I would be remiss if I failed to mention that case-insensitive string comparison functions are widely available as nonstandard extensions to the standard C library. They typically have names like `stricmp` or `strcmpi`, and they typically offer no more support for internationalization than the functions we've developed in this Item. If you're willing to sacrifice some portability, you know that your strings never contain embedded nulls, and you don't care about internationalization, you may find that the easiest way to implement a case-insensitive string comparison doesn't use the STL at all. Instead, it converts both strings to `const char*` pointers (see [Item 16](#)) and then uses `stricmp` or `strcmpi` on the pointers:

```
int ciStringCompare(const string& s1, const string& s2)
{
    return stricmp(s1.c_str(), s2.c_str());           // the function name on
}                                                       // your system might
                                                       // not be stricmp
```

Some may call this a hack, but `stricmp/strcmpi`, being optimized to do exactly one thing, typically run *much* faster on long strings than do the general-purpose algorithms `mismatch` and `lexicographical_compare`. If that's an important consideration for you, you may not care that you're trading standard STL algorithms for nonstandard C functions. Sometimes the most effective way to use the STL is to realize that other approaches are superior.

Item 36: Understand the proper implementation of `copy_if`.

One of the more interesting aspects of the STL is that although there are 11 algorithms with “copy” in their names,

<code>copy</code>	<code>copy_backward</code>
<code>replace_copy</code>	<code>reverse_copy</code>
<code>replace_copy_if</code>	<code>unique_copy</code>
<code>remove_copy</code>	<code>rotate_copy</code>
<code>remove_copy_if</code>	<code>partial_sort_copy</code>
<code>uninitialized_copy</code>	

none of them is `copy_if`. That means you can `replace_copy_if`, you can `remove_copy_if`, you can both `copy_backward` and `reverse_copy`, but if

you simply want to copy the elements of a range that satisfy a predicate, you're on your own.

For example, suppose you have a function to determine whether a Widget is defective:

```
bool isDefective(const Widget& w);
```

and you'd like to write all the defective Widgets in a vector to cerr. If copy_if existed, you could simply do this:

```
vector<Widget> widgets;
...
copy_if(widgets.begin(), widgets.end(),           // this won't compile;
        ostream_iterator<Widget>(cerr, "\n"),    // there is no copy_if
        isDefective);                             // in the STL
```

Ironically, copy_if was part of the original Hewlett Packard STL that formed the basis for the STL that is now part of the standard C++ library. In one of those quirks that occasionally makes history interesting, during the process of winnowing the HP STL into something of a size manageable for standardization, copy_if was one of the things that got left on the cutting room floor.

In *The C++ Programming Language* [7], Stroustrup remarks that it's trivial to write copy_if, and he's right, but that doesn't mean that it's necessarily easy to come up with the correct trivia. For example, here's a reasonable-looking copy_if that many people (including me) have been known to come up with:

```
template< typename InputIterator,           // a not-quite-right
          typename OutputIterator,         // implementation of
          typename Predicate>              // copy_if
OutputIterator copy_if( InputIterator begin,
                       InputIterator end,
                       OutputIterator destBegin,
                       Predicate p)
{
    return remove_copy_if(begin, end, destBegin, not1(p));
}
```

This approach is based on the observation that although the STL doesn't let you say "copy everything where this predicate is true," it does let you say "copy everything except where this predicate is *not* true." To implement copy_if, then, it seems that all we need to do is throw a not1 in front of the predicate we'd like to pass to copy_if, then

pass the resulting predicate to `remove_copy_if`. The result is the code above.

If the above reasoning were valid, we could write out our defective `Widgets` this way:

```
copy_if(widgets.begin(), widgets.end(),           // well-intentioned code
        ostream_iterator<Widget>(cerr, "\n"),    // that will not compile
        isDefective);
```

Your STL platforms will take a jaundiced view of this code, because it tries to apply `not1` to `isDefective`. (The application takes place inside `copy_if`.) As [Item 41](#) tries to make clear, `not1` can't be applied directly to a function pointer; the function pointer must first be passed through `ptr_fun`. To call this implementation of `copy_if`, you must pass not just a function object, but an *adaptable* function object. That's easy enough to do, but clients of a would-be STL algorithm shouldn't have to. Standard STL algorithms never require that their functors be adaptable, and neither should `copy_if`. The above implementation is decent, but it's not decent enough.

Here's the correct trivial implementation of `copy_if`:

```
template< typename InputIterator,           // a correct
          typename OutputIterator,         // implementation of
          typename Predicate>              // copy_if
OutputIterator copy_if( InputIterator begin,
                       InputIterator end,
                       OutputIterator destBegin,
                       Predicate p)
{
    while (begin != end) {
        if (p(*begin)) *destBegin++ = *begin;
        ++begin;
    }
    return destBegin;
}
```

Given how useful `copy_if` is, plus the fact that new STL programmers tend to expect it to exist anyway, there's a good case to be made for putting `copy_if` — the correct one! — into your local STL-related utility library and using it whenever it's appropriate.

Item 37: Use `accumulate` or `for_each` to summarize ranges.

Sometimes you need to boil an entire range down to a single number, or, more generally, a single object. For commonly needed information, special-purpose algorithms exist to do the jobs. `count` tells you how

many elements with a particular value are in a range, for example, while `count_if` tells you how many elements satisfy a predicate. The minimum and maximum values in a range are available via `min_element` and `max_element`.

At times, however, you need to summarize a range in some custom manner, and in those cases, you need something more flexible than `count`, `count_if`, `min_element`, or `max_element`. For example, you might want the sum of the lengths of the strings in a container. You might want the product of a range of numbers. You might want the average coordinates of a range of points. In each of these cases, you need to *summarize* a range, but you need to be able to define the summary you want. Not a problem. The STL has the algorithm for you. It's called `accumulate`. You might not be familiar with `accumulate`, because, unlike most algorithms, it doesn't live in `<algorithm>`. Instead, it's located with three other "numeric algorithms" in `<numeric>`. (The three others are `inner_product`, `adjacent_difference`, and `partial_sum`.)

Like many algorithms, `accumulate` exists in two forms. The form taking a pair of iterators and an initial value returns the initial value plus the sum of the values in the range demarcated by the iterators:

```
list<double> ld;                // create a list and put
...                             // some doubles into it
double sum = accumulate(ld.begin(), ld.end(), 0.0); // calculate their sum,
                                                    // starting at 0.0
```

In this example, note that the initial value is specified as `0.0`, not simply `0`. That's important. The type of `0.0` is `double`, so `accumulate` internally uses a variable of type `double` to store the sum it's computing. Had the call been written like this,

```
double sum = accumulate(ld.begin(), ld.end(), 0); // calculate their sum,
                                                    // starting at 0; this
                                                    // is not correct!
```

the initial value would be the `int` `0`, so `accumulate` would internally use an `int` to store the value it was computing. That `int` would ultimately become `accumulate`'s return value, and it would be used to initialize the variable `sum`. The code would compile and run, but `sum`'s value would be incorrect. Instead of holding the true sum of a list of doubles, it would hold the result of adding all the doubles together, but converting the result to an `int` after each addition.

`accumulate` requires only input iterators, so you can use it even with `istream_iterators` and `istreambuf_iterators` (see [Item 29](#)):

```
cout << "The sum of the ints on the standard input is " // print the sum of
    << accumulate(istream_iterator<int>(cin),          // the ints in cin
                  istream_iterator<int>(),
                  0);
```

It is this default behavior of `accumulate` that causes it to be labeled a numeric algorithm. But when `accumulate` is used in its alternate form, one taking an initial summary value and an arbitrary summarization function, it becomes much more general.

As an example, consider how to use `accumulate` to calculate the sum of the lengths of the strings in a container. To compute the sum, `accumulate` needs to know two things. First, as above, it must know the starting sum. In our case, it is zero. Second, it must know how to update this sum each time a new string is seen. To do that, we write a function that takes the sum so far and the new string and returns the updated sum:

```
string::size_type          // see below for info
stringLengthSum(string::size_type sumSoFar, // on string::size_type
                 const string& s)
{
    return sumSoFar + s.size();
}
```

The body of this function reveals that what is going on is trivial, but you may find yourself bogged down in the appearances of `string::size_type`. Don't let that happen. Every standard STL container has a typedef called `size_type` that is the container's type for counting things. This is the type returned by the container's `size` function, for example. For all the standard containers, `size_type` must be `size_t`, but, in theory, nonstandard STL-compatible containers might use a different type for `size_type` (though I have a hard time imagining why they'd want to). For standard containers, you can think of `Container::size_type` as a fancy way of writing `size_t`.

`stringLengthSum` is representative of the summarization functions `accumulate` works with. It takes a summary value for the range so far as well as the next element of the range, and it returns the new summary value. In general, that means the function will take parameters of different types. That's what it does here. The summary so far (the sum of the lengths of the strings already seen) is of type `string::size_type`, while the type of the elements being examined is `string`. As is typically the case, the return type here is the same as that of the function's first parameter, because it's the updated summary value (the one taking the latest element into account).

We can use `stringLengthSum` with `accumulate` like this:

```
set<string> ss;          // create container of strings,
...                      // and populate it
```

```

string::size_type lengthSum =           // set lengthSum to the result
    accumulate(ss.begin(), ss.end(),    // of calling stringLengthSum on
        static_cast<string::size_type>(0), // each element in ss, using 0
        stringLengthSum);              // as the initial summary value

```

Nifty, huh? Calculating the product of a range of numbers is even easier, because we don't have to write our own summation function. We can use the standard multiplies functor class:

```

vector<float> vf;                        // create container of floats
...                                     // and populate it
float product =                         // set product to the result of
    accumulate(vf.begin(), vf.end(),    // calling multiplies<float> on
        1.0f, multiplies<float>());    // each element in vf, using 1.0f
// as the initial summary value

```

The only tricky thing here is remembering to use one (as a float, not as an int!) as the initial summary value instead of zero. If we used zero as the starting value, the result would always be zero, because zero times anything is zero, right?

Our final example is a bit more ambitious. It involves finding the average of a range of points, where a point looks like this:

```

struct Point {
    Point(double initX, double initY): x(initX), y(initY) {}
    double x, y;
};

```

The summation function will be an object of a functor class called PointAverage, but before we look at PointAverage, let's look at its use in the call to accumulate:

```

list<Point> lp;
...
Point avg =                             // average the points in lp
    accumulate(lp.begin(), lp.end(),
        Point(0, 0), PointAverage());

```

Simple and straightforward, the way we like it. In this case, the initial summary value is a Point object located at the origin, and all we need to remember is not to take that point into account when computing the average of the range.

PointAverage works by keeping track of the number of points it has seen, as well as the sum of their x and y components. Each time it is called, it updates these values and returns the average coordinates of the points so far examined. Because it is called exactly once for each point in the range, it divides the x and y sums by the number of points

in the range; the initial point value passed to `accumulate` is ignored, as it should be:

```
class PointAverage:
    public binary_function<Point, Point, Point> {           // see Item 40
    public:
        PointAverage(): numPoints(0), xSum(0), ySum(0) {}
        const Point operator()(const Point& avgSoFar, const Point& p)
        {
            ++numPoints;
            xSum += p.x;
            ySum += p.y;
            return Point(xSum/numPoints, ySum/numPoints);
        }
    private:
        size_t numPoints;
        double xSum;
        double ySum;
};
```

This works fine, and it is only because I sometimes associate with an inordinately demented group of people (many of them on the Standardization Committee) that I can envision STL implementations where it could fail. Nevertheless, `PointAverage` runs afoul of paragraph 2 of section 26.4.1 of the Standard, which, as I'm sure you recall, forbids side effects in the function passed to `accumulate`. Modification of the member variables `numPoints`, `xSum`, and `ySum` constitutes a side effect, so, technically speaking, the code I've just shown you yields undefined results. In practice, it's hard to imagine it not working, but I'm surrounded by menacing language lawyers as I write this, so I've no choice but to spell out the fine print on this matter.

That's okay, because it gives me a chance to mention `for_each`, another algorithm that can be used to summarize ranges and one that isn't constrained by the restrictions imposed on `accumulate`. Like `accumulate`, `for_each` takes a range and a function (typically a function object) to invoke on each element of the range, but the function passed to `for_each` receives only a single argument (the current range element), and `for_each` returns its function when it's done. (Actually, it returns a *copy* of its function — see [Item 38](#).) Significantly, the function passed to (and later returned from) `for_each` *may* have side effects.

Ignoring the side effects issue, `for_each` differs from `accumulate` in two primary ways. First, the name `accumulate` suggests an algorithm that produces a summary of a range. `for_each` sounds like you simply want to do something to every element of a range, and, of course, that is the

algorithm's primary application. Using `for_each` to summarize a range is legitimate, but it's not as clear as `accumulate`.

Second, `accumulate` returns the summary we want directly, while `for_each` returns a function object, and we must extract the summary information we want from this object. In C++ terms, that means we must add a member function to the functor class to let us retrieve the summary information we're after.

Here's the last example again, this time using `for_each` instead of `accumulate`:

```
struct Point { ... };                                // as before
class PointAverage:
public unary_function<Point, void> {                 // see Item 40
public:
    PointAverage(): numPoints(0), xSum(0), ySum(0) {}
    void operator()(const Point& p)
    {
        ++numPoints;
        xSum += p.x;
        ySum += p.y;
    }
    Point result() const
    {
        return Point(xSum/numPoints, ySum/numPoints);
    }
private:
    size_t numPoints;
    double xSum;
    double ySum;
};
list<Point> lp;
...
Point avg = for_each(lp.begin(), lp.end(), PointAverage()).result();
```

Personally, I prefer `accumulate` for summarizing, because I think it most clearly expresses what is going on, but `for_each` works, too, and the issue of side effects doesn't dog `for_each` as it does `accumulate`. Both algorithms can be used to summarize ranges. Use the one that suits you best.

You may be wondering why `for_each`'s function parameter is allowed to have side effects while `accumulate`'s is not. This is a probing question, one that strikes at the heart of the STL. Alas, gentle reader, there are some mysteries meant to remain beyond our ken. Why the difference between `accumulate` and `for_each`? I've yet to hear a convincing explanation.

6

Functors, Functor Classes, Functions, etc.

Like it or not, functions and function-like objects — *functors* — pervade the STL. Associative containers use them to keep their elements in order, some algorithms (e.g., `find_if`) use them to control their behavior, others (e.g., `for_each` and `transform`) are meaningless without them, and adapters like `not1` and `bind2nd` actively produce them.

Yes, everywhere you look in the STL, you see functors and functor classes. Including in your source code. It's not possible to make effective use of the STL without knowing how to write well-behaved functors. That being the case, most of this chapter is devoted to explaining how to make your functors behave the way the STL expects them to. One Item, however, is devoted to a different topic, one sure to appeal to those who've wondered about the need to litter their code with `ptr_fun`, `mem_fun`, and `mem_fun_ref`. Start with that Item ([Item 41](#)), if you like, but please don't stop there. Once you understand those functions, you'll need the information in the remaining Items to ensure that your functors work properly with them, as well as with the rest of the STL.

Item 38: Design functor classes for pass-by-value.

Neither C nor C++ allows you to truly pass functions as parameters to other functions. Instead, you must pass *pointers* to functions. For example, here's a declaration for the standard library function `qsort`:

```
void qsort( void *base, size_t nmemb, size_t size,  
           int (*cmpfcn)(const void*, const void*));
```

[Item 46](#) explains why the sort algorithm is typically a better choice than the `qsort` function, but that's not at issue here. What is at issue is the declaration for `qsort`'s parameter `cmpfcn`. Once you've squinted past all the asterisks, it becomes clear that the argument passed as `cmpfcn`, which is a pointer to a function, is copied (i.e., passed by

value) from the call site to `qsort`. This is representative of the rule followed by the standard libraries for both C and C++, namely, that function pointers are passed by value.

STL function objects are modeled after function pointers, so the convention in the STL is that function objects, too, are passed by value (i.e., copied) when passed to and from functions. This is perhaps best demonstrated by the Standard's declaration of `for_each`, an algorithm which both takes and returns a function object by value:

```
template<class InputIterator,
        class Function>
Function for_each(InputIterator first,           // note return-by-value
                 InputIterator last,
                 Function f);                  // note pass-by-value
```

In truth, the pass-by-value case is not quite this iron-clad, because `for_each`'s caller could explicitly specify the parameter types at the point of the call. For example, the following would cause `for_each` to pass and return its functor by reference:

```
class DoSomething:
public unary_function<int, void> {           // Item 40 explains base class
public:
    void operator()(int x) { ... }
    ...
};
typedef deque<int>::iterator DequeIntIter; // convenience typedef
deque<int> di;
...
DoSomething d;                             // create a function object
...
for_each<DequeIntIter,                      // call for_each with type
        DoSomething&>(di.begin(),          // parameters of DequeIntIter
                     di.end(),             // and DoSomething&;
                     d);                  // this forces d to be
                                         // passed and returned
                                         // by reference
```

Users of the STL almost never do this kind of thing, however, and some implementations of some STL algorithms won't even compile if function objects are passed by reference. For the remainder of this Item, I'm going to pretend that function objects are always passed by value. In practice, that's virtually always true.

Because function objects are passed and returned by value, the onus is on you to make sure that your function objects behave well when

passed that way (i.e., copied). This implies two things. First, your function objects need to be small. Otherwise they will be too expensive to copy. Second, your function objects must be *monomorphic* (i.e., not polymorphic) — they must not use virtual functions. That's because derived class objects passed by value into parameters of base class type suffer from the slicing problem: during the copy, their derived parts are removed. (For another example of how the slicing problem affects your use of the STL, see [Item 3](#).)

Efficiency is important, of course, and so is avoiding the slicing problem, but not all functors are small, and not all are monomorphic, either. One of the advantages of function objects over real functions is that functors can contain as much state as you need. Some function objects are naturally hefty, and it's important to be able to pass such functors to STL algorithms with the same ease as we pass their anorxic counterparts.

The prohibition on polymorphic functors is equally unrealistic. C++ supports inheritance hierarchies and dynamic binding, and these features are as useful when designing functor classes as anywhere else. Functor classes without inheritance would be like, well, like C++ without the “++”. Surely there's a way to let function objects be big and/or polymorphic, yet still allow them to mesh with the pass-functors-by-value convention that pervades the STL.

There is. Take the data and/or the polymorphism you'd like to put in your functor class, and move it into a different class. Then give your functor class a pointer to this new class. For example, if you'd like to create a polymorphic functor class containing lots of data,

```
template<typename T>
class BPFC:                                     // BPFC = " Big Polymorphic
public:                                         //      Functor Class"
    unary_function<T, void> {                 // Item 40 explains this
                                              // base class

private:
    Widget w;                                // this class has lots of data,
    int x;                                    // so it would be inefficient
    ...                                       // to pass it by value

public:
    virtual void operator()(const T& val) const; // this is a virtual function,
    ...                                       // so slicing would be bad
};
```

create a small, monomorphic class that contains a pointer to an implementation class, and put all the data and virtual functions in the implementation class:

```

template<typename T>                                // new implementation class
class BPFCImpl:                                     // for modified BPFC
{
public unary_function<T, void> {
private:
    Widget w;                                       // all the data that used to
    int x;                                         // be in BPFC are now here
    ...

    virtual ~BPFCImpl();                          // polymorphic classes need
                                                    // virtual destructors

    virtual void operator()(const T& val) const;

friend class BPFC<T>;                             // let BPFC access the data
};

template<typename T>
class BPFC:                                         // small, monomorphic
{
public unary_function<T, void> {                  // version of BPFC
private:
    BPFCImpl<T> *pImpl;                          // this is BPFC's only data

public:
    void operator()(const T& val) const           // this is now nonvirtual;
    {                                             // it forwards to BPFCImpl
        pImpl->operator()(val);
    }
    ...
};

```

The implementation of `BPFC::operator()` exemplifies how all BPFC would-be virtual functions are implemented: they call their truly virtual counterparts in `BPFCImpl`. The result is a functor class (BPFC) that's small and monomorphic, yet has access to a large amount of state and acts polymorphically.

I'm glossing over a fair number of details here, because the basic technique I've sketched is well known in C++ circles. *Effective C++* treats it in Item 34. In *Design Patterns* by Gamma *et al.* [6], this is called the “Bridge Pattern.” Sutter calls this the “Pimpl Idiom” in his *Exceptional C++* [8].

From an STL point of view, the primary thing to keep in mind is that functor classes using this technique must support copying in a reasonable fashion. If you were the author of BPFC above, you'd have to make sure that its copy constructor did something reasonable about the `BPFCImpl` object it points to. Perhaps the simplest reasonable thing would be to reference count it, using something like Boost's `shared_ptr`, which you can read about in [Item 50](#).

In fact, for purposes of this Item, the *only* thing you'd have to worry about would be the behavior of BPFC's copy constructor, because function objects are always copied — passed by value, remember? — when

passed to or returned from functions in the STL. That means two things. Make them small, and make them monomorphic.

Item 39: Make predicates pure functions.

I hate to do this to you, but we have to start with a short vocabulary lesson.

- A *predicate* is a function that returns `bool` (or something that can be implicitly converted to `bool`). Predicates are widely used in the STL. The comparison functions for the standard associative containers are predicates, and predicate functions are commonly passed as parameters to algorithms like `find_if` and the various sorting algorithms. (For an overview of the sorting algorithms, turn to [Item 31](#).)
- A *pure function* is a function whose return value depends only on its parameters. If `f` is a pure function and `x` and `y` are objects, the return value of `f(x, y)` can change only if the value of `x` or `y` changes.

In C++, all data consulted by pure functions are either passed in as parameters or are constant for the life of the function. (Naturally, such constant data should be declared `const`.) If a pure function consulted data that might change between calls, invoking the function at different times with the same parameters might yield different results, and that would be contrary to the definition of a pure function.

That should be enough to make it clear what it means to make predicates pure functions. All I have to do now is convince you that the advice is well founded. To help me do that, I hope you'll forgive me for burdening you with one more term.

- A *predicate class* is a functor class whose `operator()` function is a predicate, i.e., its `operator()` returns `true` or `false` (or something that can be implicitly converted to `true` or `false`). As you might expect, any place the STL expects a predicate, it will accept either a real predicate or an object of a predicate class.

That's it, I promise! Now we're ready to study why this Item offers guidance worth following.

[Item 38](#) explains that function objects are passed by value, so you should design function objects to be copied. For function objects that are predicates, there is another reason to design them to behave well when they are copied. Algorithms may make copies of functors and hold on to them a while before using them, and some algorithm imple-

mentations take advantage of this freedom. A critical repercussion of this observation is that *predicate functions must be pure functions*.

To appreciate why this is the case, let's suppose you were to violate this constraint. Consider the following (badly designed) predicate class. Regardless of the arguments that are passed, it returns true exactly once: the third time it is called. The rest of the time it returns false.

```
class BadPredicate:                                // see Item 40 for info
    public unary_function<Widget, bool> {          // on the base class
public:
    BadPredicate(): timesCalled(0) {}               // init timesCalled to 0
    bool operator()(const Widget&)
    {
        return ++timesCalled == 3;
    }
private:
    size_t timesCalled;
};
```

Suppose we use this class to eliminate the third Widget from a vector<Widget>:

```
vector<Widget> vw;                                // create vector and put some
...                                                // Widgets into it
vw.erase(remove_if(vw.begin(),                    // eliminate the third Widget;
                   vw.end(),                       // see Item 32 for info on how
                   BadPredicate()),                // erase and remove_if relate
          vw.end());
```

This code looks quite reasonable, but with many STL implementations, it will eliminate not just the third element from vw, it will also eliminate the sixth!

To understand how this can happen, it's helpful to see how `remove_if` is often implemented. Bear in mind that `remove_if` does not *have* to be implemented this way.

```
template <typename FwdIterator, typename Predicate>
FwdIterator remove_if(FwdIterator begin, FwdIterator end, Predicate p)
{
    begin = find_if(begin, end, p);
    if (begin == end) return begin;
    else {
        FwdIterator next = begin;
        return remove_copy_if(++next, end, begin, p);
    }
}
```

The details of this code are not important, but note that the predicate `p` is passed first to `find_if`, then later to `remove_copy_if`. In both cases, of course, `p` is passed by value — is *copied* — into those algorithms. (Technically, this need not be true, but practically, it is true. For details, see [Item 38](#).)

The initial call to `remove_if` (the one in the client code that wants to eliminate the third element from `vw`) creates an anonymous `BadPredicate` object, one with its internal `timesCalled` member set to zero. This object (known as `p` inside `remove_if`) is then copied into `find_if`, so `find_if` also receives a `BadPredicate` object with a `timesCalled` value of 0. `find_if` “calls” that object until it returns true, so it calls it three times. `find_if` then returns control to `remove_if`. `remove_if` continues to execute and ultimately calls `remove_copy_if`, passing as a predicate another copy of `p`. But `p`’s `timesCalled` member is still 0! `find_if` never called `p`, it called only a *copy* of `p`. As a result, the third time `remove_copy_if` calls its predicate, it, too, will return true. And that’s why `remove_if` will ultimately remove two `Widgets` from `vw` instead of just one.

The easiest way to keep yourself from tumbling into this linguistic crevasse is to declare your `operator()` functions `const` in predicate classes. If you do that, your compilers won’t let you change any class data members:

```
class BadPredicate:
    public unary_function<Widget, bool> {
public:
    bool operator()(const Widget&) const
    {
        return ++timesCalled == 3;           // error! can't change local data
    }                                         // in a const member function
    ...
};
```

Because this is such a straightforward way of preventing the problem we just examined, I very nearly entitled this Item “Make `operator()` `const` in predicate classes.” But that doesn’t go far enough. Even `const` member functions may access mutable data members, non-`const` local static objects, non-`const` class static objects, non-`const` objects at namespace scope, and non-`const` global objects. A well-designed predicate class ensures that its `operator()` functions are independent of those kinds of objects, too. Declaring `operator()` `const` in predicate classes is *necessary* for correct behavior, but it’s not *sufficient*. A well-behaved `operator()` is certainly `const`, but it’s more than that. It’s also a pure function.

Earlier in this Item, I remarked that any place the STL expects a predicate function, it will accept either a real function or an object of a predicate class. That's true in both directions. Any place the STL will accept an object of a predicate class, a predicate function (possibly modified by `ptr_fun` — see [Item 41](#)) is equally welcome. We now understand that `operator()` functions in predicate classes should be pure functions, so this restriction extends to predicate functions, too. This function is as bad a predicate as the objects generated from the `BadPredicate` class:

[illegible]

Regardless of how you write your predicates, they should always be pure functions.

Item 40: Make functor classes adaptable.

Suppose I have a list of Widget* pointers and a function to determine whether such a pointer identifies a Widget that is interesting:

```
list<Widget*> widgetPtrs;  
bool isInteresting(const Widget *pw);
```

If I'd like to find the first pointer to an interesting Widget in the list, it'd be easy:

```
list<Widget*>::iterator i = find_if(widgetPtrs.begin(), widgetPtrs.end(),
                                   isInteresting);

if (i != widgetPtrs.end()) {
    ...                                     // process the first
}                                           // ptr to an interesting
                                           // Widget
```

If I'd like to find the first pointer to a `Widget` that is *not* interesting, however, the obvious approach fails to compile:

```
list<Widget*>::iterator i =
    find_if(widgetPtrs.begin(), widgetPtrs.end(),
            not1(isInteresting));           // error! won't compile
```

Instead, I must apply `ptr_fun` to `isInteresting` before applying `not1`:

[illegible]

That leads to some questions. *Why* do I have to apply `ptr_fun` to `isInteresting` before applying `not1`? What does `ptr_fun` do for me, and how does it make the above work?

The answer is somewhat surprising. The only thing `ptr_fun` does is make some typedefs available. That's it. These typedefs are required by `not1`, and that's why applying `not1` to `ptr_fun` works, but applying `not1` to `isInteresting` directly doesn't work. Being a lowly function pointer, `isInteresting` lacks the typedefs that `not1` demands.

`not1` isn't the only component in the STL making such demands. Each of the four standard function adapters (`not1`, `not2`, `bind1st`, and `bind2nd`) requires the existence of certain typedefs, as do any non-standard STL-compatible adapters written by others (e.g., those available from SGI and Boost — see [Item 50](#)). Function objects that provide the necessary typedefs are said to be *adaptable*, while function objects lacking these typedefs are not adaptable. Adaptable function objects can be used in more contexts than can function objects that are not adaptable, so you should make your function objects adaptable whenever you can. It costs you nothing, and it may buy clients of your functor classes a world of convenience.

I know, I know, I'm being coy, constantly referring to “certain typedefs” without telling you what they are. The typedefs in question are `argument_type`, `first_argument_type`, `second_argument_type`, and `result_type`, but it's not quite that straightforward, because different kinds of functor classes are expected to provide different subsets of these names. In all honesty, unless you're writing your own adapters (a topic not covered in this book), you don't need to know anything about these typedefs. That's because the conventional way to provide them is to inherit them from a base class, or, more precisely, a base struct. For functor classes whose `operator()` takes one argument, the struct to inherit from is `std::unary_function`. For functor classes whose `operator()` takes two arguments, the struct to inherit from is `std::binary_function`.

Well, sort of. `unary_function` and `binary_function` are templates, so you can't inherit from them directly. Instead, you must inherit from structs they generate, and that requires that you specify some type arguments. For `unary_function`, you must specify the type of parameter taken by your functor class's `operator()`, as well as its return type. For `binary_function`, you specify three types: the types of your `operator()`'s first and second parameters, and your `operator()`'s return type.

Here are a couple of examples:

```
template<typename T>
class MeetsThreshold: public std::unary_function<Widget, bool> {
private:
    const T threshold;
public:
    MeetsThreshold(const T& threshold);
    bool operator()(const Widget&) const;
    ...
};

struct WidgetNameCompare:
    public std::binary_function<Widget, Widget, bool> {
    bool operator()(const Widget& lhs, const Widget& rhs) const;
};
```

In both cases, notice how the types passed to `unary_function` or `binary_function` are the same as the types taken and returned by the functor class's `operator()`, though it is a bit of an oddity that `operator()`'s return type is passed as the last argument to `unary_function` or `binary_function`.

You may have noticed that `MeetsThreshold` is a class, while `WidgetNameCompare` is a struct. `MeetsThreshold` has internal state (its `threshold` data member), and a class is the logical way to encapsulate such information. `WidgetNameCompare` has no state, hence no need to make anything private. Authors of functor classes where everything is public often declare structs instead of classes, probably for no other reason than to avoid typing “public” in front of the base class and the `operator()` function. Whether to declare such functors as classes or structs is purely a matter of personal style. If you're still refining your personal style and would like to emulate the pros, note that stateless functor classes within the STL itself (e.g., `less<T>`, `plus<T>`, etc.) are generally written as structs.

Look again at `WidgetNameCompare`:

```
struct WidgetNameCompare:
    public std::binary_function<Widget, Widget, bool> {
    bool operator()(const Widget& lhs, const Widget& rhs) const;
};
```

Even though `operator()`'s arguments are of type `const Widget&`, the type passed to `binary_function` is `Widget`. In general, non-pointer types passed to `unary_function` or `binary_function` have `const`s and references stripped off. (Don't ask why. The reasons are neither terribly good nor terribly interesting. If you're dying to know anyway, write some pro-

grams where you don't strip them off, then dissect the resulting compiler diagnostics. If, having done that, you're *still* interested in the matter, visit boost.org (see [Item 50](#)) and check out their work on call traits and function object adapters.)

The rules change when `operator()` takes pointer parameters. Here's a struct analogous to `WidgetNameCompare`, but this one works with `Widget*` pointers:

```
struct PtrWidgetNameCompare:
    public std::binary_function<const Widget*, const Widget*, bool> {
        bool operator()(const Widget* lhs, const Widget* rhs) const;
    };

```

Here, the types passed to `binary_function` are the *same* as the types taken by `operator()`. The general rule for functor classes taking or returning pointers is to pass to `unary_function` or `binary_function` whatever types `operator()` takes or returns.

Let's not forget the fundamental reason for all this `unary_function` and `binary_function` base class gobbledegook. These classes supply type-defs that are required by function object adapters, so inheritance from those classes yields adaptable function objects. That lets us do things like this:

```
list<Widget> widgets;
...
list<Widget>::reverse_iterator i1 =          // find the last widget
    find_if(widgets.rbegin(), widgets.rend(), // that fails to meet the
        not1(MeetsThreshold<int>(10)));      // threshold of 10
                                              // (whatever that means)

Widget w( constructor arguments );

list<Widget>::iterator i2 =                  // find the first widget
    find_if(widgets.begin(), widgets.end(),  // that precedes w in the
        bind2nd(WidgetNameCompare(), w));  // sort order defined by
                                              // WidgetNameCompare

```

Had we failed to have our functor classes inherit from `unary_function` or `binary_function`, neither of these examples would compile, because `not1` and `bind2nd` both work only with adaptable function objects.

STL function objects are modeled on C++ functions, and a C++ function has only one set of parameter types and one return type. As a result, the STL implicitly assumes that each functor class has only one `operator()` function, and it's the parameter and return types for this function that should be passed to `unary_function` or `binary_function` (in accord with the rules for reference and pointer types we just dis-

cussed). This means that, tempting though it might be, you shouldn't try to combine the functionality of `WidgetNameCompare` and `PtrWidgetNameCompare` by creating a single struct with two `operator()` functions. If you did, the functor would be adaptable with respect to at most one of its calling forms (whichever one you used when passing parameters to `binary_function`), and a functor that's adaptable only half the time might just as well not be adaptable at all.

Sometimes it makes sense to give a functor class multiple invocation forms (thereby abandoning adaptability), and Items 7, 20, 23, and 25 give examples of situations where that is the case. Such functor classes are the exception, however, not the rule. Adaptability is important, and you should strive to facilitate it each time you write a functor class.

Item 41: Understand the reasons for `ptr_fun`, `mem_fun`, and `mem_fun_ref`.

What is it with this `ptr_fun`/`mem_fun`/`mem_fun_ref` stuff? Sometimes you have to use these functions, sometimes you don't, and what do they do, anyway? They just seem to sit there, pointlessly hanging around function names like ill-fitting garments. They're unpleasant to type, annoying to read, and resistant to comprehension. Are these things additional examples of STL artifacts (such as the ones described in Items 10 and 18), or just some syntactic joke foisted on us by members of a Standardization Committee with too much free time and a twisted sense of humor?

Calm yourself. The names are less than inspired, but `ptr_fun`, `mem_fun`, and `mem_fun_ref` do important jobs, and as far as syntactic jokes go, one of the primary tasks of these functions is to paper over one of C++'s inherent syntactic inconsistencies.

If I have a function `f` and an object `x`, I wish to invoke `f` on `x`, and I'm outside `x`'s member functions, C++ gives me three different syntaxes for making the call:

<code>f(x);</code>	<code>// Syntax #1: When f is a // non-member function</code>
<code>x.f();</code>	<code>// Syntax #2: When f is a member // function and x is an object or // a reference to an object</code>
<code>p->f();</code>	<code>// Syntax #3: When f is a member // function and p is a pointer to x</code>

Now, suppose I have a function that can test Widgets,

```
void test(Widget& w);           // test w and mark it as "failed" if
                               // it doesn't pass the test
```

and I have a container of Widgets:

```
vector<Widget> vw;             // vw holds widgets
```

To test every Widget in vw, I can use `for_each` in the obvious manner:

```
for_each(vw.begin(), vw.end(), test); // Call #1 (compiles)
```

But imagine that `test` is a member function of `Widget` instead of a non-member function, i.e., that `Widget` supports self-testing:

```
class Widget {
public:
    ...
    void test();                // perform a self-test; mark *this
    ...                          // as "failed" if it doesn't pass
};
```

In a perfect world, I'd also be able to use `for_each` to invoke `Widget::test` on each object in `vw`:

```
for_each(vw.begin(), vw.end(),
         &Widget::test);      // Call #2 (won't compile)
```

In fact, if the world were really perfect, I'd be able to use `for_each` to invoke `Widget::test` on a container of `Widget*` pointers, too:

```
list<Widget*> lpw;             // lpw holds pointers to widgets
for_each( lpw.begin(), lpw.end(),
         &Widget::test);      // Call #3 (also won't compile)
```

But think of what would have to happen in this perfect world. Inside the `for_each` function in **Call #1**, we'd be calling a non-member function with an object, so we'd have to use **Syntax #1**. Inside the `for_each` function in **Call #2**, we'd have to use **Syntax #2**, because we'd have an object and a member function. And inside the `for_each` function in **Call #3**, we'd need to use **Syntax #3**, because we'd be dealing with a member function and a pointer to an object. We'd therefore need three different versions of `for_each`, and how perfect would that world be?

In the world we do have, we possess only one version of `for_each`. It's not hard to envision an implementation:

```
template<typename InputIterator, typename Function>
Function for_each(InputIterator begin, InputIterator end, Function f)
{
    while (begin != end) f(*begin++);
}
```

Here I've highlighted the fact that `for_each` uses Syntax #1 when making the call. This is a universal convention in the STL: functions and function objects are always invoked using the syntactic form for non-member functions. This explains why Call #1 compiles while Calls #2 and #3 don't. It's because STL algorithms (including `for_each`) hard-wire in Syntax #1, and only Call #1 is compatible with that syntax.

Perhaps it's now clear why `mem_fun` and `mem_fun_ref` exist. They arrange for member functions (which must ordinarily be called using Syntax #2 or #3) to be called using Syntax #1.

The way `mem_fun` and `mem_fun_ref` do this is simple, though it's a little clearer if you take a look at a declaration for one of these functions. They're really function templates, and several variants of the `mem_fun` and `mem_fun_ref` templates exist, corresponding to different numbers of parameters and the constness (or lack thereof) of the member functions they adapt. Seeing one declaration is enough to understand how things are put together:

```
template<typename R, typename C>    // declaration for mem_fun for
mem_fun_t<R,C>                    // non-const member functions
mem_fun(R (C::*pmf)());            // taking no parameters. C is the
                                   // class, R is the return type of the
                                   // pointed-to member function
```

`mem_fun` takes a pointer to a member function, `pmf`, and returns an object of type `mem_fun_t`. This is a functor class that holds the member function pointer and offers an `operator()` that invokes the pointed-to member function on the object passed to `operator()`. For example, in this code,

```
list<Widget*> lpw;                  // same as above
...
for_each(lpw.begin(), lpw.end(),
         mem_fun(&Widget::test));  // this will now compile
```

`for_each` receives an object of type `mem_fun_t` holding a pointer to `Widget::test`. For each `Widget*` pointer in `lpw`, `for_each` "calls" the `mem_fun_t` object using Syntax #1, and that object immediately invokes `Widget::test` on the `Widget*` pointer using Syntax #3.

Overall, `mem_fun` adapts Syntax #3, which is what `Widget::test` requires when used with a `Widget*` pointer, to Syntax #1, which is what `for_each` uses. It's thus no wonder that classes like `mem_fun_t` are known as *function object adapters*. It should not surprise you to learn that, completely analogously with the above, the `mem_fun_ref` functions adapt Syntax #2 to Syntax #1 and generate adapter objects of type `mem_fun_ref_t`.

The objects produced by `mem_fun` and `mem_fun_ref` do more than allow STL components to assume that all functions are called using a single syntax. They also provide important typedefs, just like the objects produced by `ptr_fun`. The story behind these typedefs is told in [Item 40](#), so I won't repeat it here. However, this puts us in a position to understand why this call compiles,

```
for_each(vw.begin(), vw.end(), test);           // as above, Call #1;
                                                // this compiles
```

while these do not:

```
for_each(vw.begin(), vw.end(), &Widget::test); // as above, Call #2;
                                                // doesn't compile

for_each(lpw.begin(), lpw.end(), &Widget::test); // as above, Call #3;
                                                // doesn't compile
```

The first call (Call #1) passes a real function, so there's no need to adapt its calling syntax for use by `for_each`; the algorithm will inherently call it using the proper syntax. Furthermore, `for_each` makes no use of the typedefs that `ptr_fun` adds, so it's not necessary to use `ptr_fun` when passing `test` to `for_each`. On the other hand, adding the typedefs can't hurt anything, so this will do the same thing as the call above:

```
for_each(vw.begin(), vw.end(), ptr_fun(test)); // compiles and behaves
                                                // like Call #1 above
```

If you get confused about when to use `ptr_fun` and when not to, consider using it every time you pass a function to an STL component. The STL won't care, and there is no runtime penalty. About the worst that can be said is that some people reading your code might raise an eyebrow when they see an unnecessary use of `ptr_fun`. How much that bothers you depends, I suppose, on your sensitivity to raised eyebrows.

An alternative strategy with respect to `ptr_fun` is to use it only when you're forced to. If you omit it when the typedefs are necessary, your compilers will balk at your code. Then you'll have to go back and add it.

The situation with `mem_fun` and `mem_fun_ref` is fundamentally different. You must employ them whenever you pass a member function to an STL component, because, in addition to adding typedefs (which may or may not be necessary), they adapt the calling syntaxes from the ones normally used with member functions to the one used everywhere in the STL. If you don't use them when passing member function pointers, your code will never compile.

Which leaves just the names of the member function adapters, and here, finally, we have a genuine historical STL artifact. When the need for these kinds of adapters first became apparent, the people working on the STL focused on containers of pointers. (Given the drawbacks of such containers described in Items 7, 20, and 33, this might seem surprising, but remember that containers of pointers support polymorphism, while containers of objects do not.) They needed an adapter for member functions, so they chose `mem_fun`. Only later did they realize that they needed a different adapter for containers of objects, so they hacked up the name `mem_fun_ref` for that. No, it's not very elegant, but these things happen. Tell me *you've* never given any of your components a name that you later realized was, er, difficult to generalize... .

Item 42: Make sure `less<T>` means `operator<`.

As all Widget-savvy people are aware, Widgets have both a weight and a maximum speed:

```
class Widget {
public:
    ...
    size_t weight() const;
    size_t maxSpeed() const;
    ...
};
```

Also well known is that the natural way to sort Widgets is by weight. `operator<` for Widgets reflects this:

```
bool operator<(const Widget& lhs, const Widget& rhs)
{
    return lhs.weight() < rhs.weight();
}
```

But suppose we'd like to create a `multiset<Widget>` where the Widgets are sorted by maximum speed. We know that the default comparison function for `multiset<Widget>` is `less<Widget>`, and we know that `less<Widget>`, by default, does its work by calling `operator<` for Widgets. That being the case, it seems clear that one way to get a `multiset<Widget>` sorted by maximum speed is to sever the tie between `less<Widget>` and `operator<` by specializing `less<Widget>` to look only at a Widget's maximum speed:

```
template<>                                     // This is a specialization
struct std::less<Widget>:                     // of std::less for Widget;
    public                                     // it's also a very bad idea
    std::binary_function<Widget,
                        Widget,
                        bool> {                // See Item 40 for info
                                                // on this base class

    bool operator()(const Widget& lhs, const Widget& rhs) const
    {
        return lhs.maxSpeed() < rhs.maxSpeed();
    }
};
```

This both looks ill-advised and is ill-advised, but it may not be ill-advised for the reason you think. Does it surprise you that it compiles at all? Many programmers point out that the above isn't just a specialization of a template, it's a specialization of a template in the `std` namespace. "Isn't `std` supposed to be sacred, reserved for library implementers and beyond the reach of mere programmers?" they ask. "Shouldn't compilers reject this attempt to tamper with the workings of the C++ immortals?" they wonder.

As a general rule, trying to modify components in `std` is indeed forbidden (and doing so typically transports one to the realm of undefined behavior), but under some circumstances, tinkering is allowed. Specifically, programmers are allowed to specialize templates in `std` for user-defined types. Almost always, there are alternatives that are superior to specializing `std` templates, but on rare occasions, it's a reasonable thing to do. For instance, authors of smart pointer classes often want their classes to act like built-in pointers for sorting purposes, so it's not uncommon to see specializations of `std::less` for smart pointer types. The following, for example, is part of the Boost library's `shared_ptr`, a smart pointer you can read about in [Items 7](#) and [50](#):

```
namespace std {  
    template<typename T>  
        struct less< boost::shared_ptr<T> >:  
            public  
                binary_function<boost::shared_ptr<T>,  
                                boost::shared_ptr<T>, // this is the customary  
                                bool> {               // base class (see Item 40)  
  
                bool operator()( const boost::shared_ptr<T>& a,  
                                const boost::shared_ptr<T>& b) const  
                {  
                    return less<T*>()(a.get(),b.get()); // shared_ptr::get returns  
                                                         // the built-in pointer that's  
                                                         // in the shared_ptr object  
                }  
};  
}
```

This isn't unreasonable, and it certainly serves up no surprises, because this specialization of `less` merely ensures that sorting smart pointers behaves the same as sorting their built-in brethren. Alas, our tentative specialization of `less` for `Widget` *does* serve up surprises.

C++ programmers can be forgiven certain assumptions. They assume that copy constructors copy, for example. (As [Item 8](#) attests, failure to adhere to this convention can lead to astonishing behavior.) They assume that taking the address of an object yields a pointer to that object. (Turn to [Item 18](#) to read about what can happen when this isn't true.) They assume that adapters like `bind1st` and `not2` may be applied to function objects. ([Item 40](#) explains how things break when this isn't the case.) They assume that `operator+` adds (except for strings, but there's a long history of using "+" to mean string concatenation), that `operator-` subtracts, that `operator==` compares. And they assume that using `less` is tantamount to using `operator<`.

`operator<` is more than just the default way to implement `less`, it's *what programmers expect less to do*. Having `less` do something other than call `operator<` is a gratuitous violation of programmers' expectations. It runs contrary to what has been called "the principle of least astonishment." It's callous. It's mean. It's bad. You shouldn't do it.

Especially when there's no reason to. There's not a place in the STL using `less` where you can't specify a different comparison type instead. Returning to our original example of a `multiset<Widget>` ordered by maximum speed, all we need to do to get what we want is create a functor class called almost anything *except* `less` that performs the comparison we're interested in. Why, here's one now:

```
struct MaxSpeedCompare:
    public binary_function<Widget, Widget, bool> {
        bool operator()(const Widget& lhs, const Widget& rhs) const
        {
            return lhs.maxSpeed() < rhs.maxSpeed();
        }
    };
```

To create our `multiset`, we use `MaxSpeedCompare` as the comparison type, thus avoiding use of the default comparison type (which is, of course, `less<Widget>`):

```
multiset<Widget, MaxSpeedCompare> widgets;
```

This code says exactly what it means. It creates a `multiset` of `Widgets` sorted as defined by the functor class `MaxSpeedCompare`.

Contrast that with this:

```
multiset<Widget> widgets;
```

This says that `widgets` is a multiset of `Widgets` sorted in the default manner. Technically, that means it uses `less<Widget>`, but virtually everybody is going to assume that really means it's sorted by `operator<`.

Don't mislead all those programmers by playing games with the definition of `less`. If you use `less` (explicitly or implicitly), make sure it means `operator<`. If you want to sort objects using some other criterion, create a special functor class that's *not* called `less`. It's really as simple as that.

7

Programming with the STL

It's traditional to summarize the STL as consisting of containers, iterators, algorithms, and function objects, but *programming* with the STL is much more than that. Programming with the STL is knowing when to use loops, when to use algorithms, and when to use container member functions. It's knowing when `equal_range` is a better way to search than `lower_bound`, knowing when `lower_bound` is preferable to `find`, and knowing when `find` beats `equal_range`. It's knowing how to improve algorithm performance by substituting functors for functions that do the same thing. It's knowing how to avoid unportable or incomprehensible code. It's even knowing how to read compiler error messages that run to thousands of characters. And it's knowing about Internet resources for STL documentation, STL extensions, even complete STL implementations.

Yes, programming with the STL involves knowing many things. This chapter gives you much of the knowledge you need.

Item 43: Prefer algorithm calls to hand-written loops.

Every algorithm takes at least one pair of iterators that specify a range of objects over which to do something. `min_element` finds the smallest value in the range, for example, while `accumulate` summarizes some information about the range as a whole (see [Item 37](#)) and `partition` separates all the elements of a range into those that do and do not satisfy some criterion (see [Item 31](#)). For algorithms to do their work, they must examine every object in the range(s) they are passed, and they do this in the way you'd expect: they loop from the beginning of the range(s) to the end. Some algorithms, such as `find` and `find_if`, may return before they complete the traversal, but even these algorithms internally contain a loop. After all, even `find` and `find_if` must look at every element of a range before they can conclude that what they are looking for is *not* present.

Internally, then, algorithms are loops. Furthermore, the breadth of STL algorithms means that many tasks you might naturally code as loops could also be written using algorithms. For example, if you have a `Widget` class that supports redrawing,

```
class Widget {
public:
    ...
    void redraw() const;
    ...
};
```

and you'd like to redraw all the `Widgets` in a list, you could do it with a loop, like this,

```
list<Widget> lw;
...
for (list<Widget>::iterator i = lw.begin(); i != lw.end(); ++i) {
    i->redraw();
}
```

but you could also do it with the `for_each` algorithm:

```
for_each(lw.begin(), lw.end(),                // see Item 41 for info
         mem_fun_ref(&Widget::redraw));       // on mem_fun_ref
```

For many C++ programmers, writing the loop is more natural than calling the algorithm, and reading the loop is more comfortable than making sense of `mem_fun_ref` and the taking of `Widget::redraw`'s address. Yet this Item argues that the algorithm call is preferable. In fact, this Item argues that calling an algorithm is usually preferable to *any* hand-written loop. Why?

There are three reasons:

- **Efficiency:** Algorithms are often more efficient than the loops programmers produce.
- **Correctness:** Writing loops is more subject to errors than is calling algorithms.
- **Maintainability:** Algorithm calls often yield code that is clearer and more straightforward than the corresponding explicit loops.

The remainder of this Item lays out the case for algorithms.

From an efficiency perspective, algorithms can beat explicit loops in three ways, two major, one minor. The minor way involves the elimination of redundant computations. Look again at the loop we just saw:

```
for (list<Widget>::iterator i = lw.begin(); i != lw.end(); ++i) {
    i->redraw();
}
```

I've highlighted the loop termination test to emphasize that each time around the loop, `i` will be checked against `lw.end()`. That means that each time around the loop, the function `list::end` will be invoked. But we don't need to call `end` more than once, because we're not modifying the list. A single call to `end` would suffice, and, if we look again at the algorithm invocation, we'll see that that's exactly how many times `end` is evaluated:

```
for_each(lw.begin(), lw.end(),                // this call evaluates
         mem_fun_ref(&Widget::redraw));       // lw.end() exactly once
```

To be fair, STL implementers understand that `begin` and `end` (and similar functions, such as `size`) are used frequently, so they're likely to design them for maximal efficiency. They'll almost certainly inline them and strive to code them so that most compilers will be able to avoid repeated computations by hoisting their results out of loops like the one above. Experience shows that implementers don't always succeed, however, and when they don't, the avoidance of repeated computations is enough to give the algorithm a performance edge over the hand-written loop.

But that's the minor efficiency argument. The first major argument is that library implementers can take advantage of their knowledge of container implementations to optimize traversals in a way that no library user ever could. For example, the objects in a deque are typically stored (internally) in one or more fixed-size arrays. Pointer-based traversals of these arrays are faster than iterator-based traversals, but only library implementers can use pointer-based traversals, because only they know the size of the internal arrays and how to move from one array to the next. Some STLs contain algorithm implementations that take their deque's internal data structures into account, and such implementations have been known to clock in at more than 20% faster than the "normal" implementations of the algorithms.

The point is not that STL implementations are optimized for deques (or any other specific container type), but that implementers know more about their implementations than you do, and they can take advantage of this knowledge in algorithm implementations. If you shun algorithm calls in favor of your own loops, you forgo the opportunity to benefit from any implementation-specific optimizations they may have provided.

The second major efficiency argument is that all but the most trivial STL algorithms use computer science algorithms that are more sophisticated — sometimes *much* more sophisticated — than anything the average C++ programmer will be able to come up with. It's next to

This looks like a double win, because it not only increments the iterator specifying the insertion position, it also eliminates the need to call `begin` each time around the loop; that eliminates the minor efficiency hit we discussed earlier. Alas, this approach runs into a different problem: it yields undefined results. Each time `deque::insert` is called, it invalidates all iterators into the deque, and that includes `insertLocation`. After the first call to `insert`, `insertLocation` is invalidated, and subsequent loop iterations are allowed to head straight to looneyland.

Once you puzzle this out (possibly with the aid of STLport's debug mode, which is described in [Item 50](#)), you might come up with the following:

```
deque<double>::iterator insertLocation =
    d.begin();                                // as before
for (size_t i = 0; i < numDoubles; ++i) {      // update insertLocation
    insertLocation =                          // each time insert is
    d.insert(insertLocation, data[i] + 41);    // called to keep the
    ++insertLocation;                        // iterator valid, then
}                                              // increment it
```

This code finally does what you want, but think about how much work it took to get here! Compare that to the following call to `transform`:

```
transform(data, data + numDoubles,            // copy all elements
    inserter(d, d.begin()),                   // from data to the front
    bind2nd(plus<double>(), 41));            // of d, adding 41 to each
```

The “`bind2nd(plus<double>(), 41)`” might take you a couple of minutes to get right (especially if you don't use STL's binders very often), but the only iterator-related worries you have are specifying the beginning and end of the source range (which was never a problem) and being sure to use `inserter` as the beginning of the destination range (see [Item 30](#)). In practice, figuring out the correct initial iterators for source and destination ranges is usually easy, or at least a lot easier than making sure the body of a loop doesn't inadvertently invalidate an iterator you need to keep using.

This example is representative of a broad class of loops that are difficult to write correctly, because you have to be on constant alert for iterators that are incorrectly manipulated or are invalidated before you're done using them. To see a different example of how inadvertent iterator invalidation can lead to trouble, turn to [Item 9](#), which describes the subtleties involved in writing loops that call `erase`.

Given that using invalidated iterators leads to undefined behavior, and given that undefined behavior has a nasty habit of failing to show

itself during development and testing, why run the risk if you don't have to? Turn the iterators over to the algorithms, and let *them* worry about the vagaries of iterator manipulation.

I've explained why algorithms can be more efficient than hand-written loops, and I've described why such loops must navigate a thicket of iterator-related difficulties that algorithms avoid. With luck, you are now an algorithm believer. Yet luck is fickle, and I'd prefer a more secure conviction before I rest my case. Let us therefore move on to the issue of code clarity. In the long run, the best software is the clearest software, the software that is easiest to understand, the software that can most readily be enhanced, maintained, and molded to fit new circumstances. The familiarity of loops notwithstanding, algorithms have an advantage in this long-term competition.

The key to their edge is the power of a known vocabulary. There are 70 algorithm names in the STL — a total of over 100 different function templates, once overloading is taken into account. Each of those algorithms carries out some well-defined task, and *it is reasonable to expect professional C++ programmers to know (or be able to look up) what each does*. Thus, when a programmer sees a transform call, that programmer recognizes that some function is being applied to every object in a range, and the results of those calls are being written somewhere. When the programmer sees a call to replace_if, he or she knows that all the objects in a range that satisfy some predicate are being modified. When the programmer comes across an invocation of partition, she or he understands that the objects in a range are being moved around so that all the objects satisfying a predicate are grouped together (see [Item 31](#)). The names of STL algorithms convey a lot of semantic information, and that makes them clearer than any random loop can hope to be.

When you see a for, while, or do, all you know is that some kind of loop is coming up. To acquire even the faintest idea of what that loop does, you have to examine it. Not so with algorithms. Once you see a call to an algorithm, the name alone sketches the outline of what it does. To understand exactly what will happen, of course, you must inspect the arguments being passed to the algorithm, but that's often less work than trying to divine the intent of a general looping construct.

Simply put, algorithm names suggest what they do. “for,” “while,” and “do” don't. In fact, this is true of any component of the standard C or C++ library. Without doubt, you could write your own implementations of strlen, memset, or bsearch, if you wanted to, but you don't. Why

not? Because (1) somebody has already written them, so there's no point in your doing it again; (2) the names are standard, so everybody knows what they do; and (3) you suspect that your library implementer knows some efficiency tricks you don't know, and you're unwilling to give up the possible optimizations a skilled library implementer might provide. Just as you don't write your own versions of *strlen* *et al.*, it makes no sense to write loops that duplicate functionality already present in STL algorithms.

I wish that were the end of the story, because I think it's a strong finish. Alas, this is a tale that refuses to go gentle into that good night.

Algorithm names are more meaningful than bare loops, it's true, but specifying what to do during an iteration can be clearer using a loop than using an algorithm. For example, suppose you'd like to identify the first element in a vector whose value is greater than some *x* and less than some *y*. Here's how you could do it using a loop:

```
vector<int> v;
int x, y;
...
vector<int>::iterator i = v.begin();    // iterate from v.begin() until an
for( ; i != v.end(); ++i) {             // appropriate value is found or
    if (*i > x && *i < y) break;          // v.end() is reached
}
...                                     // i now points to the value or is
                                       // the same as v.end()
```

It is possible to pass this same logic to `find_if`, but it requires that you use a nonstandard function object adapter like SGI's `compose2` (see [Item 50](#)):

```
vector<int>::iterator i =
    find_if(v.begin(), v.end(),           // find the first value val/
        compose2(logical_and<bool>(),    // where the "and" of
            bind2nd(greater<int>(), x),    // val > x and
            bind2nd(less<int>(), y)));      // val < y
                                           // is true
```

Even if this didn't use nonstandard components, many programmers would object that it's nowhere near as clear as the loop, and I have to admit to being sympathetic to that view (see [Item 47](#)).

The `find_if` call can be made less imposing by moving the test logic into a separate functor class,

```

template<typename T>
class BetweenValues:
public unary_function<T, bool> {           // see Item 40
public:
    BetweenValues(const T& lowValue,
                  const T& highValue)      // have the ctor save the
: lowVal(lowValue), highVal(highValue)    // values to be between
    {}

    bool operator()(const T& val) const    // return whether
    {                                     // val is between the
        return val > lowVal && val < highVal; // saved values
    }

private:
    T lowVal;
    T highVal;
};

...
vector<int>::iterator i = find_if(v.begin(), v.end(),
                                BetweenValues<int>(x, y));

```

but this has its own drawbacks. First, creating the `BetweenValues` template is a lot more work than writing the loop body. Just count the lines. Loop body: one; `BetweenValues` template: fourteen. Not a very good ratio. Second, the details of what `find_if` is looking for are now physically separate from the call. To really understand the call to `find_if`, one must look up the definition of `BetweenValues`, but `BetweenValues` must be defined outside the function containing the call to `find_if`. If you try to declare `BetweenValues` *inside* the function containing the call to `find_if`, like this,

```

{                                           // beginning of function
...
    template <typename T>
    class BetweenValues: public unary_function<T, bool> { ... };
    vector<int>::iterator i = find_if(v.begin(), v.end(),
                                      BetweenValues<int>(x, y));

} ...                                     // end of function

```

you'll discover that it won't compile, because templates can't be declared inside functions. If you try to avoid that restriction by making `BetweenValues` a class instead of a template,

```
{ // beginning of function
...
class BetweenValues: public unary_function<int, bool> { ... };
vector<int>::iterator i = find_if(v.begin(), v.end(),
                                BetweenValues(x, y));
} // end of function
```

you'll find that you're still out of luck, because classes defined inside functions are known as *local classes*, and local class types can't be bound to template type arguments (such as the type of the functor taken by `find_if`). Sad as it may seem, functor classes and functor class templates are not allowed to be defined inside functions, no matter how convenient it would be to be able to do it.

In the ongoing tussle between algorithm calls and hand-written loops, the bottom line on code clarity is that it all depends on what you need to do inside the loop. If you need to do something an algorithm already does, or if you need to do something very similar to what an algorithm does, the algorithm call is clearer. If you need a loop that does something fairly simple, but would require a confusing tangle of binders and adapters or would require a separate functor class if you were to use an algorithm, you're probably better off just writing the loop. Finally, if you need to do something fairly long and complex inside the loop, the scales tilt back toward algorithms, because long, complex computations should generally be moved into separate functions, anyway. Once you've moved the loop body into a separate function, you can almost certainly find a way to pass that function to an algorithm (often `for_each`) such that the resulting code is direct and straightforward.

If you agree with this Item that algorithm calls are generally preferable to hand-written loops, and if you also agree with [Item 5](#) that range member functions are preferable to loops that iteratively invoke single-element member functions, an interesting conclusion emerges: well-crafted C++ programs using the STL contain far fewer loops than equivalent programs not using the STL. This is a good thing. Any time we can replace low-level words like `for`, `while`, and `do` with higher-level terms like `insert`, `find`, and `for_each`, we raise the level of abstraction in our software and thereby make it easier to write, document, enhance, and maintain.

Item 44: Prefer member functions to algorithms with the same names.

Some containers have member functions with the same names as STL algorithms. The associative containers offer `count`, `find`, `lower_bound`, `upper_bound`, and `equal_range`, while `list` offers `remove`, `remove_if`, `unique`, `sort`, `merge`, and `reverse`. Most of the time, you'll want to use the member functions instead of the algorithms. There are two reasons for this. First, the member functions are faster. Second, they integrate better with the containers (especially the associative containers) than do the algorithms. That's because algorithms and member functions that share the same name typically do *not* do exactly the same thing.

We'll begin with an examination of the associative containers. Suppose you have a `set<int>` holding a million values and you'd like to find the first occurrence of the value 727, if there is one. Here are the two most obvious ways to perform the search:

```

set<int> s;                                // create set, put
...                                       // 1,000,000 values
                                       // into it

set<int>::iterator i = s.find(727);        // use find member
if (i != s.end()) ...                     // function

set<int>::iterator i = find(s.begin(), s.end(), 727); // use find algorithm
if (i != s.end()) ...
```

The `find` member function runs in logarithmic time, so, regardless of whether 727 is in the set, `set::find` will perform no more than about 40 comparisons looking for it, and usually it will require only about 20. In contrast, the `find` algorithm runs in linear time, so it will have to perform 1,000,000 comparisons if 727 isn't in the set. Even if 727 is in the set, the `find` algorithm will perform, on average, 500,000 comparisons to locate it. The efficiency score is thus

```

Member find:    About 40 (worst case) to about 20 (average case)
Algorithm find: 1,000,000 (worst case) to 500,000 (average case)
```

As in golf, the low score wins, and as you can see, this matchup is not much of a contest.

I have to be a little cagey about the number of comparisons required by member `find`, because it's partially dependent on the implementation used by the associative containers. Most implementations use red-black trees, a form of balanced tree that may be out of balance by up to a factor of two. In such implementations, the maximum number of comparisons needed to search a set of a million values is 38, but for the vast majority of searches, no more than 22 comparisons are

required. An implementation based on perfectly balanced trees would never require more than 21 comparisons, but in practice, the overall performance of such perfectly balanced trees is inferior to that of red-black trees. That's why most STL implementations use red-black trees.

Efficiency isn't the only difference between member and algorithm `find`. As [Item 19](#) explains, STL algorithms determine whether two objects have the same value by checking for equality, while associative containers use equivalence as their “sameness” test. Hence, the `find` algorithm searches for 727 using equality, while the `find` member function searches using equivalence. The difference between equality and equivalence can be the difference between a successful search and an unsuccessful search. For example, [Item 19](#) shows how using the `find` algorithm to look for something in an associative container could fail even when the corresponding search using the `find` member function would succeed! You should therefore prefer the member form of `find`, `count`, `lower_bound`, etc., over their algorithm eponyms when you work with associative containers, because they offer behavior that is consistent with the other member functions of those containers. Due to the difference between equality and equivalence, algorithms don't offer such consistent behavior.

This difference is especially pronounced when working with maps and multimaps, because these containers hold pair objects, yet their member functions look only at the key part of each pair. Hence, the `count` member function counts only pairs with matching keys (a “match,” naturally, is determined by testing for equivalence); the value part of each pair is ignored. The member functions `find`, `lower_bound`, `upper_bound`, and `equal_range` behave similarly. If you use the `count` algorithm, however, it will look for matches based on (a) equality and (b) both components of the pair; `find`, `lower_bound`, etc., do the same thing. To get the algorithms to look at only the key part of a pair, you have to jump through the hoops described in [Item 23](#) (which would also allow you to replace equality testing with equivalence testing).

On the other hand, if you are really concerned with efficiency, you may decide that [Item 23](#)'s gymnastics, in conjunction with the logarithmic-time lookup algorithms of [Item 34](#), are a small price to pay for an increase in performance. Then again, if you're *really* concerned with efficiency, you'll want to consider the non-standard hashed containers described in [Item 25](#), though there you'll again confront the difference between equality and equivalence.

For the standard associative containers, then, choosing member functions over algorithms with the same names offers several benefits.

First, you get logarithmic-time instead of linear-time performance. Second, you determine whether two values are “the same” using equivalence, which is the natural definition for associative containers. Third, when working with maps and multimaps, you automatically deal only with key values instead of with (key, value) pairs. This triumvirate makes the case for preferring member functions pretty iron-clad.

Let us therefore move on to list member functions that have the same names as STL algorithms. Here the story is almost completely about efficiency. Each of the algorithms that list specializes (remove, remove_if, unique, sort, merge, and reverse) copies objects, but list-specific versions copy nothing; they simply manipulate the pointers connecting list nodes. The algorithmic complexity of the algorithms and the member functions is the same, but, under the assumption that manipulating pointers is less expensive than copying objects, list’s versions of these functions should offer better performance.

It’s important to bear in mind that the list member functions often behave differently from their algorithm counterparts. As [Item 32](#) explains, calls to the algorithms remove, remove_if, and unique must be followed by calls to erase if you really want to eliminate objects from a container, but list’s remove, remove_if, and unique member functions honestly get rid of elements; no subsequent call to erase is necessary.

A significant difference between the sort algorithm and list’s sort function is that the former can’t be applied to lists. Being only bidirectional iterators, list’s iterators can’t be passed to sort. A gulf also exists between the behavior of the merge algorithm and list’s merge. The algorithm isn’t permitted to modify its source ranges, but list::merge *always* modifies the lists it works on.

So there you have it. When faced with a choice between an STL algorithm or a container member function with the same name, you should prefer the member function. It’s almost certain to be more efficient, and it’s likely to be better integrated with the container’s usual behavior, too.

Item 45: Distinguish among count, find, binary_search, lower_bound, upper_bound, and equal_range.

So you want to look for something, and you have a container or you have iterators demarcating a range where you think it’s located. How do you conduct the search? Your quiver is fairly bursting with arrows: count, count_if, find, find_if, binary_search, lower_bound, upper_bound, and equal_range. Decisions, decisions! How do you choose?

Easy. You reach for something that's fast and simple. The faster and simpler, the better.

For the time being, we'll assume that you have a pair of iterators specifying a range to be searched. Later, we'll consider the case where you have a container instead of a range.

In selecting a search strategy, much depends on whether your iterators define a sorted range. If they do, you can get speedy (usually logarithmic-time — see [Item 34](#)) lookups via `binary_search`, `lower_bound`, `upper_bound`, and `equal_range`. If the iterators don't demarcate a sorted range, you're limited to the linear-time algorithms `count`, `count_if`, `find`, and `find_if`. In what follows, I'll ignore the `_if` variants of `count` and `find`, just as I'll ignore the variants of `binary_search`, `lower_` and `upper_bound`, and `equal_range` taking a predicate. Whether you rely on the default search predicate or you specify your own, the considerations for choosing a search algorithm are the same.

If you have an unsorted range, your choices are `count` or `find`. They answer slightly different questions, so it's worth taking a closer look at them. `count` answers the question, "Is the value there, and if so, how many copies are there?" while `find` answers the question, "Is it there, and if so, where is it?"

Suppose all you want to know is whether some special `Widget` value `w` is in a list. Using `count`, the code looks like this:

```
list<Widget> lw;           // list of Widgets
Widget w;                 // special Widget value

...
if (count(lw.begin(), lw.end(), w)) {
    ...                   // w is in lw
} else {
    ...                   // it's not
}
```

This demonstrates a common idiom: using `count` as an existence test. `count` returns either zero or a positive number, so we rely on the conversion of nonzero values to true and zero to false. It would arguably be clearer to be more explicit about what we are doing,

```
if (count(lw.begin(), lw.end(), w) != 0) ...
```

and some programmers do write it that way, but it's quite common to rely on the implicit conversion, as in the original example.

Compared to that original code, using `find` is slightly more complicated, because you have to test `find`'s return value against the list's end iterator:

```

if (find(lw.begin(), lw.end(), w) != lw.end()) {
    ...
} else {
    ...
}

```

For existence testing, the idiomatic use of `count` is slightly simpler to code. At the same time, it's also less efficient when the search is successful, because `find` stops once it's found a match, while `count` must continue to the end of the range looking for additional matches. For most programmers, `find`'s edge in efficiency is enough to justify the slight increase in usage complexity.

Often, knowing whether a value is in a range isn't enough. Instead, you'll want to know the first object in the range with the value. For example, you might want to print the object, you might want to insert something in front of it, or you might want to erase it (but see [Item 9](#) for guidance on erasing while iterating). When you need to know not just whether a value exists but also which object (or objects) has that value, you need `find`:

```

list<Widget>::iterator i = find(lw.begin(), lw.end(), w);
if (i != lw.end()) {
    ...                               // found it, i points to the first one
} else {
    ...                               // didn't find it
}

```

For sorted ranges, you have other choices, and you'll definitely want to use them. `count` and `find` run in linear time, but the search algorithms for sorted ranges (`binary_search`, `lower_bound`, `upper_bound`, and `equal_range`) run in logarithmic time.

The shift from unsorted ranges to sorted ranges leads to another shift: from using equality to determine whether two values are the same to using equivalence. [Item 19](#) comprises a discourse on equality versus equivalence, so I won't repeat it here. Instead, I'll simply note that the `count` and `find` algorithms both search using equality, while `binary_search`, `lower_bound`, `upper_bound`, and `equal_range` employ equivalence.

To test for the existence of a value in a sorted range, use `binary_search`. Unlike `bsearch` in the standard C library (and hence also in the standard C++ library), `binary_search` returns only a `bool`: whether the value was found. `binary_search` answers the question, "Is it there?," and its answer is either yes or no. If you need more information than that, you need a different algorithm.

Here's an example of `binary_search` applied to a sorted vector. (You can read about the virtues of sorted vectors in [Item 23](#).)

```
vector<Widget> vw;                // create vector, put
...                               // data into it, sort the
sort(vw.begin(), vw.end());      // data
Widget w;                        // value to search for
...
if (binary_search(vw.begin(), vw.end(), w)) {
    ...                          // w is in vw
} else {
    ...                          // it's not
}
```

If you have a sorted range and your question is, “Is it there, and if so, where is it?” you want `equal_range`, but you may think you want `lower_bound`. We'll discuss `equal_range` shortly, but first, let's examine `lower_bound` as a way of locating values in a range.

When you ask `lower_bound` to look for a value, it returns an iterator pointing to either the first copy of that value (if it's found) or to the proper insertion location for that value (if it's not). `lower_bound` thus answers the question, “Is it there? If so, where is the first copy, and if it's not, where would it go?” As with `find`, the result of `lower_bound` must be tested to see if it's pointing to the value you're looking for. Unlike `find`, you can't just test `lower_bound`'s return value against the end iterator. Instead, you must test the object `lower_bound` identifies to see if that's the value you want.

Many programmers use `lower_bound` like this:

```
vector<Widget>::iterator i = lower_bound(vw.begin(), vw.end(), w);
if (i != vw.end() && *i == w) {    // make sure i points to an object;
    ...                            // make sure the object has the
    ...                            // correct value; this has a bug!
    ...                            // found the value, i points to the
    ...                            // first object with that value
} else {
    ...                            // not found
}
```

This works most of the time, but it's not really correct. Look again at the test to determine whether the desired value was found:

```
if (i != vw.end() && *i == w) ...
```

This is an *equality* test, but `lower_bound` searched using *equivalence*. Most of the time, tests for equivalence and equality yield the same results, but [Item 19](#) demonstrates that it's not that hard to come up

with situations where equality and equivalence are different. In such situations the code above is wrong.

To do things properly, you must check to see if the iterator returned from `lower_bound` points to an object with a value that is equivalent to the one you searched for. You could do that manually ([Item 19](#) shows you how, and [Item 24](#) provides an example of when it can be worthwhile), but it can get tricky, because you have to be sure to use the same comparison function that `lower_bound` used. In general, that could be an arbitrary function (or function object). If you passed a comparison function to `lower_bound`, you'd have to be sure to use the same comparison function in your hand-coded equivalence test. *That* would mean that if you changed the comparison function you passed to `lower_bound`, you'd have to make the corresponding change in your check for equivalence. Keeping the comparison functions in sync isn't rocket science, but it is another thing to remember, and I suspect you already have plenty you're expected to keep in mind.

There is an easier way: use `equal_range`. `equal_range` returns a *pair* of iterators, the first equal to the iterator `lower_bound` would return, the second equal to the one `upper_bound` would return (i.e., the one-past-the-end iterator for the range of values equivalent to the one searched for). `equal_range`, then, returns a pair of iterators that demarcate the range of values equivalent to the one you searched for. A well-named algorithm, no? (`equivalent_range` would be better, of course, but `equal_range` is still pretty good.)

There are two important observations about `equal_range`'s return value. First, if the two iterators are the same, that means the range of objects is empty; the value wasn't found. That observation is the key to using `equal_range` to answer the question, "Is it there?" You use it like this:

```
vector<Widget> vw;
...
sort(vw.begin(), vw.end());
typedef vector<Widget>::iterator VWIter; // convenience typedefs
typedef pair<VWIter, VWIter> VWIterPair;
VWIterPair p = equal_range(vw.begin(), vw.end(), w);
if (p.first != p.second) {                               // if equal_range didn't return
                                                         // an empty range...
    ...                                                    // found it, p.first points to the
                                                         // first one and p.second
                                                         // points to one past the last
} else {
```

```

...                                     // not found, both p.first and
                                     // p.second point to the
}                                     // insertion location for
                                     // the value searched for

```

This code uses only equivalence, so it is always correct.

The second thing to note about `equal_range`'s return value is that the distance between its iterators is the number of objects in the range, i.e., the objects with a value equivalent to the target of the search. For types where equality and equivalence yield the same results, then, `equal_range` does the job of both `find` and `count` for sorted ranges. To locate the Widgets in `vw` with a value equivalent to `w` and then print out how many such Widgets exist, for example, you could do this:

```

VWIterPair p = equal_range(vw.begin(), vw.end(), w);
cout << "There are " << distance(p.first, p.second)
    << " elements in vw equivalent to w.";

```

So far, our discussion has assumed we want to search for a value in a range, but sometimes we're more interested in finding a *location* in a range. For example, suppose we have a `Timestamp` class and a vector of `Timestamps` that's sorted so that older timestamps come first:

```

class Timestamp { ... };
bool operator<(const Timestamp& lhs,      // returns whether lhs
               const Timestamp& rhs);    // precedes rhs in time
vector<Timestamp> vt;                    // create vector, fill it with
...                                     // data, sort it so that older
sort(vt.begin(), vt.end());              // times precede newer ones

```

Now suppose we have a special timestamp, `ageLimit`, and we want to remove from `vt` all the timestamps that are older than `ageLimit`. In this case, we don't want to search `vt` for a `Timestamp` equivalent to `ageLimit`, because there might not be any elements with that exact value. Instead, we need to find a *location* in `vt`: the first element that is no older than `ageLimit`. This is as easy as can be, because `lower_bound` will give us precisely what we need:

```

Timestamp ageLimit;
...
vt.erase(vt.begin(), lower_bound(vt.begin(),      // eliminate from vt all
                                vt.end(),          // objects that precede
                                ageLimit));         // ageLimit's value

```

If our requirements change slightly so that we want to eliminate all the timestamps that are at least as old as `ageLimit`, we need to find the location of the first timestamp that is *younger* than `ageLimit`. That's a job tailor-made for `upper_bound`:

```

vt.erase(vt.begin(), upper_bound(vt.begin(),           // eliminate from vt all
                                  vt.end(),             // objects that precede
                                  ageLimit));           // or are equivalent
                                                        // to ageLimit's value

```

`upper_bound` is also useful if you want to insert things into a sorted range so that objects with equivalent values are stored in the order in which they were inserted. For example, we might have a sorted list of `Person` objects, where the objects are sorted by name:

```

class Person {
public:
    ...
    const string& name() const;
    ...
};

struct PersonNameLess:
    public binary_function<Person, Person, bool> {    // see Item 40
    bool operator()(const Person& lhs, const Person& rhs) const
    {
        return lhs.name() < rhs.name();
    }
};

list<Person> lp;
...
lp.sort(PersonNameLess());                          // sort lp using
                                                        // PersonNameLess

```

To keep the list sorted the way we desire (by name, with equivalent names stored in the order in which they are inserted), we can use `upper_bound` to specify the insertion location:

```

Person newPerson;
...
lp.insert( upper_bound(lp.begin(),                     // insert newPerson after
                      lp.end(),                       // the last object in lp
                      newPerson,                      // that precedes or is
                      PersonNameLess()),              // equivalent to
          newPerson);                                 // newPerson

```

This works fine and is quite convenient, but it's important not to be misled by this use of `upper_bound` into thinking that we're magically looking up an insertion location in a list in logarithmic time. We're not. [Item 34](#) explains that because we're working with a list, the lookup takes linear time, but it performs only a logarithmic number of comparisons.

Up to this point, we have considered only the case where you have a pair of iterators defining a range to be searched. Often you have a container, not a range. In that case, you must distinguish between the sequence and associative containers. For the standard sequence containers (vector, string, deque, and list), you follow the advice we've outlined in this Item, using the containers' begin and end iterators to demarcate the range.

The situation is different for the standard associative containers (set, multiset, map, and multimap), because they offer member functions for searching that are generally better choices than the STL algorithms. [Item 44](#) goes into the details of why they are better choices, but briefly, it's because they're faster and they behave more naturally. Fortunately, the member functions usually have the same names as the corresponding algorithms, so where the foregoing discussion recommends you choose algorithms named count, find, equal_range, lower_bound, or upper_bound, you simply select the same-named member functions when searching associative containers. binary_search calls for a different strategy, because there is no member function analogue to this algorithm. To test for the existence of a value in a set or map, use count in its idiomatic role as a test for membership:

```
set<Widget> s;                // create set, put data into it
...
Widget w;                    // w still holds the value to search for
...
if (s.count(w)) {
    ...                        // a value equivalent to w exists
} else {
    ...                        // no such value exists
}
```

To test for the existence of a value in a multiset or multimap, find is generally superior to count, because find can stop once it's found a single object with the desired value, while count, in the worst case, must examine every object in the container. (This isn't a problem when working with sets and maps, because sets don't allow duplicate values and maps don't allow duplicate keys.)

However, count's role for counting things in associative containers is secure. In particular, it's a better choice than calling equal_range and applying distance to the resulting iterators. For one thing, it's clearer: count means "count." For another, it's easier; there's no need to create a pair and pass its components to distance. For a third, it's probably a little faster.

Given everything we've considered in this Item, where do we stand? The following table says it all.

What You Want to Know	Algorithm to Use		Member Function to Use	
	On an Unsorted Range	On a Sorted Range	With a set or map	With a multiset or multimap
Does the desired value exist?	find	binary_search	count	find
Does the desired value exist? If so, where is the first object with that value?	find	equal_range	find	find or lower_bound (see below)
Where is the first object with a value not preceding the desired value?	find_if	lower_bound	lower_bound	lower_bound
Where is the first object with a value succeeding the desired value?	find_if	upper_bound	upper_bound	upper_bound
How many objects have the desired value?	count	equal_range, then distance	count	count
Where are all the objects with the desired value?	find (iteratively)	equal_range	equal_range	equal_range

In the column summarizing how to work with sorted ranges, the frequency with which `equal_range` occurs may be surprising. That frequency arises from the importance of testing for equivalence when searching. With `lower_bound` and `upper_bound`, it's too easy to fall back on equality tests, but with `equal_range`, testing only for equivalence is the natural thing to do. In the second row for sorted ranges, `equal_range` beats out `find` for an additional reason: `equal_range` runs in logarithmic time, while `find` takes linear time.

For multisets and multimaps, the table lists both `find` and `lower_bound` as candidates when you're looking for the first object with a particular value. `find` is a common choice for this job, and you may have noticed that it's the one listed in the table for sets and maps. For the multi containers, however, `find` is not guaranteed to identify the *first* element in the container with a given value if more than one is present; its character is only to identify *one* of those elements. If you really need to find the *first* object with a given value, you'll want to employ `lower_bound`, and you'll have to manually perform the second half of the equivalence test described in [Item 19](#) to confirm that you've found the value you were looking for. (You could avoid the manual equivalence test by using `equal_range`, but calling `equal_range` is more expensive than calling `lower_bound`.)

Selecting among `count`, `find`, `binary_search`, `lower_bound`, `upper_bound`, and `equal_range` is easy. Choose the algorithm or member function that offers you the behavior and performance you need and that requires the least amount of work when you call it. Follow that advice (or consult the table), and you should never get confused.

Item 46: Consider function objects instead of functions as algorithm parameters.

One of the complaints about programming in high-level languages is that as the level of abstraction gets higher, the efficiency of the generated code gets lower. In fact, Alexander Stepanov (the inventor of the STL) once produced a small benchmark suite that tried to measure the “abstraction penalty” of C++ vis-à-vis C. Among other things, the results of that benchmark revealed that it's nearly universal for the code generated for manipulating a class containing a double to be less efficient than the corresponding code for manipulating a double directly. It may thus come as a surprise to learn that passing STL function objects — objects masquerading as functions — to algorithms typically yields code that is *more* efficient than passing real functions.

For example, suppose you need to sort a vector of doubles in descending order. The straightforward STL way to do it is via the `sort` algorithm and a function object of type `greater<double>`:

```
vector<double> v;  
...  
sort(v.begin(), v.end(), greater<double>());
```

If you're wary of the abstraction penalty, you might decide to eschew the function object in favor of a real function, a function that's not only real, it's inline:

```
inline
bool doubleGreater(double d1, double d2)
{
    return d1 > d2;
}

...
sort(v.begin(), v.end(), doubleGreater);
```

Interestingly, if you were to compare the performance of the two calls to `sort` (one using `greater<double>`, one using `doubleGreater`), you'd almost certainly find that the one using `greater<double>` was faster. For instance, I timed the two calls to `sort` on a vector of a million doubles using four different STL platforms, each set to optimize for speed, and the version using `greater<double>` was faster every time. At worst, it was 50% faster, at best it was 160% faster. So much for the abstraction penalty.

The explanation for this behavior is simple: inlining. If a function object's `operator()` function has been declared inline (either explicitly via inline or implicitly by defining it in its class definition), the body of that function is available to compilers, and most compilers will happily inline that function during template instantiation of the called algorithm. In the example above, `greater<double>::operator()` is an inline function, so compilers inline-expand it during instantiation of `sort`. As a result, `sort` contains zero function calls, and compilers are able to perform optimizations on this call-free code that are otherwise not usually attempted. (For a discussion of the interaction between inlining and compiler optimization, see Item 33 of *Effective C++* and chapters 8–10 of Bulka and Mayhew's *Efficient C++* [10].)

The situation is different for the call to `sort` using `doubleGreater`. To see how it's different, we must recall that there's no such thing as passing a function as a parameter to another function. When we try to pass a function as a parameter, compilers silently convert the function into a *pointer* to that function, and it's the pointer we actually pass. Hence, the call

```
sort(v.begin(), v.end(), doubleGreater);
```

doesn't pass `doubleGreater` to `sort`, it passes a pointer to `doubleGreater`. When the `sort` template is instantiated, this is the declaration for the function that is generated:

```
void sort(vector<double>::iterator first,      // beginning of range
         vector<double>::iterator last,      // end of range
         bool (*comp)(double, double));      // comparison function
```

Because `comp` is a pointer to a function, each time it's used inside `sort`, compilers make an indirect function call — a call through a pointer. Most compilers won't try to inline calls to functions that are invoked through function pointers, even if, as in this example, such functions have been declared inline and the optimization appears to be straightforward. Why not? Probably because compiler vendors have never felt that it was worthwhile to implement the optimization. You have to have a little sympathy for compiler vendors. They have lots of demands on their time, and they can't do everything. Not that this should stop you from asking them for it.

The fact that function pointer parameters inhibit inlining explains an observation that long-time C programmers often find hard to believe: C++'s `sort` virtually always embarrasses C's `qsort` when it comes to speed. Sure, C++ has function and class templates to instantiate and funny-looking operator() functions to invoke while C makes a simple function call, but all that C++ "overhead" is absorbed during compilation. At runtime, `sort` makes inline calls to its comparison function (assuming the comparison function has been declared inline and its body is available during compilation) while `qsort` calls its comparison function through a pointer. The end result is that `sort` runs faster. In my tests on a vector of a million doubles, it ran up to 670% faster, but don't take my word for it, try it yourself. It's easy to verify that when comparing function objects and real functions as algorithm parameters, there's an abstraction *bonus*.

There's another reason to prefer function objects to functions as algorithm parameters, and it has nothing to do with efficiency. It has to do with getting your programs to compile. For whatever reason, it's not uncommon for STL platforms to reject perfectly valid code, either through shortcomings in the compiler or the library or both. For example, one widely used STL platform rejects the following (valid) code to print to `cout` the length of each string in a set:

```
set<string> s;
...
transform(s.begin(), s.end(),
         ostream_iterator<string::size_type>(cout, "\n"),
         mem_fun_ref(&string::size));
```

The cause of the problem is that this particular STL platform has a bug in its handling of const member functions (such as `string::size`). A workaround is to use a function object instead:

```
struct StringSize:
    public unary_function<string, string::size_type> {           // see Item 40
    string::size_type operator()(const string& s) const
    {
        return s.size();
    }
};

transform(s.begin(), s.end(),
          ostream_iterator<string::size_type>(cout, "\n"),
          StringSize());
```

There are other workarounds for this problem, but this one does more than just compile on every STL platform I know. It also facilitates inlining the call to `string::size`, something that would almost certainly not take place in the code above where `mem_fun_ref(&string::size)` is passed to `transform`. In other words, creation of the functor class `StringSize` does more than sidestep compiler conformance problems, it's also likely to lead to an increase in performance.

Another reason to prefer function objects to functions is that they can help you avoid subtle language pitfalls. Occasionally, source code that looks reasonable is rejected by compilers for legitimate, but obscure, reasons. There are situations, for example, when the name of an instantiation of a function template is not equivalent to the name of a function. Here's one such situation:

```
template<typename FPType>                // return the average
FPType average(FPType val1, FPType val2) // of 2 floating point
{                                         // numbers
    return (val1 + val2) / 2;
}

template<typename InputIter1,
         typename InputIter2>
void writeAverages(InputIter1 begin1,        // write the pairwise
                  InputIter1 end1,          // averages of 2
                  InputIter2 begin2,        // sequences to a
                  ostream& s)               // stream
{
    transform(
        begin1, end1, begin2,
        ostream_iterator<typename iterator_traits<InputIter1>::value_type>(s, "\n"),
```

```

        average<typename iterator_traits<InputIter1>::value_type> // error?
    );
}

```

Many compilers accept this code, but the C++ Standard appears to forbid it.[†] The reasoning is that there could, in theory, be another function template named `average` that takes a single type parameter. If there were, the expression `average<typename iterator_traits<InputIter1>::value_type>` would be ambiguous, because it would not be clear which template to instantiate. In this particular example, no ambiguity is present, but some compilers reject the code anyway, and they are allowed to do that. No matter. The solution to the problem is to fall back on a function object:

```

template<typename FType>
struct Average:
    public binary_function<FType, FType, FType> {           // see Item 40
    FType operator()(FType val1, FType val2) const
    {
        return average(val1, val2);
    }
};

template<typename InputIter1, typename InputIter2>
void writeAverages(InputIter1 begin1, InputIter1 end1,
                  InputIter2 begin2, ostream& s)
{
    transform(
        begin1, end1, begin2,
        ostream_iterator<typename iterator_traits<InputIter1>::value_type>(s, "\n"),
        Average<typename iterator_traits<InputIter1>::value_type>());
};
}

```

Every compiler should accept this revised code. Furthermore, calls to `Average::operator()` inside `transform` are eligible for inlining, something that would not be true for an instantiation of `average` above, because `average` is a template for functions, not function *objects*.

Function objects as parameters to algorithms thus offer more than greater efficiency. They're also more robust when it comes to getting your code to compile. Real functions are useful, of course, but when it comes to effective STL programming, function objects are frequently more useful.

[†] In October 2003, the Standardization Committee clarified that this code should be valid. This decision doesn't officially take effect until the next version of the C++ standard is adopted (anticipated to be around 2011), but compiler vendors are likely to implement the revised rules sooner.

Item 47: Avoid producing write-only code.

Suppose you have a `vector<int>`, and you'd like to get rid of all the elements in the vector whose value is less than `x` and that occur after the last value at least as big as `y`. Does the following instantly spring to mind?

```
vector<int> v;
int x, y;

...
v.erase(
    remove_if(find_if(v.rbegin(), v.rend(),
                      bind2nd(greater_equal<int>(), y)).base(),
              v.end(),
              bind2nd(less<int>(), x)),
    v.end());
```

One statement, and the job is done. Clear and straightforward. No problem. Right?

Well, let's step back for a moment. Does this strike you as reasonable, maintainable code? "No!" shriek most C++ programmers, fear and loathing in their voices. "Yes!" squeal a few, delight evident in theirs. And therein lies the problem. One programmer's vision of expressive purity is another programmer's demonic missive from Hell.

As I see it, there are two causes for concern in the code above. First, it's a rat's nest of function calls. To see what I mean, here is the same statement, but with all the function names replaced by *fn*, each *n* corresponding to one of the functions:

```
v.f1(f2(f3(v.f4(), v.f5(), f6(f7(), y)).f8(), v.f9(), f6(f10(), x)), v.f9());
```

This looks unnaturally complicated, because I've removed the indentation present in the original example, but I think it's safe to say that any statement involving twelve function calls to ten different functions would be considered excessive by most C++ software developers. Programmers weaned on functional languages such as Scheme might feel differently, however, and my experience has been that the majority of programmers who view the original code without raising an eyebrow have a strong functional programming background. Most C++ programmers lack this background, so unless your colleagues are versed in the ways of deeply nested function calls, code like the `erase` call above is almost sure to confound the next person who is forced to make sense of what you have written.

The second drawback of the code is the significant STL background needed to understand it. It uses the less common `_if` forms of `find` and

remove, it uses reverse iterators (see [Item 26](#)), it converts reverse_iterators to iterators (see [Item 28](#)), it uses bind2nd, it creates anonymous function objects, and it employs the erase-remove idiom (see [Item 32](#)). Experienced STL programmers can swallow that combination without difficulty, but far more C++ developers will have their eyes glaze over before they've taken so much as a bite. If your colleagues are well-steeped in the ways of the STL, using erase, remove_if, find_if, base, and bind2nd in a single statement may be fine, but if you want your code to be comprehensible by C++ programmers with a more mainstream background, I encourage you to break it down into more easily digestible chunks.

Here's one way you could do it. (The comments aren't just for this book. I'd put them in the code, too.)

```
typedef vector<int>::iterator VecIntIter;

// Initialize rangeBegin to point to the element following the last
// occurrence of a value greater than or equal to y. If there is no such
// value, initialize rangeBegin to v.begin(). If the last occurrence of the
// value is the last element in v, initialize rangeBegin to v.end().
VecIntIter rangeBegin = find_if(v.rbegin(), v.rend(),
                               bind2nd(greater_equal<int>(), y)).base();

// from rangeBegin to v.end(), erase everything with a value less than x
v.erase(remove_if(rangeBegin, v.end(), bind2nd(less<int>(), x)), v.end());
```

This is still likely to confuse some people, because it relies on an understanding of the erase-remove idiom, but between the comments in the code and a good STL reference (e.g., Josuttis' *The C++ Standard Library* [3] or SGI's STL web site [21]), every C++ programmer should be able to figure out what's going on without too much difficulty.

When transforming the code, it's important to note that I didn't abandon algorithms and try to write my own loops. [Item 43](#) explains why that's generally an inferior option, and its arguments apply here. When writing source code, the goal is to come up with code that is meaningful to both compilers and humans and that offers acceptable performance. Algorithms are almost always the best way to achieve that goal. However, [Item 43](#) also explains how the increased use of algorithms naturally leads to an increased tendency to nest function calls and to throw in binders and other functor adapters. Look again at the problem specification that opened this Item:

Suppose you have a vector<int>, and you'd like to get rid of all the elements in the vector whose value is less than x and that occur after the last value at least as big as y.

The outline of a solution *does* spring to mind:

- Finding the last occurrence of a value in a vector calls for some application of `find` or `find_if` with reverse iterators.
- Getting rid of elements calls for either `erase` or the erase-remove idiom.

Put those two ideas together, and you get this pseudocode, where “something” indicates a placeholder for code that hasn’t yet been fleshed out:

```
v.erase(remove_if(find_if(v.rbegin(), v.rend(), something).base(),
                      v.end(),
                      something)),
        v.end());
```

Once you’ve got that, figuring out the *somethings* isn’t terribly difficult, and the next thing you know, you have the code in the original example. That’s why this kind of statement is often known as “write-only” code. As you write the code, it seems straightforward, because it’s a natural outgrowth of some basic ideas (e.g., the erase-remove idiom plus the notion of using `find` with reverse iterators). *Readers*, however, have great difficulty in decomposing the final product back into the ideas on which it is based. That’s the calling card of write-only code: it’s easy to write, but it’s hard to read and understand.

Whether code is write-only depends on who’s reading it. As I noted, some C++ programmers think nothing of the code in this Item. If that’s typical in the environment in which you work and you expect it to be typical in the future, feel free to unleash your most advanced STL programming inclinations. However, if your colleagues are less comfortable with a functional programming style and are less experienced with the STL, scale back your ambitions and write something more along the lines of the two-statement alternative I showed earlier.

It’s a software engineering truism that code is read more often than it is written. Equally well established is that software spends far more time in maintenance than it does in development. Software that cannot be read and understood cannot be maintained, and software that cannot be maintained is hardly worth having. The more you work with the STL, the more comfortable you’ll become with it, and the more you’ll feel the pull to nest function calls and create function objects on the fly. There’s nothing wrong with that, but always bear in mind that the code you write today will be read by somebody — possibly you — someday in the future. Prepare for that day.

Use the STL, yes. Use it well. Use it effectively. But avoid producing write-only code. In the long run, such code is anything *but* effective.

Item 48: Always #include the proper headers.

Among the minor frustrations of STL programming is that it is easy to create software that compiles on one platform, yet requires additional `#include` directives on others. This annoyance stems from the fact that the Standard for C++ (unlike the Standard for C) fails to dictate which standard headers must or may be `#included` by other standard headers. Given such flexibility, different implementers have chosen to do different things.

To give you some idea of what this means in practice, I sat down one day with five STL platforms (let's call them A, B, C, D, and E), and I spent a little time throwing toy programs at them to see which standard headers I could omit and still get a successful compilation. This indirectly told me which headers `#include` other headers. This is what I found:

- With A and C, `<vector>` `#includes` `<string>`.
- With C, `<algorithm>` `#includes` `<string>`.
- With C and D, `<iostream>` `#includes` `<iterator>`.
- With D, `<iostream>` `#includes` `<string>` and `<vector>`.
- With D and E, `<string>` `#includes` `<algorithm>`.
- With all five implementations, `<set>` `#includes` `<functional>`.

Except for the case of `<set>` `#includeing` `<functional>`, I didn't find a way to get a program with a missing header past implementation B. According to Murphy's Law, then, you will always develop under a platform like A, C, D, or E and you will always be porting to a platform like B, especially when the pressure for the port is greatest and the time to accomplish it is least. Naturally.

But don't blame your compilers or library implementations for your porting woes. It's your fault if you're missing required headers. Any time you refer to elements of namespace `std`, you are responsible for having `#included` the appropriate headers. If you omit them, your code might compile anyway, but you'll still be missing necessary headers, and other STL platforms may justly reject your code.

To help you remember what's required when, here's a quick summary of what's in each standard STL-related header:

- Almost all the containers are declared in headers of the same name, i.e., `vector` is declared in `<vector>`, `list` is declared in `<list>`, etc. The exceptions are `<set>` and `<map>`. `<set>` declares both `set` and `multiset`, and `<map>` declares both `map` and `multimap`.

- All but four algorithms are declared in `<algorithm>`. The exceptions are `accumulate` (see [Item 37](#)), `inner_product`, `adjacent_difference`, and `partial_sum`. Those algorithms are declared in `<numeric>`.
- Special kinds of iterators, including `istream_iterators` and `istreambuf_iterators` (see [Item 29](#)), are declared in `<iterator>`.
- Standard functors (e.g., `less<T>`) and functor adapters (e.g., `not1`, `bind2nd`) are declared in `<functional>`.

Any time you use any of the components in a header, be sure to provide the corresponding `#include` directive, even if your development platform lets you get away without it. Your diligence will pay off in reduced stress when you find yourself porting to a different platform.

Item 49: Learn to decipher STL-related compiler diagnostics.

It's perfectly legal to define a vector with a particular size,

```
vector<int> v(10);           // create a vector of size 10
```

and strings act a lot like vectors, so you might expect to be able to do this:

```
string s(10);               // attempt to create a string of size 10
```

This won't compile. There is no string constructor taking an int argument. One of my STL platforms tells me that like this:

```
example.cpp(20) : error C2664: '___thiscall std::basic_string<char,struct
std::char_traits<char>,class std::allocator<char>
>::std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> >(const class std::allocator<char> &)' : cannot convert
parameter 1 from 'const int' to 'const class std::allocator<char> &'
Reason: cannot convert from 'const int' to 'const class std::allocator<char>'
No constructor could take the source type, or constructor overload
resolution was ambiguous
```

Isn't that wonderful? The first part of the message looks as if a cat walked across the keyboard, the second part mysteriously refers to an allocator never mentioned in the source code, and the third part says the constructor call is bad. The third part is accurate, of course, but let's first focus our attention on the result of the purported feline stroll, because it's representative of diagnostics you'll frequently see when using strings.

string isn't a class, it's a typedef. In particular, it's a typedef for this:

```
basic_string<char, char_traits<char>, allocator<char> >
```

That's because the C++ notion of a string has been generalized to mean sequences of arbitrary character types with arbitrary character characteristics ("traits") and stored in memory allocated by arbitrary allocators. All string-like objects in C++ are really instantiations of the template `basic_string`, and that's why most compilers refer to the type `basic_string` when they issue diagnostics about programs making erroneous use of strings. (A few compilers are kind enough to use the name `string` in diagnostics, but most aren't.) Often, such diagnostics will explicitly note that `basic_string` (and the attendant helper templates `char_traits` and `allocator`) are in the `std` namespace, so it's not uncommon to see errors involving strings yield diagnostics that mention this type:

```
std::basic_string<char, std::char_traits<char>, std::allocator<char> >
```

This is quite close to what's used in the compiler diagnostic above, but different compilers use variations on the theme. Another STL platform I use refers to strings this way,

```
basic_string<char, string_char_traits<char>, __default_alloc_template<false,0> >
```

The names `string_char_traits` and `__default_alloc_template` are nonstandard, but that's life. Some STL implementations deviate from the standard. If you don't like the deviations in your current STL implementation, consider replacing it with a different one. [Item 50](#) gives examples of places you can go for alternative implementations.

Regardless of how a compiler diagnostic refers to the string type, the technique for reducing the diagnostic to something meaningful is the same: globally replace the `basic_string` gobbledegook with the text "string". If you're using a command-line compiler, it's usually easy to do this with a program like `sed` or a scripting language like `perl`, `python`, or `ruby`. (You'll find an example of such a script in Zolman's article, "An STL Error Message Decryptor for Visual C++" [\[26\]](#).) In the case of the diagnostic above, we globally replace

```
std::basic_string<char, struct std::char_traits<char>, class std::allocator<char> >
```

with `string` and we end up with this:

```
example.cpp(20) : error C2664: '__thiscall string::string(const class
std::allocator<char> &)' : cannot convert parameter 1 from 'const int' to
'const class std::allocator<char> &'
```

This makes clear (or at least clearer) that the problem is in the type of the parameter passed to the `string` constructor, and even though the mysterious reference to `allocator<char>` remains, it should be easy to

look up the constructor forms for `string` to see that none exists taking only a size.

By the way, the reason for the mysterious reference to an allocator is that each standard container has a constructor taking only an allocator. In the case of `string`, it's one of three constructors that can be called with one argument, but for some reason, this compiler figures that the one taking an allocator is the one you're trying to call. The compiler figures wrong, and the diagnostic is misleading. Oh well.

As for the constructor taking only an allocator, please don't use it. That constructor makes it easy to end up with containers of the same type but with inequivalent allocators. In general, that's bad. Very bad. To find out why, turn to [Item 11](#).

Now let's tackle a more challenging diagnostic. Suppose you're implementing an email program that allows users to refer to people by nicknames instead of by email addresses. For example, such a program would make it possible to use "The Big Cheese" as a synonym for the email address of the President of the United States (which happens to be president@whitehouse.gov). Such a program might use a map from nicknames to email addresses, and it might offer a member function `showEmailAddress` that displays the email address associated with a given nickname:

```
class NiftyEmailProgram {
private:
    typedef map<string, string> NicknameMap;
    NicknameMap nicknames;           // map from nicknames to
                                    // email addresses
public:
    ...
    void showEmailAddress(const string& nickname) const;
};
```

Inside `showEmailAddress`, you'll need to find the map entry associated with a particular nickname, so you might write this:

```
void NiftyEmailProgram::showEmailAddress(const string& nickname) const
{
    ...
    NicknameMap::iterator i = nicknames.find(nickname);
    if (i != nicknames.end()) ...
    ...
}
```

Compilers don't like this, and with good reason, but the reason isn't obvious. To help you figure it out, here's what one STL platform helpfully emits:

```
example.cpp(17) : error C2440: 'initializing' : cannot convert from 'class std::_Tree<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,struct
std::pair<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > const ,class std::basic_string<char,struct
std::char_traits<char>,class std::allocator<char> > >,struct std::map<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,struct
std::less<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > >,class std::allocator<class std::basic_string<char,struct
std::char_traits<char>,class std::allocator<char> > > >::_Kfn,struct std::less<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> > >,class
std::allocator<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > > >::const_iterator' to 'class std::_Tree<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,struct
std::pair<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > const ,class std::basic_string<char,struct
std::char_traits<char>,class std::allocator<char> > >,struct std::map<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> >,struct
std::less<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > >,class std::allocator<class std::basic_string<char,struct
std::char_traits<char>,class std::allocator<char> > > >::_Kfn,struct std::less<class
std::basic_string<char,struct std::char_traits<char>,class std::allocator<char> > >,class
std::allocator<class std::basic_string<char,struct std::char_traits<char>,class
std::allocator<char> > > >::iterator'
```

No constructor could take the source type, or constructor overload resolution was ambiguous

At 2095 characters long, this message looks fairly gruesome, but I've seen worse. One of my favorite STL platforms produces a diagnostic of 4812 characters for this example. As you might guess, features other than its error messages are what have engendered my fondness for it.

Let's reduce this mess to something manageable. We begin with the replacement of the `basic_string` gibberish with `string`. That yields this:

```
example.cpp(17) : error C2440: 'initializing' : cannot convert from 'class
std::_Tree<class string,struct std::pair<class string const ,class string
>,struct std::map<class string,class string,struct std::less<class string
>,class std::allocator<class string > >::_Kfn,struct std::less<class string
>,class std::allocator<class string > >::const_iterator' to 'class
std::_Tree<class string,struct std::pair<class string const ,class string
>,struct std::map<class string,class string,struct std::less<class string
>,class std::allocator<class string > >::_Kfn,struct std::less<class string
>,class std::allocator<class string > >::iterator'
```

No constructor could take the source type, or constructor overload resolution was ambiguous

Much better. Now a svelte 745 characters long, we can start to actually look at the message. One of the things that is likely to catch our eye is the mention of the template `std::_Tree`. The Standard says nothing about a template called `_Tree`, but the leading underscore in the name followed by a capital letter jogs our memory that such names

are reserved for implementers. This is an internal template used to implement some part of the STL.

In fact, almost all STL implementations use some kind of underlying template to implement the standard associative containers (set, multiset, map, and multimap). In the same way that source code using string typically leads to diagnostics mentioning `basic_string`, source code using a standard associative container often leads to diagnostics mentioning some underlying tree template. In this case, it's called `_Tree`, but other implementations I know use `__tree` or `__rb_tree`, the latter reflecting the use of red-black trees, the most common type of balanced tree used in STL implementations.

Setting `_Tree` aside for a moment, the message above mentions a type we should recognize: `std::map<class string,class string,struct std::less<class string >,class std::allocator<class string > >`. This is precisely the type of map we are using, except that the comparison and allocator types (which we chose not to specify when we defined the map) are shown. The error message will be easier to understand if we replace that type with our typedef for it, `NicknameMap`. That leads to this:

```
example.cpp(17) : error C2440: 'initializing' : cannot convert from 'class
std::_Tree<class string,struct std::pair<class string const ,class string
>,struct NicknameMap::_Kfn,struct std::less<class string >,class
std::allocator<class string > >::const_iterator' to 'class std::_Tree<class
string,struct std::pair<class string const ,class string >,struct
NicknameMap::_Kfn,struct std::less<class string >,class std::allocator<class
string > >::iterator'
```

No constructor could take the source type, or constructor overload resolution was ambiguous

This message is shorter, but not much clearer. We need to do something with `_Tree`. Because `_Tree` is an implementation-specific template, the only way to know the meaning of its template parameters is to read the source code, and there's no reason to go rummaging through implementation-specific source code if we don't have to. Let's try simply replacing all the stuff passed to `_Tree` with `SOMETHING` to see what we get. This is the result:

```
example.cpp(17) : error C2440: 'initializing' : cannot convert from 'class
std::_Tree<SOMETHING>::const_iterator' to 'class
std::_Tree<SOMETHING>::iterator'
```

No constructor could take the source type, or constructor overload resolution was ambiguous

This is something we can work with. The compiler is complaining that we're trying to convert some kind of `const_iterator` into an iterator, a

clear violation of `const` correctness. Let's look again at the offending code, where I've highlighted the line raising the compiler's ire:

```
class NiftyEmailProgram {
private:
    typedef map<string, string> NicknameMap;
    NicknameMap nicknames;
public:
    ...
    void showEmailAddress(const string& nickname) const;
};

void NiftyEmailProgram::showEmailAddress(const string& nickname) const
{
    ...
    NicknameMap::iterator i = nicknames.find(nickname);
    if (i != nicknames.end()) ...
    ...
}
```

The only interpretation that makes any sense is that we're trying to initialize `i` (which is an iterator) with a `const_iterator` returned from `map::find`. That seems odd, because we're calling `find` on `nicknames`, and `nicknames` is a non-`const` object. `find` should thus return a non-`const` iterator.

Look again. Yes, `nicknames` is declared as a non-`const` map, but `showEmailAddress` is a *const* member function, and inside a `const` member function, all non-static data members of the class become `const`! Inside `showEmailAddress`, `nicknames` is a `const` map. Suddenly the error message makes sense. We're trying to generate an iterator into a map we've promised not to modify. To fix the problem, we must either make `i` a `const_iterator` or we must make `showEmailAddress` a non-`const` member function. Both solutions are probably less challenging than ferreting out the meaning of the error message.

In this Item, I've shown textual substitutions to reduce the complexity of error messages, but once you've practiced a little, you'll be able to perform the substitutions in your head most of the time. I'm no musician (I have trouble turning on the radio), but I'm told that good musicians can sight-read several bars at a glance; they don't need to look at individual notes. Experienced STL programmers develop a similar skill. They can internally translate things like `std::basic_string<char,struct std::char_traits<char>,class std::allocator<char>>` into `string` without thinking about it. You, too, will develop this skill, but until you do, remember that you can almost always reduce compiler diagnostics to something comprehensible by replacing lengthy template-based type names with shorter mnemonics. In

many cases, all you have to do is replace typedef expansions with typedef names you're already using. That's what we did when we replaced `std::map<class string, class string, struct std::less<class string >, class std::allocator<class string >>` with `NicknameMap`.

Here are a few other hints that should help you make sense of STL-related compiler messages:

- For vector and string, iterators are sometimes pointers, so compiler diagnostics may refer to pointer types if you've made a mistake with an iterator. For example, if your source code refers to `vector<double>::iterators`, compiler messages will sometimes mention `double*` pointers. (A noteworthy exception is when you're using the STL implementation from STLport and you're running in debug mode. In that case, vector and string iterators are definitely not pointers. For more on STLport and its debug mode, turn to [Item 50](#).)
- Messages mentioning `back_inserter`, `front_inserter`, or `insert_iterator` almost always mean you've made a mistake calling `back_inserter`, `front_inserter`, or `inserter`, respectively. (`back_inserter` returns an object of type `back_inserter_iterator`, `front_inserter` returns an object of type `front_inserter_iterator`, and `inserter` returns an object of type `insert_iterator`. For information on the use of these inserters, consult [Item 30](#).) If you didn't call these functions, some function you called (directly or indirectly) did.
- Similarly, if you get a message mentioning `binder1st` or `binder2nd`, you've probably made a mistake using `bind1st` or `bind2nd`. (`bind1st` returns an object of type `binder1st`, and `bind2nd` returns an object of type `binder2nd`.)
- Output iterators (e.g., `ostream_iterators`, `ostreambuf_iterators` (see [Item 29](#)), and the iterators returned from `back_inserter`, `front_inserter`, and `inserter`) do their outputting or inserting work inside assignment operators, so if you've made a mistake with one of these iterator types, you're likely to get a message complaining about something inside an assignment operator you've never heard of. To see what I mean, try compiling this code:


```
vector<string*> v;                                // try to print a container
copy(v.begin(), v.end(),                        // of string* pointers as
    ostream_iterator<string>(cout, "\n"));      // string objects
```
- If you get an error message originating from inside the implementation of an STL algorithm (i.e., the source code giving rise to the error is in `<algorithm>`), there's probably something wrong with the types you're trying to use with that algorithm. For example, you

may be passing iterators of the wrong category. To see how such usage errors are reported, edify (and amuse!) yourself by feeding this to your compilers:

```
list<int>::iterator i1, i2;           // pass bidirectional iterators to
sort(i1, i2);                       // an algorithm requiring random
                                   // access iterators
```

- If you're using a common STL component like vector, string, or the for_each algorithm, and a compiler says it has no idea what you're talking about, you've probably failed to #include a required header file. As [Item 48](#) explains, this problem can befall code that has been compiling smoothly for quite some time if you port it to a new platform.

Item 50: Familiarize yourself with STL-related web sites.

The Internet is rife with STL information. Ask your favorite search engine to look for “STL”, and it's sure to return hundreds of links, some of which may actually be relevant. For most STL programmers, however, no searching is necessary. The following sites are likely to rise to the top of almost everybody's most-frequently-used list:

- **The SGI STL site**, <http://www.sgi.com/tech/stl/>.
- **The STLport site**, <http://www.stlport.org/>.
- **The Boost site**, <http://www.boost.org/>.

What follows are brief descriptions of why these sites are worth bookmarking.

The SGI STL Web Site

SGI's STL web site tops the list, and for good reason. It offers comprehensive documentation on every component of the STL. For many programmers, this site is their on-line reference manual, regardless of which STL platform they are using. (The reference documentation was put together by Matt Austern, who later extended and polished it for his *Generic Programming and the STL* [4].) The material here covers more than just the STL components themselves. *Effective STL's* discussion of thread safety in STL containers (see [Item 12](#)), for example, is based on the treatment of the topic at the SGI STL web site.

The SGI site offers something else for STL programmers: a freely downloadable implementation of the STL. This implementation has been ported to only a handful of compilers, but the SGI distribution is

also the basis for the widely ported STLport distribution, about which I write more in a moment. Furthermore, the SGI implementation of the STL offers a number of nonstandard components that can make STL programming even more powerful, flexible, and fun. Foremost among these are the following:

- The **hashed associative containers** `hash_set`, `hash_multiset`, `hash_map`, and `hash_multimap`. For more information about these containers, turn to [Item 25](#).
- A **singly linked list container**, `slist`. This is implemented as you'd imagine, and iterators point to the list nodes you'd expect them to point to. Unfortunately, this makes it expensive to implement the insert and erase member functions, because both require adjustment of the next pointer of the node *preceding* the node pointed to by the iterator. In a doubly linked list (such as the standard list container), this isn't a problem, but in a singly linked list, going "back" one node is a linear-time operation. For SGI's `slist`, insert and erase take linear instead of constant time, a considerable drawback. SGI addresses the problem through the nonstandard (but constant-time) member functions `insert_after` and `erase_after`. Notes SGI,

If you find that `insert_after` and `erase_after` aren't adequate for your needs and that you often need to use insert and erase in the middle of the list, you should probably use `list` instead of `slist`.

Dinkumware also offers a singly linked list container called `slist`, but it uses a different iterator implementation that preserves the constant-time performance of insert and erase. For more information on Dinkumware, consult [Appendix B](#).

- A **string-like container for very large strings**. The container is called `rope`, because a rope is a heavy-duty string, don't you see? SGI describes ropes this way:

Ropes are a scalable string implementation: they are designed for efficient operations that involve the string as a whole. Operations such as assignment, concatenation, and substring take time that is nearly independent of the length of the string. Unlike C strings, ropes are a reasonable representation for very long strings, such as edit buffers or mail messages.

Under the hood, ropes are implemented as trees of reference-counted substrings, and each substring is stored as a char array. One interesting aspect of the rope interface is that the begin and

end member functions always return `const_iterator`s. This is to discourage clients from performing operations that change individual characters. Such operations are expensive. ropes are optimized for actions that involve entire strings (e.g., assignment, concatenation, and taking substrings, as mentioned above); single-character operations perform poorly.

- A variety of **nonstandard function objects and adapters**. The original HP STL implementation included more functor classes than made it into standard C++. Two of the more widely missed by old-time STL hackers are `select1st` and `select2nd`, because they are so useful for working with maps and multimaps. Given a pair, `select1st` returns its first component and `select2nd` returns its second. These nonstandard functor class templates can be used as follows:

```
map<int, string> m;

...
// write all the map keys to cout
transform(m.begin(), m.end(),
          ostream_iterator<int>(cout, "\n"),
          select1st<map<int, string>::value_type>());

// create a vector and copy all the values in the map into it
vector<string> v;
transform(m.begin(), m.end(), back_inserter(v),
          select2nd<map<int, string>::value_type>());
```

As you can see, `select1st` and `select2nd` make it easy to use algorithm calls in places where you might otherwise have to write your own loops (see [Item 43](#)), but, if you use these functors, the fact that they are nonstandard leaves you open to the charge that you are writing unportable and unmaintainable code (see [Item 47](#)). Diehard STL aficionados don't care. They consider it an injustice that `select1st` and `select2nd` didn't make it into the Standard in the first place.

Other nonstandard function objects that are part of the SGI implementation include `identity`, `project1st`, `project2nd`, `compose1` and `compose2`. To find out what these do, you'll have to visit the web site, though you'll find an example use of `compose2` on [page 187](#) of this book. By now, I hope it's clear that visiting the SGI web site will certainly be rewarding.

SGI's library implementation goes beyond the STL. Their goal is the development of a complete implementation of the standard C++ library, except for the parts inherited from C. (SGI assumes you

already have a standard C library at your disposal.) As a result, another noteworthy download available from SGI is an implementation of the C++ iostreams library. As you might expect, this implementation integrates well with SGI's implementation of the STL, but it also features performance that's superior to that of many iostream implementations that ship with C++ compilers.

The STLport Web Site

STLport's primary selling point is that it offers a modified version of SGI's STL implementation (including iostreams, etc.) that's been ported to more than 20 compilers. Like SGI's library, STLport's is available for free download. If you're writing code that has to work on multiple platforms, you may be able to save yourself a wheelbarrow of grief by standardizing on the STLport implementation and using it with all your compilers.

Most of STLport's modifications to SGI's code base focus on improved portability, but STLport's STL is also the only implementation I know that offers a "debug mode" to help detect improper use of the STL — uses that compile but lead to undefined runtime behavior. For example, [Item 30](#) uses this example in its discussion of the common mistake of writing beyond the end of a container:

```
int transmogrify(int x);                // this function produces
                                       // some new value from x

vector<int> values;

...                                   // put data into values

vector<int> results;
transform(values.begin(), values.end(),    // this will attempt to
          results.end(),                  // write beyond the
          transmogrify);                  // end of results!
```

This will compile, but when run, it yields undefined results. If you're lucky, something horrible will happen inside the call to `transform`, and debugging the problem will be relatively straightforward. If you're not lucky, the call to `transform` will trash data somewhere in your address space, but you won't discover that until later. At that point, determining the cause of the memory corruption will be — shall we say? — *challenging*.

STLport's debug mode all but eliminates the challenge. When the above call to `transform` is executed, the following message is generated (assuming STLport is installed in the directory `C:\STLport`):

```
C:\STLport\stlport\stl\debug\_iterator.h:265 STL assertion failure :
_Dereferenceable(*this)
```

The program then stops, because STLport debug mode calls abort if it encounters a usage error. If you'd prefer to have an exception thrown instead, you can configure STLport to do things your way.

Admittedly, the above error message isn't as clear as it might be, and it's unfortunate that the reported file and line correspond to the location of the internal STL assertion instead of the line calling transform, but this is still a lot better than running past the call to transform, then trying to figure out why your data structures are corrupt. With STLport's debug mode, all you need to do is fire up your debugger and walk the call stack back into the code you wrote, then determine what you did wrong. Finding the offending source line is generally not a problem.

STLport's debug mode detects a variety of common errors, including passing invalid ranges to algorithms, attempting to read from an empty container, using an iterator from one container as the argument to a second container's member function, etc. It accomplishes this magic by having iterators and their containers track one another. Given two iterators, it's thus possible to check to see if they come from the same container, and when a container is modified, it's possible to invalidate the appropriate set of iterators.

Because STLport uses special iterator implementations in debug mode, iterators for vector and string are class objects instead of raw pointers. Hence, using STLport and compiling in debug mode is a good way to make sure that nobody is getting sloppy about the difference between pointers and iterators for these container types. That alone may be reason enough to give STLport's debug mode a try.

The Boost Web Site

In 1997, when the closing bell rang on the process that led to the International Standard for C++, some people were disappointed that library features they'd advocated hadn't made the cut. Some of these people were members of the Committee itself, so they set out to lay the foundation for additions to the standard library during the second round of standardization. The result is Boost, a web site whose mission is to "provide free, peer-reviewed, C++ libraries. The emphasis is on portable libraries which work well with the C++ Standard Library." Behind the mission is a motive:

To the extent a library becomes "existing practice", the likelihood increases that someone will propose it for future standardization. Submitting a library to Boost.org is one way to establish existing practice... .

In other words, Boost offers itself as a vetting mechanism to help separate the sheep from the goats when it comes to potential additions to the standard C++ library. This is a worthy service, and we should all be grateful.

Another reason to be grateful is the collection of libraries you'll find at Boost. I won't attempt to describe them all here, not least because new ones will doubtless have been added by the time you read these words. For STL users, however, two kinds of libraries are particularly relevant. The first is the smart pointer library featuring `shared_ptr`, the template for reference-counted smart pointers that, unlike the standard library's `auto_ptr`, may safely be stored in STL containers (see [Item 8](#)). Boost's smart pointer library also offers `shared_array`, a reference-counted smart pointer for dynamically allocated arrays, but [Item 13](#) argues that dynamically allocated arrays are inferior to vectors and strings, and I hope you find its argument persuasive.

Boost's second attraction for STL fans is its bevy of STL-related function objects and associated facilities. These libraries comprise a fundamental redesign and reimplementations of the ideas behind STL function objects and adapters, and the results eliminate a host of restrictions that artificially limit the utility of the standard functors. As an example of such a restriction, you'll find that if you try to use `bind2nd` with `mem_fun` or `mem_fun_ref` (see [Item 41](#)) to bind an object to a member function's parameter and that member function takes its parameter by reference, your code is unlikely to compile. You'll find the same if you try to use `not1` or `not2` with `ptr_fun` and a function declaring a by-reference parameter. The cause in both cases is that during template instantiation, most STL platforms generate a reference to a reference, and references to references are not legal in C++. (The Standardization Committee is mulling over a change in the Standard to address this matter.) Here's an example of what has become known as "the reference-to-reference problem:"

```
class Widget {  
public:  
    ...  
    int readStream(istream& stream);           // readStream takes  
    ...                                       // its parameter by  
};                                           // reference  
vector<Widget*> vw;  
...
```

```
for_each(                                     // most STL platforms
    vw.begin(), vw.end(),                   // try to generate a
    bind2nd(mem_fun(&Widget::readStream), cin) // reference to a
);                                           // reference in this
                                           // call; such code
                                           // won't compile
```

Boost's function objects avoid this and other issues, plus they expand the expressiveness of function objects considerably.

If you're intrigued by the potential of STL function objects and you want to explore it further, hurry over to Boost right away. If you abhor function objects and think they exist only to pacify a vocal minority of Lisp apologists turned C++ programmers, hurry over to Boost anyway. Boost's function object libraries are important, but they make up only a small part of what you'll find at the site.

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Bibliography

Most of the publications below are cited in this book, though many citations occur only in the Acknowledgments. Uncited publications are bulleted instead of having a number.

Given the instability of Internet URLs, I hesitated before including them in this bibliography. In the end, I decided that even if a URL is broken by the time you try it, knowing where a document used to reside should help you find it at a different URL.

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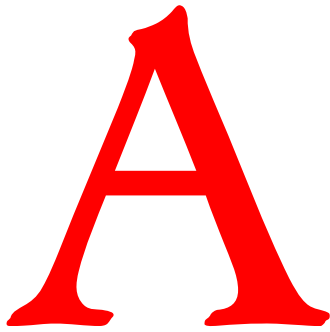
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- [5] ISO/IEC, *International Standard, Programming Languages — C++* (second edition), Reference Number ISO/IEC 14882:2003(E), 2003. The official document describing C++. Available in PDF from ANSI at <http://webstore.ansi.org/ansidocstore/default.asp>.
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- [7] Bjarne Stroustrup, *The C++ Programming Language* (third edition), Addison-Wesley, 1997, ISBN 0-201-88954-4. The “resource acquisition is initialization” idiom I mention in [Item 12](#) is discussed in section 14.4.1 of this book, and the code I refer to in [Item 36](#) is on page 530.
- [8] Herb Sutter, *Exceptional C++: 47 Engineering Puzzles, Programming Problems, and Solutions*, Addison-Wesley, 2000, ISBN 0-201-61562-2. An exemplary complement to my *Effective* books, I’d laud this even if Herb hadn’t asked me to write the Foreword for it.
- [9] Herb Sutter, *More Exceptional C++: 40 New Engineering Puzzles, Programming Problems, and Solutions*, Addison-Wesley, 2001, ISBN 0-201-70434-X. This book is every bit as good as its predecessor [8].
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- [11] Matt Austern, “How to Do Case-Insensitive String Comparison,” *C++ Report*, May 2000. This article is so useful, it’s reproduced as [Appendix A](#) of this book.

- [12] Herb Sutter, "When Is a Container Not a Container?," *C++ Report*, May 1999. Available at <http://www.gotw.ca/publications/mill09.htm>. Revised and updated as Item 6 in *More Exceptional C++* [9].
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- [21] The SGI STL web site, <http://www.sgi.com/tech/stl/>. [Item 50](#) summarizes the material at this important site. The page on thread safety in STL containers (which motivated [Item 12](#)) is at http://www.sgi.com/tech/stl/thread_safety.html.
- [22] The Boost web site, <http://www.boost.org/>. [Item 50](#) summarizes the material at this important site.
- [23] Nicolai M. Josuttis, "User-Defined Allocator," <http://www.josuttis.com/cppcode/allocator.html>. This page is part of the web site for Josuttis' excellent book on C++'s standard library [3].

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 - Matt Austern, "The Standard Librarian: Bitsets and Bit Vectors," *C/C++ Users Journal's C++ Experts Forum* (an on-line extension to the magazine), May 2001, <http://www.ddj.com/cpp/184401382>. This article provides information on bitsets and how they compare to `vector<bool>`, topics I briefly examine in Item 18.

Things I Had To Write (But Wish I Hadn't)

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- The *Effective C++ CD* Errata List, <http://www.aristeia.com/BookErrata/cd1e-errata.html>.
- [29] The *More Effective C++* `auto_ptr` Update page, http://www.aristeia.com/BookErrata/auto_ptr-update.html.



Locales and Case-Insensitive String Comparisons

[Item 35](#) explains how to use `mismatch` and `lexicographical_compare` to implement case-insensitive string comparisons, but it also points out that a truly general solution to the problem must take locales into account. This book is about the STL, not internationalization, so I have next to nothing to say about locales. However, Matt Austern, author of *Generic Programming and the STL* [\[4\]](#), addressed the locale aspects of case-insensitive string comparison in a column in the May 2000 *C++ Report* [\[11\]](#). In the interest of telling the whole story on this important topic, I am pleased to reprint his column here, and I am grateful to Matt and to 101communications for granting me permission to do so.

How to Do Case-Insensitive String Comparison

by Matt Austern

If you've ever written a program that uses strings (and who hasn't?), chances are you've sometimes needed to treat two strings that differ only by case as if they were identical. That is, you've needed comparisons — equality, less than, substring matches, sorting — to disregard case. And, indeed, one of the most frequently asked questions about the Standard C++ Library is how to make strings case insensitive. This question has been answered many times. Many of the answers are wrong.

First of all, let's dispose of the idea that you should be looking for a way to write a case-insensitive string class. Yes, it's technically possible, more or less. The Standard Library type `std::string` is really just an alias for the template `std::basic_string<char, std::char_traits<char>, std::allocator<char>>`. It uses the traits parameter for all comparisons, so, by providing a traits parameter with equality and less than redefined appropriately, you can instantiate `basic_string` in such a way so that the `<` and `==` operators are case insensitive. You can do it, but it isn't worth the trouble.

- *You won't be able to do I/O*, at least not without a lot of pain. The I/O classes in the Standard Library, like `std::basic_istream` and `std::basic_ostream`, are templated on character type and traits just like `std::basic_string` is. (Again, `std::ostream` is just an alias for `std::basic_ostream<char, char_traits<char> >`.) The traits parameters have to match. If you're using `std::basic_string<char, my_traits_class>` for your strings, you'll have to use `std::basic_ostream<char, my_traits_class>` for your stream output. You won't be able to use ordinary stream objects like `cin` and `cout`.
- *Case insensitivity isn't about an object, it's about how you use an object*. You might very well need to treat a string as case-sensitive in some contexts and case-insensitive in others. (Perhaps depending on a user-controlled option.) Defining separate types for those two uses puts an artificial barrier between them.
- *It doesn't quite fit*. Like all traits classes,[†] `char_traits` is small, simple, and stateless. As we'll see later in this column, proper case-insensitive comparisons are none of those things.
- *It isn't enough*. Even if all of `basic_string`'s own member functions were case insensitive, that still wouldn't help when you need to use nonmember generic algorithms like `std::search` and `std::find_end`. It also wouldn't help if you decided, for reasons of efficiency, to change from a container of `basic_string` objects to a string table.

A better solution, one that fits more naturally into the design of the Standard Library, is to ask for case-insensitive comparison when that's what you need. Don't bother with member functions like `string::find_first_of` and `string::rfind`; all of them duplicate functionality that's already there in nonmember generic algorithms. The generic algorithms, meanwhile, are flexible enough to accommodate case-insensitive strings. If you need to sort a collection of strings in case-insensitive order, for example, all you have to do is provide the appropriate comparison function object:

```
std::sort(C.begin(), C.end(), compare_without_case);
```

The remainder of this column will be about how to write that function object.

A first attempt

There's more than one way to alphabetize words. The next time you're in a bookstore, check how the authors' names are arranged: does

[†] See Andrei Alexandrescu's column in the April 2000 *C++ Report* [19].

Mary McCarthy come before Bernard Malamud, or after? (It's a matter of convention, and I've seen it done both ways.) The simplest kind of string comparison, though, is the one we all learned in elementary school: lexicographic or "dictionary order" comparison, where we build up string comparison from character-by-character comparison.

Lexicographic comparison may not be suitable for specialized applications (no single method is; a library might well sort personal names and place names differently), but it's suitable much of the time and it's what string comparison means in C++ by default. Strings are sequences of characters, and if `x` and `y` are of type `std::string`, the expression `x < y` is equivalent to the expression

```
std::lexicographical_compare(x.begin(), x.end(), y.begin(), y.end()).
```

In this expression, `lexicographical_compare` compares individual characters using operator<, but there's also a version of `lexicographical_compare` that lets you choose your own method of comparing characters. That other version takes five arguments, not four; the last argument is a function object, a Binary Predicate that determines which of two characters should precede the other. All we need in order to use `lexicographical_compare` for case-insensitive string comparison, then, is to combine it with a function object that compares characters without regard to case.

The general idea behind case-insensitive comparison of characters is to convert both characters to upper-case and compare the results. Here's the obvious translation of that idea into a C++ function object, using a well known function from the standard C library:

```
struct lt_nocase
: public std::binary_function<char, char, bool> {
    bool operator()(char x, char y) const {
        return std::toupper(static_cast<unsigned char>(x)) <
               std::toupper(static_cast<unsigned char>(y));
    }
};
```

"For every complex problem there is a solution that is simple, neat, and wrong." People who write books about C++ are fond of this class, because it makes a nice, simple example. I'm as guilty as anyone else; I use it in my book a half dozen times. It's almost right, but that's not good enough. The problem is subtle.

Here's one example where you can begin to see the problem:

```

int main()
{
    const char* s1 = "GEW\334RZTRAMINER";
    const char* s2 = "gew\374rztraminer";
    printf("s1 = %s, s2 = %s\n", s1, s2);

    printf("s1 < s2: %s\n",
           std::lexicographical_compare(s1, s1 + 14, s2, s2 + 14, lt_nocase())
           ? "true" : "false");
}

```

You should try this out on your system. On my system (a Silicon Graphics O2 running IRIX 6.5), here's the output:

```

s1 = GEWÜRZTRAMINER, s2 = gewürztraminer
s1 < s2: true

```

Hm, how odd. If you're doing a case-insensitive comparison, shouldn't "gewürztraminer" and "GEWÜRZTRAMINER" be the same? And now a slight variation: if you insert the line

```

setlocale(LC_ALL, "de");

```

just before the printf statement, suddenly the output changes:

```

s1 = GEWÜRZTRAMINER, s2 = gewürztraminer
s1 < s2: false

```

Case-insensitive string comparison is more complicated than it looks. This seemingly innocent program depends crucially on something most of us would prefer to ignore: locales.

Locales

A char is really nothing more than a small integer. We can choose to interpret a small integer as a character, but there's nothing universal about that interpretation. Should some particular number be interpreted as a letter, a punctuation mark, or a nonprinting control character?

There's no single right answer, and it doesn't even make a difference as far as the core C and C++ languages are concerned. A few library functions need to make those distinctions: `isalpha`, for example, which determines whether a character is a letter, and `toupper`, which converts lowercase letters to uppercase and does nothing to uppercase letters or to characters that aren't letters. All of that depends on local cultural and linguistic conventions: distinguishing between letters and nonletters means one thing in English, another thing in Swedish. Conversion from lower to upper case means something different in the Roman alphabet than in the Cyrillic, and means nothing at all in Hebrew.

By default the character manipulation functions work with a character set that's appropriate for simple English text. The character '\374' isn't affected by `toupper` because it isn't a letter; it may look like `ü` when it's printed on some systems, but that's irrelevant for a C library routine that's operating on English text. There is no `ü` character in the ASCII character set. The line

```
setlocale(LC_ALL, "de");
```

tells the C library to start operating in accordance with German conventions. (At least it does on IRIX. Locale names are not standardized.) There is a character `ü` in German, so `toupper` changes `ü` to `Ü`.

If this doesn't make you nervous, it should. While `toupper` may look like a simple function that takes one argument, it also depends on a global variable—worse, a hidden global variable. This causes all of the usual difficulties: a function that uses `toupper` potentially depends on every other function in the entire program.

This can be disastrous if you're using `toupper` for case-insensitive string comparison. What happens if you've got an algorithm that depends on a list being sorted (`binary_search`, say), and then a new locale causes the sort order to change out from under it? Code like this isn't reusable; it's barely usable. You can't use it in libraries—libraries get used in all sorts of programs, not just programs that never call `setlocale`. You might be able to get away with using it in a large program, but you'll have a maintenance problem: maybe you can prove that no other module ever calls `setlocale`, but can you prove that no other module in next year's version of the program will call `setlocale`?

There's no good solution to this problem in C. The C library has a single global locale, and that's that. There is a solution in C++.

Locales in C++

A locale in the C++ Standard Library isn't global data buried deep within the library implementation. It's an object of type `std::locale`, and you can create it and pass it to functions just like any other object. You can create a locale object that represents the usual locale, for example, by writing

```
std::locale L = std::locale::classic();
```

or you can create a German locale by writing

```
std::locale L("de");
```

(As in the C library, names of locales aren't standardized. Check your implementation's documentation to find out what named locales are available.)

Locales in C++ are divided into *facets*, each of which handles a different aspect of internationalization, and the function `std::use_facet` extracts a specific facet from a locale object.[†] The `ctype` facet handles character classification, including case conversion. Finally, then, if `c1` and `c2` are of type `char`, this fragment will compare them in a case-insensitive manner that's appropriate to the locale `L`.

```
const std::ctype<char>& ct = std::use_facet<std::ctype<char>>(L);
bool result = ct.toupper(c1) < ct.toupper(c2);
```

There's also a special abbreviation: you can write

```
std::toupper(c, L),
```

which (if `c` is of type `char`) means the same thing as

```
std::use_facet<std::ctype<char>>(L).toupper(c).
```

It's worth minimizing the number of times you call `use_facet`, though, because it can be fairly expensive.

Just as lexicographical comparison isn't appropriate for all applications, so character-by-character case conversion isn't always appropriate. (In German, for example, the lowercase letter “ß” corresponds to the uppercase sequence “SS”.) Unfortunately, however, character-by-character case conversion is all we've got. Neither the C nor the C++ standard library provides any form of case conversion that works with anything more than one character at a time. If this restriction is unacceptable for your purposes, then you're outside the scope of the standard library.

A digression: another facet

If you're already familiar with locales in C++, you may have thought of another way to perform string comparisons: The `collate` facet exists precisely to encapsulate details of sorting, and it has a member function with an interface much like that of the C library function `strcmp`. There's even a little convenience feature: if `L` is a locale object, you can compare two strings by writing `L(x, y)` instead of going through the nuisance of calling `use_facet` and then invoking a `collate` member function.

The “classic” locale has a `collate` facet that performs lexicographical comparison, just like `string`'s `operator<` does, but other locales perform

[†] Warning: `use_facet` is a function template whose template parameter appears only in the return type, not in any of the arguments. Calling it uses a language feature called *explicit template argument specification*, and some C++ compilers don't support that feature yet. If you're using a compiler that doesn't support it, your library implementor may have provided a workaround so that you can call `use_facet` some other way.

whatever kind of comparison is appropriate. If your system happens to come with a locale that performs case-insensitive comparisons for whatever languages you're interested in, you can just use it. That locale might even do something more intelligent than character-by-character comparisons!

Unfortunately, this piece of advice, true as it may be, isn't much help for those of us who don't have such systems. Perhaps someday a set of such locales may be standardized, but right now they aren't. If someone hasn't already written a case-insensitive comparison function for you, you'll have to write it yourself.

Case-insensitive string comparison

Using `ctype`, it's straightforward to build case-insensitive string comparison out of case-insensitive character comparison. This version isn't optimal, but at least it's correct. It uses essentially the same technique as before: compare two strings using `lexicographical_compare`, and compare two characters by converting both of them to uppercase. This time, though, we're being careful to use a locale object instead of a global variable. (As an aside, converting both characters to uppercase might not always give the same results as converting both characters to lowercase: There's no guarantee that the operations are inverses. In French, for example, it's customary to omit accent marks on uppercase characters. In a French locale it might be reasonable for `toupper` to be a lossy conversion; it might convert both 'é' and 'e' to the same uppercase character, 'E'. In such a locale, then, a case-insensitive comparison using `toupper` will say that 'é' and 'E' are equivalent characters while a case-insensitive comparison using `tolower` will say that they aren't. Which is the right answer? Probably the former, but it depends on the language, on local customs, and on your application.)

```
struct lt_str_1
: public std::binary_function<std::string, std::string, bool> {
    struct lt_char {
        const std::ctype<char>& ct;

        lt_char(const std::ctype<char>& c) : ct(c) {}

        bool operator()(char x, char y) const {
            return ct.toupper(x) < ct.toupper(y);
        }
    };

    std::locale loc;
    const std::ctype<char>& ct;
```

```

lt_str_1(const std::locale& L = std::locale::classic())
    : loc(L), ct(std::use_facet<std::ctype<char>>(loc)) {}

bool operator()(const std::string& x, const std::string& y) const {
    return std::lexicographical_compare(x.begin(), x.end(),
                                         y.begin(), y.end(),
                                         lt_char(ct));
}

};

```

This isn't quite optimal yet; it's slower than it ought to be. The problem is annoying and technical: we're calling `toupper` in a loop, and the C++ standard requires `toupper` to make a virtual function call. Some optimizers may be smart enough to move the virtual function overhead out of the loop, but most aren't. Virtual function calls in loops should be avoided.

In this case, avoiding it isn't completely straightforward. It's tempting to think that the right answer is another one of `ctype`'s member functions,

```
const char* ctype<char>::toupper(char* f, char* l) const,
```

which changes the case of the characters in the range `[f, l)`. Unfortunately, this isn't quite the right interface for our purpose. Using it to compare two strings would require copying both strings to buffers and then converting the buffers to uppercase. Where do those buffers come from? They can't be fixed size arrays (how large is large enough?), but dynamic arrays would require an expensive memory allocation.

An alternative solution is to do the case conversion once for every character, and cache the result. This isn't a fully general solution—it would be completely unworkable, for example, if you were working with 32-bit UCS-4 characters. If you're working with `char`, though (8 bits on most systems), it's not so unreasonable to maintain 256 bytes of case conversion information in the comparison function object.

```

struct lt_str_2 :
    public std::binary_function<std::string, std::string, bool> {
    struct lt_char {
        const char* tab;

        lt_char(const char* t) : tab(t) {}

        bool operator()(char x, char y) const {
            return tab[x - CHAR_MIN] < tab[y - CHAR_MIN];
        }
    };

    char tab[CHAR_MAX - CHAR_MIN + 1];

```

```

lt_str_2(const std::locale& L = std::locale::classic()) {
    const std::ctype<char>& ct = std::use_facet<std::ctype<char>>(L);
    for (int i = CHAR_MIN; i <= CHAR_MAX; ++i)
        tab[i - CHAR_MIN] = (char) i;
    ct.toupper(tab, tab + (CHAR_MAX - CHAR_MIN + 1));
}

bool operator()(const std::string& x, const std::string& y) const {
    return std::lexicographical_compare(x.begin(), x.end(),
                                         y.begin(), y.end(),
                                         lt_char(tab));
}
};

```

As you can see, `lt_str_1` and `lt_str_2` aren't very different. The former has a character-comparison function object that uses a `ctype` facet directly, and the latter a character-comparison function object that uses a table of precomputed uppercase conversions. This might be slower if you're creating an `lt_str_2` function object, using it to compare a few short strings, and then throwing it away. For any substantial use, though, `lt_str_2` will be noticeably faster than `lt_str_1`. On my system the difference was more than a factor of two: it took 0.86 seconds to sort a list of 23,791 words with `lt_str_1`, and 0.4 with `lt_str_2`.

What have we learned from all of this?

- A case-insensitive string class is the wrong level of abstraction. Generic algorithms in the C++ standard library are parameterized by policy, and you should exploit that fact.
- Lexicographical string comparisons are built out of character comparisons. Once you've got a case-insensitive character comparison function object, the problem is solved. (And you can reuse that function object for comparisons of other kinds of character sequences, like `vector<char>`, or string tables, or ordinary C strings.)
- Case-insensitive character comparison is harder than it looks. It's meaningless except in the context of a particular locale, so a character comparison function object needs to store locale information. If speed is a concern, you should write the function object to avoid repeated calls of expensive facet operations.

Correct case-insensitive comparison takes a fair amount of machinery, but you only have to write it once. You probably don't want to think about locales; most people don't. (Who wanted to think about Y2K bugs in 1990?) You stand a better chance of being able to ignore locales if you get locale-dependent code right, though, than if you write code that glosses over the dependencies.

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Remarks on Microsoft's STL Platforms

The information that follows is for developers using Microsoft Visual C++ (MSVC) versions 4–6. If you’re using Visual C++ .NET, your STL platform doesn’t have the problems described below, and you may ignore this appendix.

Suppose you have two vectors of Widgets and you'd like to copy the Widgets in one vector to the end of another. That's easy. Just use vector's range insert function (see [Item 5](#)):

If you have a vector and a deque, you do the same thing:

[illegible]

In fact, you can do this no matter what the type of the container holding the objects to be copied. Even custom containers work:

```
vector<Widget> vw;
...
list<Widget> lw;
...
vw.insert(vw.begin(), lw.begin(), lw.end());           // prepend to vw a copy
                                                         // of the Widgets in lw

set<Widget> sw;
...
vw.insert(vw.begin(), sw.begin(), sw.end());           // prepend to vw a copy
                                                         // of the Widgets in sw

template<typename T,                                   // template for a custom
        typename Allocator = allocator<T> >           // STL-compatible
class SpecialContainer { ... };                       // container
SpecialContainer<Widget> scw;
...
vw.insert(vw.end(), scw.begin(), scw.end());           // append to vw a copy
                                                         // of the Widgets in scw
```

This flexibility is possible because vector's range insert function isn't a function at all. Rather, it's a *member function template* that can be instantiated with *any iterator type* to generate a specific range insert function. For vector, the Standard declares the insert template like this:

```
template <class T, class Allocator = allocator<T> >
class vector {
public:
    ...
    template <class InputIterator>
    void insert(iterator position, InputIterator first, InputIterator last);
    ...
};
```

Each standard container is required to offer this templated range insert. Similar member function templates are required for the containers' range constructors and for the range form of assign (both of which are discussed in [Item 5](#)).

MSVC Versions 4–6

Unfortunately, the STL implementation that ships with MSVC versions 4–6 declares no member function templates. This library was

originally developed for MSVC version 4, and that compiler, like most compilers of its day, lacked member function template capabilities. Between MSVC4 and MSVC6, the compiler added support for these templates, but, due to legal proceedings that affected Microsoft without directly involving them, the library remained essentially frozen.

Because the STL implementation shipping with MSVC4–6 was designed for a compiler lacking member function templates, the library's authors approximated the functionality of such templates by replacing each template with a specific function, one that accepted only iterators from the same container type. For `insert`, for example, the member function template was replaced with this:

```
void insert(iterator position,           // "iterator" is the
           iterator first, iterator last); // container's iterator type
```

This restricted form of range member functions makes it possible to perform a range `insert` from a `vector<Widget>` to a `vector<Widget>` or from a `list<int>` to a `list<int>`, but not from a `vector<Widget>` to a `list<Widget>` or from a `set<int>` to a `deque<int>`. It's not even possible to do a range `insert` (or `assign` or `construction`) from a `vector<long>` to a `vector<int>`, because `vector<long>::iterator` is not the same type as a `vector<int>::iterator`. As a result, the following perfectly valid code fails to compile using MSVC4–6:

```
istream_iterator<Widget> begin(cin), end; // create begin and end
                                           // iterators for reading
                                           // Widgets from cin
                                           // (see Item 6)

vector<Widget> vw(begin, end);             // read cin's Widgets into vw
                                           // (again, see Item 6); won't
                                           // compile with MSVC4–6

list<Widget> lw;
...
lw.assign(vw.rbegin(), vw.rend());         // assign vw's contents to lw
                                           // (in reverse order); won't
                                           // compile with MSVC4–6

SpecialContainer<Widget> scw;
...
scw.insert(scw.end(), lw.begin(), lw.end()); // insert at the end of scw a
                                           // copy of the Widgets in lw;
                                           // won't compile with
```

MSVC4–6

So what do you do if you must use MSVC4–6? That depends on the MSVC version you are using and whether you are forced to use the STL implementation that comes with the compiler.

A Workaround for MSVC4-5

Look again at the valid code examples that fail to compile with the STL that accompanies MSVC4-6:

```
vector<Widget> vw(begin, end);           // rejected by the MSVC4-6
                                         // STL implementation

list<Widget> lw;
...
lw.assign(vw.rbegin(), vw.rend());       //also rejected
SpecialContainer<Widget> scw;
...
scw.insert(scw.end(), lw.begin(), lw.end()); // ditto
```

These calls look rather different, but they all fail for the same reason: missing member function templates in the STL implementation. There's a single cure for all of them: use copy and an insert iterator (see [Item 30](#)) instead. Here, for example, are the workarounds for the examples above:

```
istream_iterator<Widget> begin(cin), end;

vector<Widget> vw;           // default-construct vw;
copy(begin, end, back_inserter(vw)); // then copy the
                                         // Widgets in cin into it

list<Widget> lw;
...
lw.clear();                 // eliminate lw's old
copy(vw.rbegin(), vw.rend(), back_inserter(lw)); // Widgets; copy over
                                         // vw's Widgets (in
                                         // reverse order)

SpecialContainer<Widget> scw;
...
copy(lw.begin(), lw.end(), // copy lw's Widgets to
    inserter(scw, scw.end())); // the end of scw
```

I encourage you to use these copy-based workarounds with the library that ships with MSVC4-5, but beware! Don't fall into the trap of becoming so comfortable with the workarounds, you forget that they are *workarounds*. As [Item 5](#) explains, using the copy algorithm is almost always inferior to using a range member function, so as soon as you have a chance to upgrade your STL platform to one that supports member function templates, stop using copy in places where range member functions are the proper approach.

An Additional Workaround for MSVC6

You can use the MSVC4–5 workaround with MSVC6, too, but for MSVC6 there is another option. The compilers that are part of MSVC4–5 offer no meaningful support for member function templates, so the fact that the STL implementation lacks them is immaterial. The situation is different with MSVC6, because MSVC6's compiler does support member function templates. It thus becomes reasonable to consider replacing the STL that ships with MSVC6 with an implementation that provides the member function templates the Standard prescribes.

[Item 50](#) explains that both SGI and STLport offer freely downloadable STL implementations, and both of those implementations count the MSVC6 compiler as one with which they'll work. You can also purchase the latest MSVC-compatible STL implementation from Dinkumware. Each option has advantages and disadvantages.

SGI's and STLport's implementations are free, and I suspect you know what that means as regards official support for the software: there isn't any. Furthermore, because SGI and STLport design their libraries to work with a variety of compilers, you will probably have to manually configure their implementations to get the most out of MSVC6. In particular, you may have to explicitly enable support for member function templates, because, working with many compilers as they do, SGI and/or STLport may not enable that by default. You may also have to worry about linking with other MSVC6 libraries (especially DLLs), including things like making sure you use the appropriate builds for threading and debugging, etc.

If that kind of thing scares you, or if you've been known to grumble that you can't afford free software, you may want to look into Dinkumware's replacement library for MSVC6. It's designed to be drop-in compatible with the native MSVC6 STL and to maximize MSVC6's adherence to the Standard as an STL platform. Since Dinkumware authored the STL that ships with MSVC6, there's a decent chance their latest STL implementation really *is* a drop-in replacement. To learn more about Dinkumware STL implementations, visit the company's web site: <http://www.dinkumware.com/>.

Regardless of whether you choose SGI's, STLport's, or Dinkumware's implementation as an STL replacement, you'll do more than gain an STL with member function templates. You'll also bypass conformance problems in other areas of the library, such as string failing to declare `push_back`. Furthermore, you'll gain access to useful STL extensions, including hashed containers (see [Item 25](#)) and singly linked lists (slists).

SGL's and STLport's implementations also offer a variety of nonstandard functor classes, such as `select1st` and `select2nd` (see [Item 50](#)).

Even if you're trapped with the STL implementation that ships with MSVC6, it's probably worth your while to visit the Dinkumware web site. That site lists known bugs in the MSVC6 library implementation and explains how to modify your copy of the library to reduce its shortcomings. Needless to say, editing your library headers is something you do at your own risk. If you run into trouble, don't blame me.

Index

The example classes and class templates declared or defined in this book are indexed under *example classes/templates*. The example functions and function templates are indexed under *example functions/templates*.

Before A

`__default_alloc_template` 211

A

Abrahams, David xvii

abstraction bonus 203

abstraction penalty 201

accumulate

function objects for 158

initial value and 157, 159

side effects and 160

adaptability

algorithm function objects and 156

definition of 170

functor classes and 169–173

overloading operator() and 173

add or update functionality, in map 107

adjacent_difference 157

Adobe, demonic software spewed by xviii

advance

efficiency of 122

to create iterators from

const_iterators 120–123

Alexandrescu, Andrei xvii, 227, 230

algorithms

accumulate 156–161

function objects for 158

initial value and 157, 159

side effects and 160

adaptable function objects and 156

adjacent_difference 157

as a vocabulary 186

binary_search 192–201

container mem funcs vs. 190–192

copy, eliminating calls to 26

copy_if, implementing 154–156

copying func objects within 166–168

count 156, 192–201

equal_range vs., for sorted ranges 197

count_if 157

efficiency, vs. explicit loops 182–184

equal_range 192–201

count vs., for sorted ranges 197

find 192–201

count in multiset, multimap vs. 199

lower_bound in multiset, multimap
vs. 201

using equivalence vs. equality 86

for_each 156–161

side effects and 160

function call syntax in 175

function parameters to 201–205

hand-written loops vs. 181–189

includes 148

inner_product 157

inplace_merge 148

lexicographical_compare 153

longest name 153

loops and 182

lower_bound 192–201

equality vs. equivalence and 195

max_element 157

merge 148, 192

min_element 157

mismatch 151

nth_element 133–138

optimizations in 183

- partial_sort 133–138
- partial_sum 157
- partition 133–138
 - containers of pointers and 145
- remove 139–143
 - on containers of pointers 143–146
- remove_copy_if 44
- remove_if 44, 144, 145
 - see also [remove](#)
 - possible implementation of 167
- set_difference 148
- set_intersection 148
- set_symmetric_difference 148, 153
- set_union 148
- sort 133–138
- sorted ranges and 146–150
- sorting 133–138
 - alternatives to 138
- stable_partition 133–138
- stable_sort 133–138
- string member functions vs. 230
- transform 129
- unique 145, 148, 149
 - see also [remove](#)
- unique_copy 148, 149
- upper_bound 197–198

[All your base are belong to us](#)

allocations

- minimizing via reserve 66–68
- minimum, in string 72

allocator

- allocate interface, vs. operator new 51

allocators

- boilerplate code for 54
- conventions and restrictions 48–54
- fraud in interface 51
- in string 69
- legitimate uses 54–58
- never being called 52
- permitted assumptions about 49
- rebind template within 53
- stateful, portability and 50, 51, 212
- summary of guidelines for 54
- typedefs within 48
- URLs for examples 227, 228

allusions

- to *Candide* 143
- to *Do Not Go Gentle into that Good Night* 187
- to Martin Luther King, Jr. 51
- to Matthew 25:32 222
- to *The Little Engine that Could* 63

amortized constant time complexity 6

ANSI Standard for C++

- see [C++ Standard, The](#)

anteater 88

argument_type 170

array-based containers, definition of 13

arrays

- as part of the STL 2
- as STL containers 64
- containers vs. 22
- vector and string vs. 63–66

assignments

- assign vs. operator= 25
- superfluous, avoiding 31
- via range member functions 33

associative containers

- see [standard associative containers](#), [hashed containers](#)

Austern, Matt xvi, xvii, 54, 226, 227, 228

- see also [Generic Programming and the STL](#)

author, contacting the xii

auto_ptr

- as container element 40–43
- semantics of copying 41
- sorting 41
- URL for update page on 228

average, finding, for a range 159–161

Avery, Katrina xvii

B

back_insert_iterator 216

back_inserter 130, 216

- push_back and 130

backwards order, inserting elements

- in 184

Barron, Carl xvi

base, see [reverse_iterator](#)

basic_ostream, relation to ostream 230

basic_string

- relation to string 210, 229
- relation to string and wstring 4

Becker, Pete xvii

Becker, Thomas xvii

begin/end, relation to rbegin/rend 123

bidirectional iterators

- binary search algorithms and 148
- definition of 5
- standard associative containers and 5

binary_function 170–172

- pointer parameters and 172
- reference parameters and 171

binary_search 199

- bsearch vs. 194
- related functions vs. 192–201

bind1st 216

- adaptability and 170

bind2nd 210, 216

- adaptability and 170
- reference-to-reference problem and 222

binder1st 216

binder2nd 216
 binders, definition of 5
 bitfields 80
 bitset
 as alternative to `vector<bool>` 81
 as part of the STL 2
 Boost 170, 172, 221–223
 shared_array 222
 shared_ptr 39, 146, 165, 178, 222
 web site URL 217, 227
 Bridge Pattern 165
 Bruce, Robin xviii
 bsearch, vs. `binary_search` 194
 bugs
 Dinkumware list for MSVC 244
 in this book, reporting xii
 Bulka, Dov 226
 see also *Efficient C++*
 Buzard, Susannah xviii

C

C++ Programming Language, The 61, 155
 bibliography entry for 226
C++ Standard Library, The 2, 6, 94, 113, 207
 bibliographic entry for 225
 C++ Standard, The
 bibliography entry for 226
 citation number omitted for 3
 guidance on choosing containers 12
 reference counting and 64
 URL for purchasing 226
 Campbell, Aaron xvii
Candide, allusion to 143
 capacity
 cost of increasing for `vector/string` 66
 minimizing in `vector` and `string` 77–79
 capacity, vs. size 66–67
 Carey, Alicia xviii
 case conversions, when not one-for-one 234
 case-insensitive
 string class 229–230
 string comparisons, see *string*
 casts
 const_iterator to iterator 120
 creating temporary objects via 98
 to references 98
 to remove constness 98
 to remove `map/multimap` constness 99
 categories, for iterators 5
 char_traits 113, 211, 230
 choosing
 among `binary_search`, `lower_bound`,
 `upper_bound`, and
 `equal_range` 192–201
 among containers 12

 among iterator types 116–119
 vector vs. `string` 64
 citations, in this book 2
 class vs. `typename` 7
 classes, vs. structs for functor classes 171
 clustering, in node-based containers 103
 COAPs, see *containers*, *auto_ptr*
 collate facet 234
 color, use in this book 8
`comp.lang.c++.moderated` xv
`comp.std.c++` xv
 comparison functions
 consistency and 149
 equal values and 92–94
 for pairs 104
 comparisons
 see also *string*, *comparisons*, *comparison functions*
 iterators with `const_iterators` 118
 lexicographic 231
 compilers
 diagnostics, deciphering 210–217
 implementations, vs. STL impls 3
 independence, value of 3
 optimizations, inlining and 202
 problems, see *workarounds*
 complexity
 amortized constant time 6
 constant time 6
 guarantees, in the STL 5–6
 linear time 6
 logarithmic time 6
 compose1 219
 compose2 187, 219
 const_iterator
 casting to iterator 120
 comparisons with iterators 118
 converting to iterator 120–123
 other iterator types vs. 116–119
 const_reverse_iterator
 other iterator types vs. 116–119
 constant time complexity 6
 amortized, see *amortized constant time*
 constness
 casting away 98
 of `map/multimap` elements 95
 of `set/multiset` elements 95
 construction, via range mem funcs 31
 contacting the author xii
 container adapters, as part of the STL 2
 container-independent code 16, 47
 illusory nature of 15–20
 containers
 arrays as 64
 arrays vs. 22
 assign vs. `operator=` in 25
 associative, see *standard associative containers*

- auto_ptrs in 40–43
- calling empty vs. size 23–24
- choosing among
 - advice from the Standard 12
 - iterator types 116–119
- contiguous memory
 - see [contiguous-memory containers](#), [vector](#), [string](#), [deque](#), [rope](#)
- converting const_iterators to
 - iterators 120–123
- criteria for selecting 11–15
- deleting pointers in 36–40
- encapsulating 19
- erasing
 - see also [erase-remove idiom](#)
 - elements in 43–48
 - iterator invalidation during 14, 45
 - relation to remove 139–143
- exception safety and 14, 37, 39
- filling from legacy APIs 77
- [hashed](#), see [hashed containers](#)
- improving efficiency via reserve 66–68
- insertions
 - in reverse order 184
 - iterator invalidation during 14
- iterators
 - casting among 121
 - invalidation, see [iterators](#), [invalidation](#)
- mem funcs vs. algorithms 190–192
- node-based
 - see [node-based containers](#), [list](#), [standard associative containers](#), [slist](#)
- object copying and 20–22
- of pointers, remove and 143–146
- of proxy objects 82
- of smart pointers 39
- range vs. single-element member
 - functions 24–33
- relation of begin/end, rbegin/rend 123
- replacing one with another 18
- requirements in the Standard 79
- resize vs. reserve 67
- rolling back insertions and erasures 14
- rope 218
- sequence
 - see [standard sequence containers](#)
- size vs. capacity 66–67
- size_type typedef 158
- slist 218
- sorted vector vs. associative 100–106
- thread safety and 58–62
- transactional semantics for insertion
 - and erasing 14
- typedefs for 18–19
- value_type typedef 36, 108
- vector<bool>, problems with 79–82

- contiguous memory
 - for string 75
 - for vector 74
- contiguous-memory containers
 - see also [vector](#), [string](#), [deque](#), [rope](#)
 - cost of erasing in 32
 - cost of insertion 28
 - definition of 13
 - iterator invalidation in, see [iterators](#), [invalidation](#)
- conventions, for allocators 48–54
- conversions
 - among iterator types 117
 - from const_iterator to iterator 120–123
 - when not one-for-one 234
- copy
 - eliminating calls to 26
 - insert iterators and 26
 - missing member templates and 242
- copy_if, implementing 154–156
- copying
 - auto_ptrs, semantics of 41
 - function objects
 - efficiency and 164
 - within algorithms 166–168
 - objects in containers 20–22
- count 156
 - as existence test 193
 - equal_range vs., for sorted ranges 197
 - related functions vs. 192–201
- count_if 157
- <cctype>, conventions of functions in 151
- <ctype.h>, conventions of functions in 151

D

- Dalla Gasperina, Marco xvi
- [debug mode](#), see [STLport](#), [debug mode](#)
- deciphering compiler diagnostics 210–217
- definitions
 - adaptable function object 170
 - amortized constant time complexity 6
 - array-based container 13
 - bidirectional iterator 5
 - binder 5
 - COAP 40
 - constant time complexity 6
 - container-independent code 16
 - contiguous-memory container 13
 - equality 84
 - equivalence 84–85
 - forward iterator 5
 - function object 5
 - functor 5
 - functor class 5
 - input iterator 5
 - linear time complexity 6

- local class 189
- logarithmic time complexity 6
- monomorphic function object 164
- node-based container 13
- output iterator 5
- predicate 166
- predicate class 166
- pure function 166
- random access iterator 5
- range member function 25
- resource acquisition is initialization 61
- stability, in sorting 135
- standard associative container 5
- standard sequence container 5
- STL platform 3
- transactional semantics 14
- delete
 - delete[] vs. 63
 - using wrong form 64
- deleting
 - objects more than once 64
 - pointers in containers 36–40
- dependent types, typename and 7–8
- deque
 - unique invalidation rules for 15
 - see also [iterators, invalidation](#)
- deque<bool>, as alternative to
 - vector<bool> 81
- dereferencing functor class 90
- Derge, Gillmer [xvii](#)
- Design Patterns* 80, 165
 - bibliography entry for 226
- Design Patterns CD*
 - bibliography entry for 226
- destination range, ensuring adequate
 - size 129–133
- destructors, calling explicitly 56
- Dewhurst, Stephen [xvii](#)
- dictionary order comparison, see [lexicographic comparison](#)
- Dinkumware
 - bug list for MSVC STL 244
 - hashed containers implementation 114
 - interface for hashed containers 113
 - slist implementation 218
 - STL replacement for MSVC6 243
 - web site for 243
- disambiguating function and object
 - declarations 35
- distance
 - declaration for 122
 - efficiency of 122
 - explicit template argument specification and 122
 - to convert const iterators to
 - iterators 120–123
- Do Not Go Gentle into that Good Night*,
 - allusion to 187

- documentation, on-line, for the STL 217
- Dr. Seuss [xi](#)

E

- Effective C++*
 - bibliography entry for 225
 - citation number omitted for 3
 - inheritance from class without a virtual
 - destructor discussion in 37
 - inlining discussion in 202
 - object copying discussions in 21
 - pointer to implementation class discussion in 165
 - slicing problem discussion in 22
 - URL for errata list for 228
- Effective C++ CD*
 - bibliography entry for 225
 - URL for errata list for 228
- Effective STL*, web site for [xii](#)
- efficiency
 - see also [optimizations](#)
 - advance and 122
 - algorithms vs. explicit loops 182–184
 - associative container vs. sorted
 - vector 100–106
 - case-insensitive string comparisons
 - and 154
 - comparative, of sorting algorithms 138
 - copy vs. range insert 26
 - copying
 - function objects and 164
 - objects and 21
 - distance and 122
 - empty vs. size 23–24
 - erasing from contiguous-memory
 - containers 32
 - function objects vs. functions 201–205
 - hashed containers, overview of 111–115
 - hashing and 101
 - “hint” form of insert and 110
 - improving via custom allocators 54
 - increasing vector/string capacity 66
 - inlining and 202
 - inserting into contiguous-memory
 - containers 28
 - istreambuf_iterators and 126–127
 - Items on 9
 - list::remove vs. the erase-remove
 - idiom 43
 - logarithmic vs. linear 190
 - map::operator[] vs. map::insert 106–111
 - mem funcs vs. algorithms 190–192
 - minimizing reallocs via reserve 66–68
 - ostreambuf_iterators and 127
 - range vs. single-element member
 - functions 24–33
 - small string optimization 71
 - sort vs. qsort 203

- sorting algorithms, overview of 133–138
 - string implementation trade-offs 68–73
 - toupper and 236–237
 - use_facet and 234
- Efficient C++* 202
 - bibliography entry for 226
- Einstein, Albert 69
- Emacs 27
- email address
 - for comments on this book xii
 - for the President of the USA 212
- embedded nulls 75, 154
- empty, vs. size 23–24
- encapsulating containers 19
- equal_range 196–197
 - count vs., for sorted ranges 197
 - related functions vs. 192–201
- equal_to 86, 112
- equality
 - definition of 84
 - equivalence vs. 83–88
 - in hashed containers 113
 - lower_bound and 195
- equivalence
 - definition of 84–85
 - equality vs. 83–88
 - in hashed containers 113
 - lower_bound and 195
- equivalent values, inserting in order 198
- erase
 - see also [erase-remove idiom](#)
 - relation to remove algorithm 139–143
 - return types for 32
 - return value for standard sequence containers 46
- erase_after 218
- erase-remove idiom 43, 47, 142, 145, 146, 184, 207
 - limitations of 46
 - list::remove vs. 43
 - remove_if variant 144
- erasing
 - see also [containers, erasing](#)
 - base iterators and 124
 - elements in containers 43–48
 - rolling back 14
 - via range member functions 32
- errata list
 - for *Effective C++* 228
 - for *Effective C++ CD* 228
 - for *More Effective C++* 228
 - for this book xii
- error messages, deciphering 210–217
- example classes/templates
 - Average 205
 - BadPredicate 167, 168
 - BetweenValues 188, 189
 - BPF 164, 165
 - BPFImpl 165
 - CStringCompare 85
 - Contestant 77
 - CustomerList 20
 - DataCompare 105
 - DeleteObject 37, 38
 - Dereference 90
 - DereferenceLess 91
 - DoSomething 163
 - Employee 95
 - Heap1 57
 - Heap2 57
 - IDNumberLess 96
 - list 52
 - list::ListNode 52
 - Lock 60
 - lt_nocase 231
 - lt_str_1 235
 - lt_str_1::lt_char 235
 - lt_str_2 236
 - lt_str_2::lt_char 236
 - MaxSpeedCompare 179
 - MeetsThreshold 171
 - NiftyEmailProgram 212, 215
 - Person 198
 - PersonNameLess 198
 - Point 159, 161
 - PointAverage 160, 161
 - PtrWidgetNameCompare 172
 - RCSP 146
 - SharedMemoryAllocator 55
 - SpecialAllocator 19, 49
 - SpecialContainer 240
 - SpecialString 37
 - SpecialWidget 21
 - SpecificHeapAllocator 57
 - std::less<Widget> 178
 - StringPtrGreater 93
 - StringPtrLess 89
 - StringSize 204
 - Timestamp 197
 - vector 240
 - Widget 7, 18, 19, 21, 35, 84, 106, 111, 143, 174, 177, 182, 222
 - WidgetNameCompare 171
- example functions/templates
 - anotherBadPredicate 169
 - average 204
 - Average::operator() 205
 - BadPredicate::BadPredicate() 167
 - BadPredicate::operator() 167, 168
 - BetweenValues::BetweenValues 188
 - BetweenValues::operator() 188
 - BPF::operator() 164, 165
 - BPFImpl::BPFImpl 165
 - BPFImpl::operator() 165

ciCharCompare 151
 ciCharLess 153
 ciStringCompare 152, 153, 154
 CiStringCompare::operator() 85
 ciStringCompareImpl 152
 copy_if 155, 156
 DataCompare::keyLess 105
 DataCompare::operator() 105
 delAndNullifyUncertified 145
 DeleteObject::operator() 37, 38
 Dereference::operator() 90
 DereferenceLess::operator() 91
 doSomething 36, 37, 38, 39, 74, 75, 77
 DoSomething::operator() 163
 doubleGreater 202
 efficientAddOrUpdate 110
 Employee::idNumber 95
 Employee::name 95
 Employee::setName 95
 Employee::setTitle 95
 Employee::title 95
 fillArray 76, 77, 184
 fillString 76
 hasAcceptableQuality 137
 Heap1::alloc 57
 Heap1::dealloc 57
 IDNumberLess::operator() 96
 isDefective 155
 isInteresting 169
 lastGreaterThanOrEqualTo 8
 Lock::Lock 60
 Lock::Lock 60
 lt_nocase::operator() 231
 lt_str_1::lt_char::lt_char 235
 lt_str_1::lt_char::operator() 235
 lt_str_1::lt_str_1 236
 lt_str_1::operator() 236
 lt_str_2::lt_char::lt_char 236
 lt_str_2::lt_char::operator() 236
 lt_str_2::lt_str_2 237
 lt_str_2::operator() 237
 MaxSpeedCompare::operator() 179
 MeetsThreshold::MeetsThreshold 171
 MeetsThreshold::operator() 171
 NiftyEmailProgram::showEmailAddress
 212, 215
 operator< for Timestamp 197
 operator< for Widget 177
 operator== for Widget 8, 84
 Person::name 198
 Person::operator() 198
 Point::Point 159
 PointAverage::operator() 160, 161
 PointAverage::PointAverage 160, 161
 PointAverage::result 161
 print 90
 PtrWidgetNameCompare::operator() 172
 qualityCompare 134
 SharedMemoryAllocator::allocate 55

SharedMemoryAllocator::dealloc 55
 SpecificHeapAllocator::allocate 57
 SpecificHeapAllocator::dealloc 57
 std::less<Widget>::operator() 178
 stringLengthSum 158
 StringPtrGreater::operator() 93
 stringPtrLess 91
 StringPtrLess::operator() 89
 StringSize::operator() 204
 test 174
 transmogrify 129, 220
 vector<bool>::operator[] 80
 vector<bool>::reference 80
 Widget::isCertified 143
 Widget::maxSpeed 177
 Widget::operator= 21, 106, 111
 Widget::readStream 222
 Widget::redraw 182
 Widget::test 174
 Widget::weight 177
 Widget::Widget 21, 106
 widgetAPCompare 41
 WidgetNameCompare::operator() 171
 writeAverages 204, 205
 exception safety 14, 37, 39, 50, 61
Exceptional C++ 14, 165
 bibliography entry for 226
 exceptions, to guidelines in this book 10
 explicit template argument specification
 distance and 122
 for_each and 163
 use_facet and 234
 extending the STL 2

F

facets, locales and 234–235
 find
 count in multiset, multimap vs. 199
 lower_bound in multiset, multimap
 vs. 201
 related functions vs. 192–201
 using equivalence vs. equality 86
 first_argument_type 170
 for_each
 declaration for 163
 explicit template argument specifica-
 tion and 163
 possible implementation of 174
 side effects and 160
 forward iterators
 definition of 5
 operator-- and 5
 fragmentation, memory, reducing 54
 fraud, in allocator interface 51
 free STL implementations 217, 220
 front_insert_iterator 216

front_inserter 130, 216
 push_front and 130
 Fuller, John xviii
 function objects
 as workaround for compiler problems 204
 definition of 5
 dereferencing, generic 90
 for accumulate 158
 functions vs. 201–205
 monomorphic, definition of 164
 pass-by-value and 162–166
 slicing 164
 functional programming 206
 functions
 calling forms 173
 calling syntax in the STL 175
 comparison
 equal values and 92–94
 for pointers 88–91
 declaration forms 33–35
 declaring templates in 188
 function objects vs. 201–205
 in <cctype>, conventions of 151
 in <cctype.h>, conventions of 151
 pointers to, as formal parameters 34
 predicates, need to be pure 166–169
 pure, definition of 166
 range vs. single-element 24–33
 functor classes
 adaptability and 169–173
 classes vs. structs 171
 definition of 5
 overloading operator() in 114
 pass-by-value and 162–166
 functor, see [function objects](#)

G

Gamma, Erich 226
 see also [Design Patterns](#)
Generic Programming and the STL 2, 94, 217, 229
 bibliography entry for 226
 gewürztraminer 232
 glass, broken, crawling on 95, 97
 Glassborow, Francis xvii
 Green, Cliff xvii
 growth factor, for vector/string 66

H

hand-written loops
 algorithms vs. 181–189
 iterator invalidation and 185
 Hansen, Karin xviii

Harrison, Doug xvi, xvii
 hashed containers 111–115
 Dinkumware interface for 113
 equality vs. equivalence in 113
 SGI interface for 112
 two implementation approaches to 114
 headers
 #includeing the proper ones 209–210
 <algorithm> 210
 <cctype> 151
 <cctype.h> 151
 <functional> 210
 <iterator> 210
 <list> 209
 <map> 209
 <numeric> 210
 <numeric> 157
 <set> 209
 <vector> 209
 failing to #include 217
 summary of 209–210
 heaps, separate, allocators and 56–58
 Helm, Richard 226
 see also [Design Patterns](#)
 Henney, Kevlin xvi
 “hint” form of insert 100, 110
How the Grinch Stole Christmas! xi

I

identifiers, reserved 213
 identity 219
 implementations
 compilers vs. the STL 3
 variations for string 68–73
 includes 148
 #includes, portability and 209–210
 inlining 202
 function pointers and 203
 inner_product 157
 inplace_merge 148
 input iterators
 definition of 5
 range insert and 29
 insert
 as member template 240
 “hint” form 100, 110
 operator[] in map vs. 106–111
 return types for 32
 insert iterators
 see also [inserter](#), [back_inserter](#), [front_inserter](#)
 container::reserve and 131
 copy algorithm and 26
 insert_after 218
 insert_iterator 216

inserter 130, 216
 inserting
 see also [containers, insertions](#)
 base iterators and 124
 equivalent values in order 198
 in reverse order 184
 rolling back 14
 via range member functions 32
 internationalization
 see also [locales](#)
 strcmp and 150
 stricmp/strcmpi and 154
 invalidation, [see iterators, invalidation](#)
 ios::skipws 126
 iostreams library, SGI implementation
 of 220
 ISO Standard for C++
 see [C++ Standard, The](#)
 istream_iterators 157, 210
 operator>> and 126
 parsing ambiguities and 35
 istreambuf_iterators 157, 210
 use for efficient I/O 126–127
 Items on efficiency, list of 9
 iterator
 comparisons with const iterators 118
 other iterator types vs. 116–119
 reverse_iterator's base and 123–125
 iterator_traits 113
 value_type typedef 42
 iterators
 see also [istreambuf_iterators](#),
 [ostreambuf_iterators](#)
 base, erasing and 124
 base, insertion and 124
 casting 120
 categories 5
 see also [input iterators](#), [output iterator](#),
 [forward iterators](#), [bidirectional it-](#)
 [erators](#), [random access iterators](#)
 in hashed containers 114
 choosing among types 116–119
 conversions among types 117
 dereferencing function object for 91
 implemented as pointers 120
 invalidation
 during deque::insert 185
 during erasing 14, 45, 46
 during insertion 14
 during vector reallocation 59
 during vector/string insert 68
 during vector/string reallocation 66
 during vector::insert 27
 in hand-written loops 185
 in standard sequence containers 17
 in STLport STL implementation 221
 predicting in vector/string 68

 undefined behavior from 27, 45, 46,
 185
 unique rules for deque 15
 pointers in vector vs. 75
 relationship between iterator and
 reverse_iterator's base 123–125
 typedefs for 18–19
 types in containers
 see also [iterator](#), [const_iterator](#),
 [reverse_iterator](#), [const_reverse_iterator](#)
 casting among 121
 types, mixing 119

J

Johnson, Curt xviii
 Johnson, Ralph 226
 see also [Design Patterns](#)
 Johnson, Tim xvii
 Jones, Jason xviii
 Josuttis, Nicolai xv, xvi, xvii, 225, 227
 see also [C++ Standard Library, The](#)

K

Kanze, James xvi
 Kasperski, Marcin xvii
 Kernighan, Brian xvii
 key_comp 85, 110
 keys, for set/multiset, modifying 95–100
 King, Martin Luther, Jr., allusion to 51
 Kreft, Klaus xvi, 228

L

Langer, Angelika xvi, xviii, 228
 leaks, [see resource leaks](#)
 Leary-Coutu, Chanda xviii
 legacy APIs
 filling containers from 77
 sorted vectors and 76
 vector and string and 74–77
 vector<bool> and 81
 lemur 88
 less 210
 operator< and 177–180
 less_equal 92
 Lewandowski, Scott xvi
 lexicographic comparison 231
 lexicographical_compare 153
 strcmp and 153
 use for case-insensitive string
 comparisons 150–154
 lhs, as parameter name 8–9

- linear time complexity 6
 - for binary search algorithms with bidi-
rectional iterators 148
 - logarithmic complexity vs. 190
- list
 - algorithm specializations 192
 - iterator invalidation in, see *iterators*,
invalidation
 - merge 192
 - remove 142–143
 - vs. the *erase-remove idiom* 43
 - sort 137
 - splice
 - exception safety of 50
 - vs. size 23–24
 - unique 143
- Little Engine that Could, The*, allusion
to 63
- local classes
 - definition of 189
 - type parameters and 189
- locales 232–233
 - case-insensitive string comparisons
and 229–237
 - facets and 234–235
- locality of reference 103
 - improving via allocators 55
- locking objects 60
- logarithmic time complexity 6
 - linear complexity vs. 190
 - meaning for binary search
algorithms 147
- longest algorithm name 153
- lookup speed, maximizing 100
- lower_bound
 - equality vs. equivalence and 195
 - related functions vs. 192–201

M

- mailing list for Scott Meyers xiii
- Manning, Jared xvii
- map
 - add or update functionality in 107
 - constness of elements 95
 - iterator invalidation in, see *iterators*,
invalidation
 - key, casting away constness 99
 - key_comp member function 110
 - value_type typedef 108
- Matthew 25:32, allusion to 222
- max_element 157
- max_size 66
- Mayhew, David 226
 - see also *Efficient C++*

- mem_fun
 - declaration for 175
 - reasons for 173–177
 - reference-to-reference problem and 222
- mem_fun_ref
 - reasons for 173–177
 - reference-to-reference problem and 222
- mem_fun_ref_t 175
- mem_fun_t 175
- member funcs, vs. algorithms 190–192
- member function templates
 - see *member templates*
- member templates
 - avoiding client redundancy with 38
 - in the STL 239–240
 - Microsoft's STL platforms and 239–244
 - vector::insert as 240
 - workaround for when missing 242
- memory fragmentation, allocators and 54
- memory layout
 - for string 69–71, 75
 - for vector 74
- memory leaks, see *resource leaks*
- memory, shared, allocators and 55–56
- merge 148, 192
- Meyers, Scott
 - mailing list for xiii
 - web site for xiii
- Microsoft's STL platforms 239–244
 - Dinkumware replacement library
for 243
- microsoft.public.vc.stl xv
- min_element 157
- mismatch 151
 - use for case-insensitive string
comparisons 150–154
- mixing iterator types 119
- modifying
 - components in std 178
 - const objects 99
 - set or multiset keys 95–100
- monomorphic function objects 164
- Moore, Aaron xvii
- More Effective C++*
 - auto_ptr and 40
 - bibliography entry for 225
 - citation number omitted for 3
 - errata list 228
 - smart pointers and 39
 - placement new discussion in 56
 - proxy objects discussion in 49, 80
 - reference counting discussion in 4, 71
 - resource acquisition is initialization
discussion in 61
 - smart pointer discussion in 39

- STL overview in [xi](#)
- URL for auto_ptr update page for [228](#)
- URL for errata list for [228](#)
- More Exceptional C++*
 - bibliography entry for [226](#)
- multimap
 - constness of elements [95](#)
 - find vs. count in [199](#)
 - find vs. lower_bound in [201](#)
 - indeterminate traversal order in [87](#)
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - key, casting away constness [99](#)
 - value_type typedef [108](#)
- multiple deletes [64](#)
- multiplies [159](#)
- multiset
 - constness of elements [95](#)
 - corrupting via element modification [97](#)
 - find vs. count in [199](#)
 - find vs. lower_bound in [201](#)
 - indeterminate traversal order in [87](#)
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - keys, modifying [95–100](#)
- multithreading
 - allocators and [54](#)
 - containers and [58–62](#)
 - reference counting and [64–65](#)
 - string and [64–65](#)

N

- Naran, Siemel [xvi](#)
- newsgroups [xv](#)
 - comp.lang.c++.moderated [xv](#)
 - comp.std.c++ [xv](#)
 - microsoft.public.vc.stl [xv](#)
- node-based containers
 - see also [standard associative containers](#), [list](#), [slist](#), [hashed containers](#)
 - allocators and [52](#)
 - clustering in [103](#)
 - definition of [13](#)
- nonstandard containers
 - see [hashed containers](#), [slist](#), [rope](#)
- not1 [155](#), [156](#), [169](#), [170](#), [172](#), [210](#), [222](#)
 - adaptability and [170](#)
- not2 [152](#), [222](#)
 - adaptability and [170](#)
- nth_element [133–138](#)
- nulls, embedded [75](#), [154](#)
- <numeric> [157](#)

O

- objects
 - copying, in containers [20–22](#)
 - for locking [60](#)
 - slicing [21–22](#), [164](#)
 - temporary, created via casts [98](#)
- One True Editor, the, see [Emacs](#)
- operator new, interface,
 - vs. allocator::allocate [51](#)
- operator() [5](#)
 - declaring const [168](#)
 - functor class and [5](#)
 - inlining and [202](#)
 - overloading
 - adaptability and [173](#)
 - in functor classes [114](#)
- operator++, side effects in [45](#)
- operator--, forward iterators and [5](#)
- operator. (“operator dot”) [49](#)
- operator<, less and [177–180](#)
- operator>>
 - istream_iterators and [126](#)
 - sentry objects and [126](#)
 - whitespace and [126](#)
- operator[], vs. insert in map [106–111](#)
- optimizations
 - algorithms and [183](#)
 - function pointers and [203](#)
 - inlining and [202](#)
 - istreambuf_iterators and [127](#)
 - range insertions and [31](#)
 - reference counting and [64](#)
 - small strings and [71](#)
 - stricmp/strcmpi and [154](#)
 - to reduce default allocator size [70](#)
- ostream, relation to basic_ostream [230](#)
- ostream_iterators [216](#)
- ostreambuf_iterators [216](#)
 - efficiency and [127](#)
- output iterator, definition of [5](#)
- overloading, operator() in functor
 - classes [114](#)

P

- page faults [102](#), [103](#)
- pair, comparison functions for [104](#)
- parameters
 - function objects vs. functions [201–205](#)
 - pointers to functions [34](#)
 - type, local classes and [189](#)
- parentheses
 - ignored, around parameter names [33](#)
 - to distinguish function and object
 - declarations [35](#)

- parse, most vexing in C++ 33–35
- parsing, objects vs. functions 33–35
- partial_sort 133–138
- partial_sum 157
- partition 133–138
 - containers of pointers and 145
 - remove vs. 141
- pass-by-value
 - function objects and 163
 - functor classes and 162–166
- penguin 88
- Perfect Woman, see [Woman, Perfect](#)
- performance, see [efficiency](#)
- Persephone xviii
- Pimpl Idiom 165
- placement new 56
- Plauger, P. J. xvi, 114, 227
- pointers
 - allocator typedef for 48
 - as iterators 120
 - as return type from `vector::begin` 75
 - assignments, avoiding superfluous 31
 - comparison functions for 88–91
 - deleting in containers 36–40
 - dereferencing function object for 91
 - destructors for 36
 - invalidation, see [iterators, invalidation](#)
 - iterators in vector vs. 75
 - parameters, `binary_function` and 172
 - parameters, `unary_function` and 172
 - returned from `reverse_iterator::base` 125
 - smart, see [smart pointers](#)
 - to bitfields 80
 - to functions, as formal parameters 34
- portability
 - #includes and 209–210
 - casting `const` iterator to iterators 121
 - `container::auto_ptr` and 41
 - explicit template argument specification and 163
 - hashed containers, code using 112
 - identity, `project1st`, `project2nd`, `compose1`, `compose2`, `select1st`, `select2nd` and 219
 - multiple compilers and 3
 - range construction with
 - `istream_iterators` and 35
 - `reverse_iterator::base` and 125
 - set/multiset key modification and 98
 - stateful allocators and 50, 51
 - STLport STL implementation and 220
 - `stricmp/strcmpl` and 154
- Potter, John xvi, xvi
- predicate class, definition of 166
- predicates
 - definition of 166
 - need to be pure functions 166–169

- predicting iterator invalidation, in vector/string 68
- principle of least astonishment, the 179
- priority_queue 138
 - as part of the STL 2
- project1st 219
- project2nd 219
- proxy objects 49
 - containers of 82
 - `vector<bool>` and 80
- `ptr_fun`, reasons for 173–177
- pure function, definition of 166
- `push_back`, `back_inserter` and 130
- `push_front`, `front_inserter` and 130

Q

- qsort 162
 - declaration for 162
 - sort vs. 203
- queue, as part of the STL 2

R

- Rabinowitz, Marty xviii
- random access iterators
 - definition of 5
 - sorting algorithms requiring 137
- range
 - destination, ensuring adequate size 129–133
 - member functions 25
 - input iterators and 29
 - single-element versions vs. 24–33
 - summary of 31–33
 - pointer assignments in list and 31
 - sorted, algorithms requiring 146–150
 - summarizing 156–161
- `raw_storage_iterator` 52
- RB trees, see [red-black trees](#)
- `rbegin/rend`, relation to `begin/end` 123
- reallocations
 - invalidation of iterators during 59
 - minimizing via `reserve` 66–68
- rebinding allocators 53
- red-black trees 190, 191, 214
- redundant computations, avoiding via
 - algorithm calls 182
- Reeves, Jack xvi, 227
- reference counting
 - disabling, for string 65
 - multithreading and 64–65
 - smart pointers 39, 146
 - see also [Boost](#), [shared_ptr](#)
 - string and 64–65

The C++ Standard and 64
 this book and 4

references

- allocator typedef for 48
- casting to 98
- invalidation, see [iterators, invalidation](#)
- parameters, [binary_function](#) and 171
- parameters, [unary_function](#) and 171
- to bitfields 80

reference-to-reference problem 222

remove 139–142

- see also [erase-remove idiom](#)
- on containers of pointers 143–146
- partition vs. 141

[remove_copy_if](#) 44

[remove_if](#) 44, 144, 145

- see also [remove](#)
- possible implementation of 167

[replace_if](#) 186

replacing STL implementations 243

reporting bugs in this book [xii](#)

reserve

- insert iterators and 131
- resize vs. 67

resize

- reallocation and 67
- reserve vs. 67

resource acquisition is initialization 61

resource leaks 36, 39, 63, 144, 145

- avoiding via smart pointers 39, 146
- preventing via classes 61

[result_type](#) 170

return type

- [allocator::allocate](#) vs. [operator new](#) 51
- for [container::begin](#) 75
- for distance 122
- for erase 32, 117
- for function objects for [accumulate](#) 158
- for insert 17, 32, 117
- for [vector::operator\[\]](#) 80

reverse order

- inserting elements in 184

[reverse_iterator](#)

- base member function 123–125
- other iterator types vs. 116–119

rhs, as parameter name 8–9

Rodgers, Mark [xv](#), [xvi](#), [xvii](#)

rolling back, insertions and erasures 14

rope 218

S

Scheme 206

[second_argument_type](#) 170

[select1st](#) 219

[select2nd](#) 219

sentry objects, [operator<<](#) and 126

separate heaps, allocators and 56–58

sequence containers

- see [standard sequence containers](#)

set

- constness of elements 95
- corrupting via element modification 97
- iterator invalidation in, see [iterators, invalidation](#)
- keys, modifying 95–100
- membership test, idiomatic 199

[set_difference](#) 148

[set_intersection](#) 148

[set_symmetric_difference](#) 148, 153

[set_union](#) 148

[sgetc](#) 127

SGI

- hashed containers implementation 114
- iostreams implementation 220
- slist implementation 218
- STL web site 94, 207, 217–220
- thread-safety definition at 58
- URL for 217, 227

shared memory, allocators and 55–56

[shared_ptr](#), see [Boost, shared_ptr](#)

shrink to fit, see [swap trick, the](#)

side effects

- [accumulate](#) and 160
- for [each](#) and 160
- in [operator++](#) 45

size

- vs. capacity 66–67
- vs. empty 23–24

[size_type](#) 158

sizeof, variations when applied to string 69

[skipws](#) 126

slicing problem 21–22, 164

slist 218

small string optimization 71

Smallberg, David [xvi](#), [xvii](#)

smart pointers

- see also [Boost, shared_ptr](#)
- avoiding resource leaks with 39, 146
- dereferencing function object for 91
- implicit conversions and 146

sort 133–138

- [qsort](#) vs. 203

sorted range

- algorithms requiring 146–150

sorted vectors

- associative containers vs. 100–106
- legacy APIs and 76

sorting

- algorithms for 133–138
- [auto_ptrs](#) 41
- consistency and 149

- stability, in sorting 135
- stable_partition 133–138
- stable_sort 133–138
- stack, as part of the STL 2
- Staley, Abbi xviii
- standard associative containers
 - see also [containers](#)
 - bidirectional iterators and 5
 - comparison funcs for pointers 88–91
 - definition of 5
 - “hint” form of insert 100, 110
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - key_comp member function 85
 - search complexity of 6
 - sorted vector vs. 100–106
 - typical implementation 52, 190
- Standard for C++
 - see [C++ Standard, The](#)
- standard sequence containers
 - see also [containers](#)
 - definition of 5
 - erase’s return value 46
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - push_back, back_inserter and 130
 - push_front, front_inserter and 130
- Standard Template Library, see [STL](#)
- Stepanov, Alexander 201
- STL
 - algorithms, vs. string member functions 230
 - arrays and 2
 - bitset and 2
 - complexity guarantees in 5–6
 - container adapters and 2
 - containers, selecting among 11–15
 - definition of 2
 - documentation, on-line 217
 - early usage problems with xi
 - extending 2
 - free implementations of 217, 220
 - function call syntax in 175
 - implementations
 - compiler implementations vs. 3
 - Dinkumware bug list for MSVC 244
 - replacing 243
 - member templates in 239–240
 - platform, see [STL platform](#)
 - priority_queue and 2
 - queue and 2
 - stack and 2
 - thread safety and 58–62
 - valarray and 2
 - web sites about 217–223
 - wide-character strings and 4
- STL platform
 - definition of 3
 - Microsoft’s, remarks on 239–244
- STLport 220–221
 - debug mode 185, 216
 - detecting invalidated iterators in 221
 - hashed containers at 112
 - URL for 217
- strcmp 152, 234
 - internationalization issues and 150
 - lexicographical_compare and 153
- strncmp 154
- streams, relation to string 230
- stricmp 154
- strict weak ordering 94
- string
 - allocators in 69
 - alternatives to 65
 - arrays vs. 63–66
 - as typedef for basic_string 65
 - c_str member function 75
 - comparisons
 - case-insensitive 150–154, 235–237
 - using locales 229–237
 - cost of increasing capacity 66
 - disabling reference counting 65
 - embedded nulls and 75, 154
 - growth factor 66
 - implementation variations 68–73
 - inheriting from 37
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - iterators as pointers 120
 - legacy APIs and 74–77
 - mem funcs vs. algorithms 230
 - memory layout for 75
 - minimum allocation for 72
 - multithreading and 64–65
 - reference counting and 64–65
 - relation to basic_string 4, 210, 229
 - relation to streams 230
 - reserve, input iterators and 131
 - resize vs. reserve 67
 - shrink to fit 78–79
 - size vs. capacity 66–67
 - size_type typedef 158
 - sizeof, variations in 69
 - small, optimization for 71
 - summing lengths of 158
 - trimming extra capacity from 77–79
 - vector vs. 64
 - vector<char> vs. 65
 - whether reference-counted 65
 - wstring and 4
- string_char_traits 211

strings

- case-insensitive 229–230
- wide-character, see *wstring*

Stroustrup, Bjarne xvi, 68, 226, 228
see also *C++ Programming Language, The*

structs, vs. classes for functor classes 171

summarizing ranges 156–161

Sutter, Herb xvi, xvii, 65, 226, 227, 228
see also *Exceptional C++*

swap trick, the 77–79

T

templates

- declaring inside functions 188
- explicit argument specification for 122, 163, 234

instantiating with local classes 189

member

in the STL 239–240

Microsoft's platforms and 239–244
parameters, declared via class vs. *typename* 7

temporary objects, created via casts 98

thread safety, in containers 58–62

see also *multithreading*

tolower 151

as inverse of *toupper* 235

toupper

- as inverse of *tolower* 235
- cost of calling 236–237

traits classes 113, 211, 230

transactional semantics 14

transform 129, 186

traversal order, in multiset, multimap 87

trees, red-black 190, 191, 214

typedefs

- allocator::pointer* 48
- allocator::reference* 48
- argument_type* 170
- container::size_type* 158
- container::value_type* 36
- first_argument_type* 170
- for container and iterator types 18–19
- mem_fun* and 176
- mem_fun_ref* and 176
- ptr_fun* and 170
- result_type* 170
- second_argument_type* 170
- string as 65
- wstring* as 65

typename

- class vs. 7
- dependent types and 7–8

U

unary_function 170–172

- pointer parameters and 172
- reference parameters and 171

undefined behavior

- accessing *v[0]* when *v* is empty 74
- applying some algorithms to ranges of unsorted values 147
- associative container comparison funcs yielding true for equal values 92
- attempting to modify objects defined to be const 99
- changing a set or multiset key 97
- deleting an object more than once 64
- deleting derived object via *ptr-to-base* with a nonvirtual destructor 38
- detecting via STLport's debug mode 220–221
- modifying components in *std* 178
- multithreading and 59
- practical meaning of 3–4
- side effects inside *accumulate*'s function object 160
- specifying uninitialized memory as destination range for algorithms 132
- using algorithms with inconsistent sort orders 149
- using the wrong form of *delete* 64
- when using invalidated iterator 27, 45, 46, 185

underscores, in identifiers 213

uninitialized_fill 52

uniq 148

unique 145, 148, 149

see also *remove*

unique_copy 148, 149

unsigned char, use in *<cctype>* and *<cctype.h>* 151

upper_bound 197–198

related functions vs. 192–201

Urbano, Nancy L., see *Perfect Woman*

URLs

- for Austern's sample allocator 228
- for *auto_ptr* update page 228
- for Boost web site 217, 227
- for Dinkumware web site 243
- for *Effective C++ CD* errata list 228
- for *Effective C++* errata list 228
- for *Effective STL* errata list xii
- for Josuttis' sample allocator 227
- for *More Effective C++* errata list 228
- for Persephone's web site 71
- for purchasing The C++ Standard 226
- for Scott Meyers' mailing list xiii
- for Scott Meyers' web site xiii
- for SGI STL web site 217, 227

- for STLport web site [217](#)
- for this book's errata list [xii](#)
- for this book's web site [xii](#)
- use_facet [234](#)
 - cost of calling [234](#)
 - explicit template argument specification and [234](#)
- Usenet newsgroups, see [newsgroups](#)

V

- valarray, as part of the STL [2](#)
- value_type typedef
 - in containers [36](#), [108](#)
 - in iterator_traits [42](#)
- vector
 - see also [vector<bool>](#), [vector<char>](#)
 - arrays vs. [63–66](#)
 - contiguous memory for [74](#)
 - cost of increasing capacity [66](#)
 - growth factor [66](#)
 - iterator invalidation in, see [iterators](#), [invalidation](#)
 - iterators as pointers [120](#)
 - legacy APIs and [74–77](#)
 - reserve, input iterators and [131](#)
 - resize vs. reserve [67](#)
 - return type from begin [75](#)
 - shrink to fit [78–79](#)
 - size vs. capacity [66–67](#)
 - sorted
 - legacy APIs and [76](#)
 - vs. associative containers [100–106](#)
 - string vs. [64](#)
 - trimming extra capacity from [77–79](#)
- vector::insert, as member template [240](#)
- vector<bool>
 - alternatives to [81](#)
 - legacy APIs and [81](#)
 - problems with [79–82](#)
 - proxy objects and [80](#)
- vector<char>, vs. string [65](#)
- virtual functions, toupper and [236–237](#)
- Visual Basic [127](#)
- Visual C++, see [Microsoft's STL platforms](#)
- Vlissides, John [226](#)
 - see also [Design Patterns](#)
- vocabulary, algorithms as [186](#)
- Von, Victor [xvii](#)

W

- Wait, John [xviii](#)
- wchar_t [4](#)
- web sites
 - see also [URLs](#)
 - for STL-related resources [217–223](#)

- whitespace, operator<< and [126](#)
- wide-character strings, see [wstring](#)
- Widget, use in this book [9](#)
- Wizard of Oz, The [83](#)
- Woman, Perfect, see [Urbano, Nancy L.](#)
- wombat [88](#)
- workarounds
 - for improperly declared iterator-related functions [119](#)
 - for Microsoft's STL platforms [242–244](#)
 - for missing member templates [242](#)
 - for use_facet [234](#)
 - function objects as [204](#)
- write-only code, avoiding [206–208](#)
- writing the author [xii](#)
- wstring
 - as typedef for basic_string [65](#)
 - relation to basic_string [4](#)
 - string and [4](#)
 - this book and [4](#)

Y

- Yelle, Dennis [xvii](#)

Z

- Zolman, Leor [xvii](#), [211](#), [228](#)

